

## Problem Set 1

### 1. [ $\epsilon$ -optimal strategies]

A strategy pair  $(x, y)$  for a zero sum game is  $\epsilon$ -optimal if a player can gain at most  $\epsilon$  by deviating from his strategy assuming that the other player does not deviate,

- (a) Given a strategy pair  $(x, y)$  suggest an efficient procedure to verify that  $(x, y)$  is  $\epsilon$ -optimal.
- (b) Show that the value  $x^t A y$  of an  $\epsilon$  optimal strategy is within  $\epsilon$  of the value of the game.

### 2. [Irrelevant Attributes]

Given a function of  $n$  variables, which is, in fact, the OR of  $r$  of them, where  $r \ll n$ . Consider an algorithm to predict the function which maintains weights  $w_1, \dots, w_n$  and predicts 1 when  $w_1 x_1 + \dots + w_n x_n \geq n$  and 0 otherwise. Each time the prediction is incorrect halve or double the appropriate weights. Complete the description of the algorithm and show it makes  $O(r \log n)$  mistakes.

### 3. [The matching game]

The matching game is played over the complete weighted bipartite graph  $G(V, E)$ . The edge player plays an edge  $e \in E$  while the vertex player plays a vertex  $v \in V$  and if  $v \in e$  the edge player pays  $1/w_e$  to the vertex player.

- (a) If the vertex player plays a uniformly random vertex what is the best response for the edge player?
- (b) If the edge player plays a uniformly random edge from the maximum weight matching what is the best response for the vertex player?
- (c) Are these two strategies optimal for this game? If not, what are a pair of strategies that are optimal for this game.
- (d) Show how to use the multiplicative weights algorithm to find a fractional matching with cost within  $(1 + \epsilon)$  factor of optimal. (A fractional matching is an assignment  $x_e$  to the edges such that for each vertex  $\sum_{e \ni u} x_e \leq 1$ .)