ANSWERS.

- 1. CLR 37-1.
- 2. CLR 31.2-2.
- 3. CLR 31.2-4.
- 4. In class we analyzed the effect of *blocking* on reducing the number of slow memory accesses in conventional matrix multiplication: we showed that if we multiply *s*-by-*s* blocks in the inner loop of blocked matrix multiplication then the number of slow memory accesses dropped from $2n^3 + 2n^2$ to $2n^3/s + 2n^2$. Extend this analysis to the case of multiplying $C = C + A \cdot B$ where A is n-by-k and B is k-by-m. You may assume that s divides m, n and k. You should turn in your algorithm, indicating when data is moved between fast and slow memory, and show your analysis counting the number of such data moves.
- 5. CLR 31.2-6.
- 6. How many *real* arithmetic operations (adds and multiplies) does it takes to multiply 2 complex *n*-by-*n* matrices, where complex arithmetic is implemented in the most s-traightforward way? Apply the idea of the last question to show how to do it more cheaply.