Performance Variability in Hadoop's Map Reduce Charles Reiss (Advisor: Randy Katz)

MOTIVATION / METHODOLOGY

Goal: Characterize performance inconsistency in MapReduce frameworks that lead to poor interactive performance and scaling.

Data source:

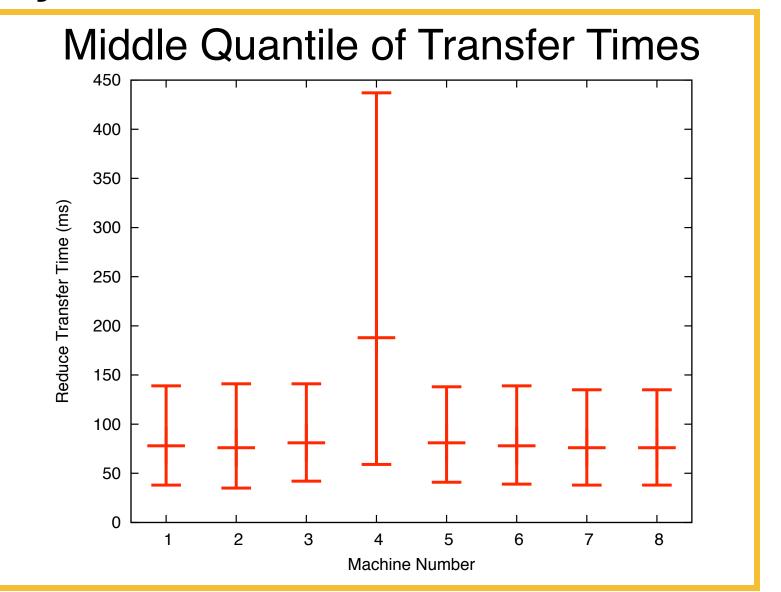
RADLab

•Sort jobs using identity map and reduce functions

•Job is mostly framework overhead

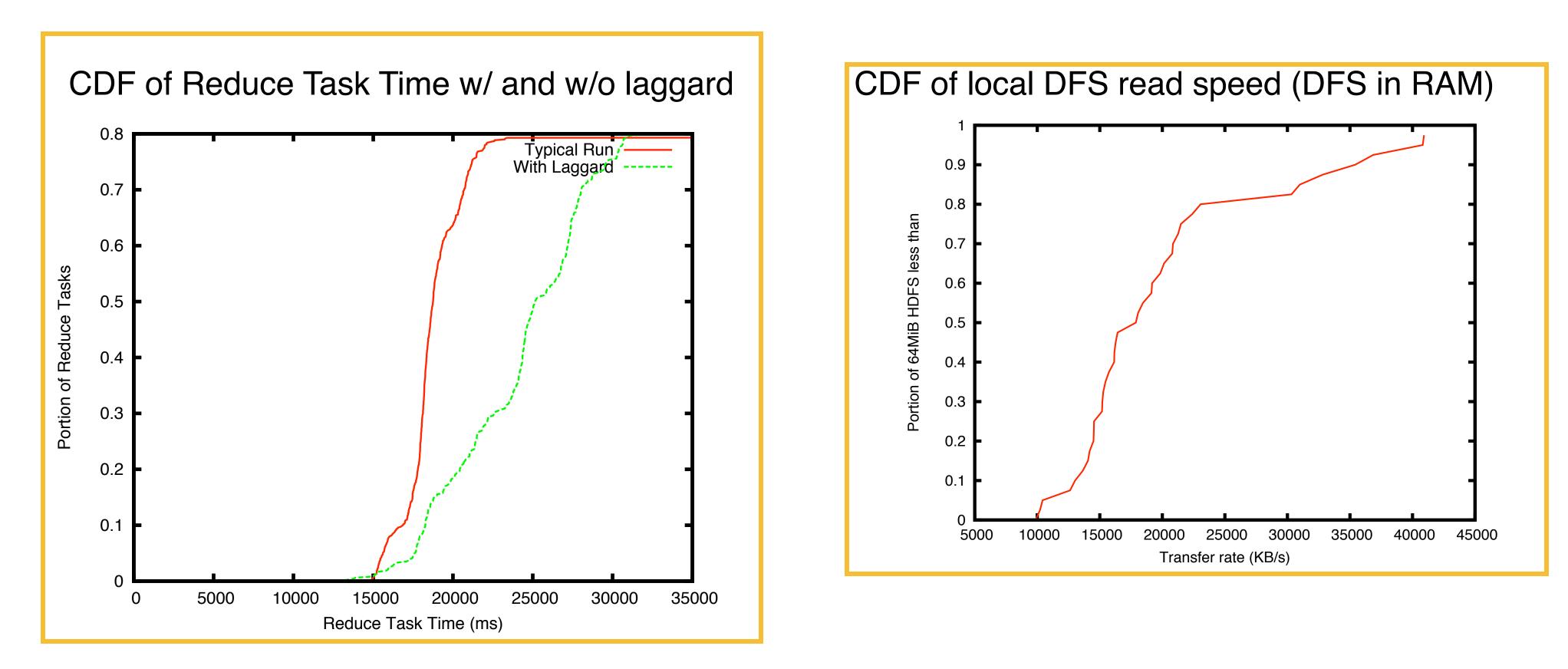
