Browsing Web 3.0 on 3.0 Watts
or Why Browsers Will Be Parallel and Implications for Education

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UC Berkeley
Milestones in computing?

What are your top 3 milestones in computing?
- enabled new ways of thinking/doing things

Some candidates:
- hard disks and databases
- languages and their compilers
- the transistor
- theory of NP-completeness
- Web/internet
- the PC
Bell’s Law of Computer Classes
Moore’s Law Enables Two Evolution Paths

**Path 1**: increasing performance, same cost and form factor

**Path 2**: constant performance, decreasing cost and form factor
Bell’s Law: a Corollary of Moore’s Law

Mainframe

Mini

WS

PC

Laptop

Handset

Ubiquitous

computing in the Cloud

price/form factor

time

slide courtesy Dave Patterson
Handheld’s Killer Apps

The laptop→handheld transition will happen if...
  - handhelds do to laptops what laptops did to desktops

Handheld killer apps that will make laptops irrelevant:

1. Device convergence
   - phone + music player + PDA + gaming
2. Laptop replacement
   - small form factor = convenience
3. Personal assistant
   - vision, hearing, memory, health improvement
Handheld as a Laptop Replacement

- When will handhelds take over?
  - when they have laptop-quality browser

- Why is a web browser enough?
  - In Web 3.0: most apps will be browser-based
  - In Web 2.0: some apps are still native (Picassa, Google Earth)
A handheld browser may be soon possible

A guy walks into a bar, asks for a cup, and starts his browser.
Display: cell phone projector or “wearable”

Texas Instruments, CES 2008
Input: idea for tablet input for a handheld

- Inspiration: mimio, a whiteboard recorder (mimio.com)

How mimio works:
- triangulates in the same way that one measures lightning distance

1. marker simulates a lightning strike: simultaneously emits light and sound signals;
2. capture bar measures sound travel time: yields marker distance to each mic;
3. the two distances determine marker location on the whiteboard; goto step 1
Dasher + picomimio = keyboard-rate input

Dasher: replacement for traditional keyboards

- Input rates up to ~30 words/minute
- Only needs 1 input axis (up/down) to work
  - can be controlled by picomimio, eyes, tilt sensor, ...

See http://www.inference.phy.cam.ac.uk/dasher/ for more info, online demo
Other input alternatives

- speech recognition
  - we’ll have enough MIPS for low-noise environments
- sensors for gestures
  - iPhone has sensors: ex. to zoom into a map, move it so
- virtual laser keyboard:
What about CPU performance?

**Display:** many alternatives

**Input:** half as many

**Network:** plenty fast soon (all we need is better providers)

**CPU speed:** no longer considered a bottleneck. Is it true?
Why is iPhone slow? Network or CPU?

Experiment: time to load+render cnn.com on two networks

<table>
<thead>
<tr>
<th>machine</th>
<th>seconds</th>
</tr>
</thead>
<tbody>
<tr>
<td>a modern desktop (2Mbps network)</td>
<td>2</td>
</tr>
<tr>
<td>T40 1.6GHz (a very old laptop; 2Mbps network)</td>
<td>7</td>
</tr>
<tr>
<td>T40 1.6Ghz (same laptop&amp;network, battery mode)</td>
<td>13</td>
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<tr>
<td>iPhone 600MHz (2Mbps network)</td>
<td>37</td>
</tr>
<tr>
<td>iPhone 600MHz (1Mbps network)</td>
<td>40</td>
</tr>
</tbody>
</table>

♫ CPU performance is critical
Uniprocessor Performance (SPECint)

Proxy Web Servers?

Recently introduced Skyfire browser uses a proxy server
  - server parses the web page, sends images to handheld
  - not unlike the X11 client/server architecture

Limitations:
  - network latency: impact on interactivity
  - radio power: may be cheaper to decompress than to receive
  - server CPU cycles needed: average browser CPU load is 5%!
End of Bell’s Law?

Undergraduate Course in PL and Compilers
Is our field slowly dying?

Observations from graduate admissions:
- AI, bio, NLP, OS much better at recruiting the brightest minds

This may have roots in undergraduate education

Enrollments at Berkeley:
- Algo=165; OS=155; AI=160; DB=110
- compare with compilers=60

The trend goes beyond Berkeley, I suspect
- Amazon rank for Silberschatz =800
- for Dragon Book=8,000
Why this lack of interest?

So much excitement
  - eg, Web languages and frameworks

So many interesting open problems remain
  - eg, parallelism, speeding up scripting languages

Recognition: 6 out of last 10 Turing Awards to our discipline
  - perhaps this is a sign of “mature and solved”?

Is AI viewed as the hero of web2.0 success (web search)?
  - while web2.0 is arguably a programming success

Education problem? Political mistakes? Marketing?
  - No longer a required “core” course?
Relevance?

My survey why people take AI not compilers
  - compilers “too hard”, “solved, old stuff”

OS textbook revised more often
  - Silberschatz in 7th edition
  - Dragon in 2nd edition

I looked at a few OS lectures
  - was stuff I wanted to learn! (eg PKK, peer-to-peer networks)
  - compare this with register allocation
End of Bell’s Law?

Course in PL and Compilers

Browsers for Handhelds
Good news, bad news

Current laptops: how powerful?

~20GOPS at 20W, more with SIMD

Good news: we should be able to get 50GOPS at 2W

- even in current 65nm technology (40pJ/op)

Bad news: the 50GOPS will come from 10-100 parallel cores

✈️✈️✈️ We must build a parallel browser
Transition to handhelds qualitatively different

Power Wall: Previous Bell steps were easy. We can now no longer wait for smaller, lower power processors. Instead, software must be parallelized.
Browser autopsy: three major components

- web page “compilation”: HTML to DOM
  - lexicing, parsing, and syntax-directed translation
  - our project: parallelize all three

- scripting
  - JavaScript (AJAX)
  - our project: a dataflow-ish web client language

- page layout and rendering:
  - similar to formatting a latex document
  - our project: design web language hand in hand with layout/rendering

- But these bottlenecks invoke Amdahl's Law
  - current browsers “inherently” sequential
  - for neither of these components, parallel successes exist
End of Bell’s Law?

Parallel Browser

Course in PL and Compilers

How to modernize the course?
Experience from Berkeley

- My students and I have been remodeling the undergraduate PL and compiler course since 2003
- Two major revisions, each offered twice
- Student survey results on “usefulness of the course”:
  
<table>
<thead>
<tr>
<th>Term</th>
<th>Rating</th>
</tr>
</thead>
<tbody>
<tr>
<td>Fall 2003</td>
<td>5.4/7.0</td>
</tr>
<tr>
<td>Fall 2004</td>
<td>5.8</td>
</tr>
<tr>
<td>Spring 2007</td>
<td>5.9</td>
</tr>
<tr>
<td>Fall 2007</td>
<td>6.3</td>
</tr>
</tbody>
</table>

  ← 15-year high for the course

No less “useful” than OS, AI. Still limping behind architecture.
How to revise course?

Our initial thinking in 2003: *Modernize it!*
- incorporate latest research: latest analyses, type systems

Then the key question arose: *Who is the target audience?*
- compiler writers or software engineers?

But do programmers need language technologies?
- surprisingly yes

But the content needs to go back to basics
Why do you want to take the course?

Write code that writes code.
- don’t write "mindless" code, write a generator
- ex: database schema → Java wrapper classes

New languages will keep coming. Be ready.
- CSS, HTML, JavaScript, JSP, PHP, Python, Ruby, XML
- know advantages of picking a particular language

Develop your own language.
- Many web languages created practically in the garage, not in a research lab: PHP, JavaScript, Ruby, perl

Learn about compilers and interpreters.
- language is the most important of all programmer tools
- Typical job ad: if you had to take one course, it’d be PL+compilers
“We design them here, but the labor is cheaper in Hell.”

languages, frameworks, APIs
Change 1: Raise Level of Abstraction

What are the source/target programming model?

Before: compile Java-like $\rightarrow$ x86-like
Now: compile a Dataflow extension to AJAX $\rightarrow$ JavaScript
Change 2: Language landscape is broader today

- Dynamic languages in wide use
- Interpreters in wide use

- Programmer choice of data structures in dynamic, interpreted makes huge difference on performance
Change 3: Back to Basics

Some course content decisions go back 30 years
- when computers were 1,000,000-times slower
- when those topics were research topics

Go back to simple parsers: CYK and Earley
- CYK can be explained 10-times faster than LALR
- easy to implement → leads to better understanding
End of Bell’s Law?

Parallel Browser

Course in PL and Compilers

Back to basics and to parallelism
End of Bell’s Law?

Parallel Browser

Course in PL and Compilers

Back to basics and to parallelism

Regular Expressions
Parsing
Web Scripting
Applications
What fraction is HTML compilation?

- 10-40% of time spent in lexing, parsing, syntax-directed translation
  - (remainder in layout/rendering)
- loading a fark.com page from disk cache; little JavaScript
- on (old) debug build of Firefox 3

HTML compilation must be parallelized
Lexing, from 10,000 feet

**Goal:** given lexical specification and input, find lexemes

Content ::= [^<]+
ETag ::= */[^>]*>
STag ::= <[^>]*>
Problem: lexing seems “inherently sequential”

- To know automaton state at input position $i+1$
- We need to know the automaton state at position $i$!
Ideal solution

- Divide input among the processors
- For each processor starting at position $i+1$
  - Ask an oracle in which state the neighbor at $i$ finished
  - Scan in parallel from next state, at $i+1$
- Finally, merge the results
Practical solution: guess! (speculate)

How can we guess from position $i+1$ the state at position $i$?