SHUFFLE STEP LAGGARDS

A single slow machine (observed "naturally" on EC2) does not send its map outputs quickly:



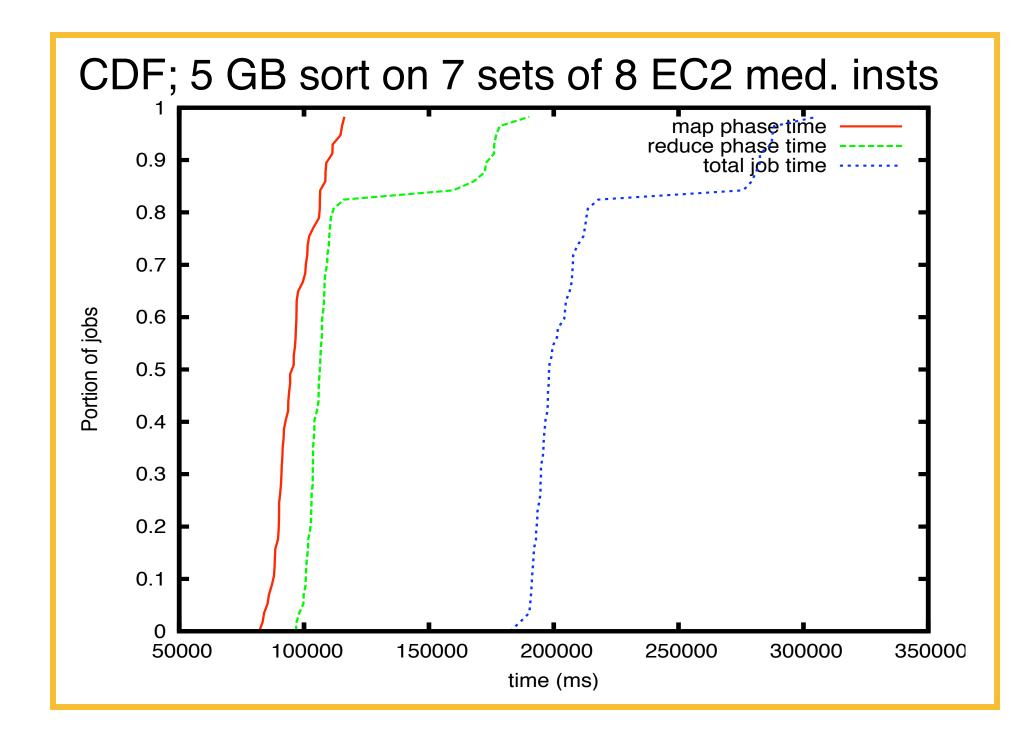
With a trivial mapper, reading the data from the distributed filesystem (DFS) dominated the map phase. This transfer is long enough to make setup overhead negligible but even local reads from a RAM disk-backed DFS are not consistent when the FS is loaded by the job:

This makes *all* reduce tasks in the job slow:



OVERALL JOB BREAKDOWN

Overall synthetic job time is relatively consistent (<10% change in completion time) on a particular set of nodes, except for the occasional set of runs with laggard nodes:

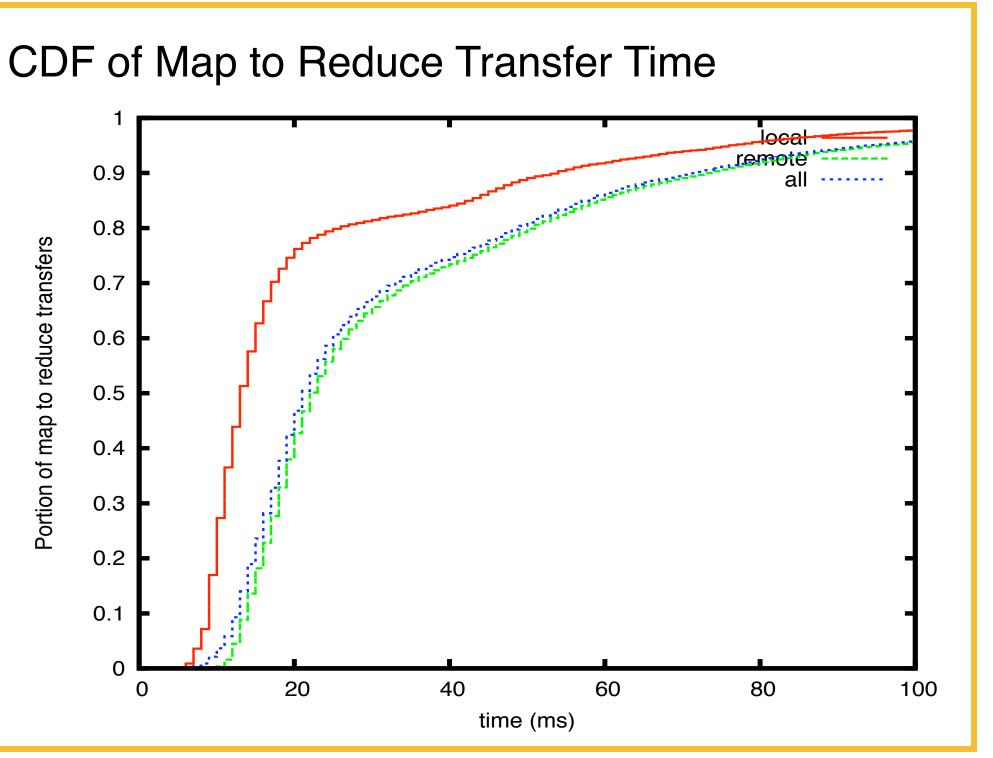


MAP READS

PROPOSED WORKAROUNDS •Make scheduler aware of time that will be necessary to pull map outputs when reducers are running •Make scheduler account aware that task slots on the same machine aren't

SHUFFLE STEP CONNECITON OVERHEAD

Hadoop's reducers "pull" from each mapper. Even after accounting for locality and performing only one transfer at a time (instead of 8), the time taken to do this is inconsistent (for approx 900KB: median ~50ms, but 90th percentile ~300ms), despite consistent, small data sizes:



independent.

•Before running reduce tasks:

•Consolidate map outputs so they can be transferred al at once;

•Move map outputs to where reducers are expected to run

•Separate scheduling of data transfer part of tasks and data processing part of tasks

OTHER FUTURE WORK

Empirical data:

-Analyze Yahoo! Traces to estimate real effects of these overheads, job configuration mistakes