1. Assume state(i) could have been any automaton state
   - 😊 the "speculation" is always correct (not really a guess)
   - 😊 can yield $O (\log n)$ algorithm [Hillis and Steele] ...
   - 😞 ... but prohibitively expensive in practice

2. Was one of a "likely set" of automaton states
   - 😊 can be more efficient than algorithm (1)
   - 😊 can fine-tune speculation based on language and workload
   - 😞 speculation can be wrong
   - 😞 still can be expensive (memory overhead, bad guesses)

But we can do better ...
Our solution (1/2)

Observation: In “real” lexers, the DFA converges to a stable, recurring state (often, the “start state”), from multiple initial states, after a small number $k$ of characters.

Lexing: < b > Berkeley ! < / b >

From:

- "start", s0
  - $s_1$ $s_4$ $s_5$

- "in STag", $s_4$
  - $s_4$ $s_4$ $s_5$

- "in ETag", $s_2$
  - $s_2$ $s_2$ $s_3$

- "in Content", $s_6$
  - $s_1$ $s_4$ $s_5$

 fsm

$k$

Only need to follow one DFA path instead of several.
Our solution (2/2)

- **Sketch of our algorithm:**
  - split input into blocks with $k$-character overlap
  - scan blocks in parallel, each starting from "good" initial state
Our solution (2/2)

- **Sketch of our algorithm:**
  - split input into blocks with $k$-character overlap
  - scan blocks in parallel, each starting from “good” initial state
  - find if blocks converge: expected in $k$-overlap
  - speculation may fail; if so, block is rescanned
Preliminary results: speedup over flex

- **flex**: optimized, single-thread lexer on fat Cell core
- Speedup computed by *flex time / cellex time*

**future page sizes**: 5 cores are 6x faster than flex

**today's page sizes**: 5 cores are 4.5x faster than flex
Regular expressions

Most teachers, happily cut RE to a single lecture.

We just saw that automata offer fun algorithmic exercises.

Are regular expressions also useful?
  – probably the most frequently embedded language
  – ex: mashups (extract info from HTML pages, …)

Embedding presents language semantics challenges
  – consider the following example …
Imagine you want to parse a config file:
```
filesToCompile=a.cpp b.cpp
```
The regex:
```
[a-zA-Z]+=.*
```
Now let’s allow an optional newline-separated 2\textsuperscript{nd} line:
```
filesToCompile=a.cpp b.cpp \[NL] d.cpp e.h
```
We extend the original regex:
```
[a-zA-Z]+=.*(\\\n.*)?
```
This regex does not match our two-line input. Why?
What compiler textbooks don’t teach you

First, the string matching problem in the Textbook World:
- “Does a regex \( r \) match the entire string \( s \)?”
- this is a clean statement and suitable for theoretical study
- here is where equivalence of regex and FSM is defined

The matching problem in the Real World:
- “Given a string \( s \) and a regex \( r \), find a substring in \( s \) matching \( r \).”

Do you see the language design issue here?
- There may be many such substrings.
- We need to decide what substring to find.

It is easy to agree where the substring should start:
- the matched substring should be the leftmost match
Two schools of Real World regexes

They differ in where it should end:

**Declarative approach:** return the longest of all matches
  - conceptually, enumerates all matches and returns longest

**Operational approach:** define behavior of *, | operators
  e*  match e maximally such that remainder of regex matches
  e|e  select leftmost choice while allowing remainder to match

```plaintext
filesToCompile=a.cpp b.cpp \[NL\] d.cpp e.h
[a-zA-Z]+=.*(\n.*)?
```
These are important differences

- We saw that a non-contrived regex behaves differently
  - personal story: I spent 3 hours debugging a similar regex
  - despite reading the manual carefully

- The operational semantics of *
  - does not guarantee longest match
  - forces the programmer to reason about backtracking

- It seems that backtracking is nice to reason about
  - because it’s local: no need to consider the entire regex
  - the cognitive load is actually higher, as it breaks composition
Lessons: WTH did things go wrong?

Likely problem:
- long time ago, someone did not know that NFA can find all matches simultaneously and/or
- NFA can be implemented efficiently

I would like to blame Perl,
- but this regex semantics seems older

The theory of finite automata is elegant
- a big success of computer science
- we should make sure that our students know it
- and design clean languages
CYK Parser

Simple context-free-language parser

- running time is $O(n^3)$, space $O(n^2)$

Shunned for many years.

“Even tabular methods [CYK, Earley] should be avoided if the language at hand has a grammar for which more efficient algorithms [LL, LALR] are available.” The Theory of Parsing …, Aho, Ullman, 1972

But in practice, running time is more like $\theta(n^{1.2})$

- plus computers are now 1,000,000-times faster than in 1972
- browser: we plan to parallelize CYK + Earley parser
Parsing, from 10,000 Feet

Doc → STag ETag
ETag → ‘</ Name ’>
STag → ‘< Name ’>
Name → [id]

<html></html>
**CYK**

- reduction = adding a non-terminal edge
- reduce until start symbol added

```
Doc  →  Open  Close
Open  →  '<'  Name  '>
Close  →  '</'  Name  '>
Name  →  [id]
```
Parallel CYK

- CYK is inherently parallel: lots of independent work
  - Bad for serial processing, great for parallel
  - Parse graph is lock-free sparse array \((n \times n \times |G|)\)

edge added when parallel results are combined.

independent work, done in parallel
Parsing in the course

Back to basics: from LALR to CYK
  - I don’t have to relearn LALR each time 😊

Students used CYK to build their own bison
  - including declarative disambiguation, accepts all grammars
  - teaches ambiguous grammars, needed for real world tasks
  - google calculator:
    - 34 knots in km/h
    - half a dozen pints * (110 Calories per 12 fl oz) / 25W in days
    - 5 in in in
  - live programming (gcac, PHP in 20 LOC)

More on parallelism in parsing
  - rethink attribute grammars
A new web client language

AJAX reactive programming is based on clunky callbacks
- too much “plumbing” in the code
- hard to parallelize

Dataflow is a cleaner abstraction

As an example, consider this “follower” program
- where a box follows mouse, with a delay
AJAX code: callbacks

```html
<div id="box" style="position:absolute; background: black;">  
  Seconds to deadline: <span id="time"> ... time ... </span>
</div>

<script>
  document.addEventListener('mousemove', function (e) {
    var left = e.pageX;
    var top = e.pageY;
    setTimeout(function() {
      document.getElementById("box").style.top = top;
      document.getElementById("box").style.left = left;
    }, 500);
  }, false);
</script>
```

This code moves the box with a delay. We need to set up two nested callbacks. We need to refer to the DOM explicitly by element ids. Code does make it clear at all what’s going on. How would you parallelize the program if multiple box were moving on the screen?
**FlapJax code: from callbacks to streams**

- Program is clearer when data flow in it directly exposed
  - in dataflow version: changing mouse coordinates are streams
  - coordinate streams adjust box position after they are delayed
  - another stream (time) adjusts text after it is formatted
Future apps

- Future web apps will be like desktop apps and more …
  - browser = the new windows manager ➔ tabs outdated
  - browser = new OS (local storage, refined security policies)
  - new usage modes (multi-touch, camera-based input, data)

- We want to identify domains that a browser can support
  - hypertext documents and media
  - office suites
  - simpler games
  - rich visualization, for data presentation (eg search results)
Example 1: Baryl Desktop Manager

- 3D desktop with physical properties
Example 2: ManyTube mockup demo

- Example of a new media app
Example 3: OS X Time Machine

An early example of visualizing time-varying data.
Example 4: Multi-touch interfaces

- http://www.youtube.com/watch?v=ysEVYwa-
Example 5: Stereoscopic displays (VR)

- May force us to rethink the desktop metaphor
What should the programming model support?

QoS:
- latency specifications for GUI responsiveness
- video frame rate, etc

Animation with physical properties (both GUI and games)
- property changes over time, stated declaratively
- trajectories: how to declare them?
- physical properties: stretching, gravity, friction, but maybe also flow, fracture

2.5D and/or 3D
- web page = logical structure + script-produced 3D view?
  - What will a 3D nytimes.com look like? Will 3D ease browsing?
- Q: how to project a part of 3D scene for 2D viewing/reading?
Power Wall: Previous Bell steps were easy. To make the net step, however, we cannot wait for smaller, lower power processors. Instead, software must be parallelized.
Conclusion: parallelism

If parallel parsing, layout, and scripting succeeds, browsing on the handheld could be as rich as browsing on the desktop …

… and the next Bell “computer class” will happen.

And all this thanks to advances in languages and compilation.
Conclusion: undergraduate course

Back to basics

Simpler to teach

Allows students invent new technology, rather than just learn about it

Need to rethink the course also for parallelism
Backup slides
Course tricks

- debug professor’s code
- submit test cases for grading on a curve
- reinvent CYK
- live programming
Applications

- Drive language development
  - both small in the course
  - and future AJAX in Par Lab

- Course
  - google calculator
  - mashups
    - animation, interaction with flickr

- Future web applications: challenge problems for web client language design