

CS61B Lecture #19

Administrative:

Today:

- Trees

Readings for Today: *Data Structures*, Chapter 5

Readings for Next Topic: *Data Structures*, Chapter 6

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A Recursive Structure

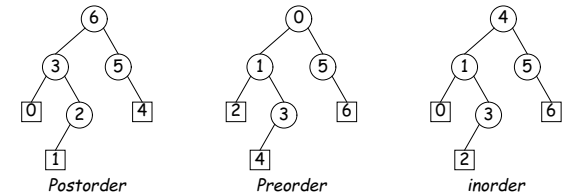
- Trees naturally represent recursively defined objects with more than one recursive instance.
- Common examples: expressions, sentences.
 - Expressions have definitions such as "an expression consists of a literal or two expressions separated by an operator."
- Also describe structures in which we recursively divide a set into multiple subsets.

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Fundamental Operation: Traversal

- *Traversing a tree* means enumerating (some subset of) its nodes.
- Typically done recursively, because that is natural description.
- As nodes are enumerated, we say they are *visited*.
- Three basic orders for enumeration (+ variations):
 - **Preorder**: visit node, traverse its children.
 - **Postorder**: traverse children, visit node.
 - **Inorder**: traverse first child, visit node, traverse second child (binary trees only).

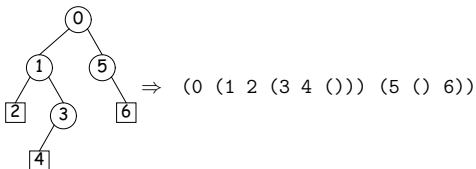


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... and Traversal Code (Preorder)

Problem: Convert



```
void printAsLisp (Tree T, PrintWriter p) {
    if (T == null)
        p.print (" (");
    else if (T.degree () == 0)
        p.print (" " + T.label ());
    else {
        p.print "(" + T.label ();
        for (int i = 0; i < T.maxChild (); i += 1)
            printAsLisp (T.child (i), p);
        p.print (")");
    }
}
```

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Preorder Traversal Code: General

```
void preorderTraverse (Tree T, Action whatToDo)
{
    if (T != null) {
        whatToDo.action (T);
        for (int i = 0; i < T.maxChild (); i += 1)
            preorderTraverse (T.child (i), whatToDo);
    }
}
```

- What is Action?

```
interface Action {
    void action (Tree T);
}

class Print implements Action {
    void action (Tree T) {
        System.out.print (T.label ());
    }
}
```

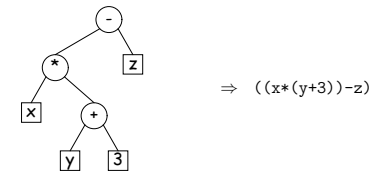
```
// preorderTraverse (myTree, new Print ());
```

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Inorder Traversal \iff Infix Expressions

Problem: Convert



```
// CORRECTED VERSION
void printExpr (Tree T, PrintWriter p) {
    if (T.degree () == 0)
        p.print (T.label ());
    else {
        p.print "(";
        printExpr (T.left (), p);
        p.print (T.label ());
        printExpr (T.right (), p);
        p.print (")");
    }
}
```

- To think about: how to get rid of all those parentheses?

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Times

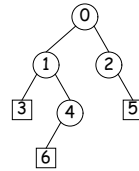
- The traversal algorithms have roughly the form of the boom example in §1.3.3 of *Data Structures*—an exponential algorithm.
- However, the role of M in that algorithm is played by the *height* of the tree, not the number of nodes.
- In fact, easy to see that tree traversal is *linear*: $\Theta(N)$, where N is the # of nodes.
- Reason: Form of the algorithm implies that there is one visit at the root, and then one visit for every edge in the tree.
- Since every node but the root has exactly one parent, and the root has none, must be $N - 1$ edges in any non-empty tree.
- In positional tree, is also one recursive call for each empty tree, but # of empty trees can be no greater than kN , where k is arity of tree.
- For k -ary tree (max # children is k), $h+1 \leq N \leq \frac{k^{h+1}-1}{k-1}$, where h is height.
- So $h \in \Omega(\log_k N) = \Omega(\lg N)$ and $h \in O(N)$.
- Many tree algorithms look at one child only. For them, time is proportional to the *height* of the tree, and this is $\Theta(\lg N)$, assuming that tree is *bushy*—each level has about as many nodes as possible.

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Level-Order (Breadth-First) Traversal

Problem: Traverse all nodes at depth 0, then depth 1, etc:



- One technique: *Iterative Deepening*. For each level, k , from 0 to h , call `doLevel(T,k)`

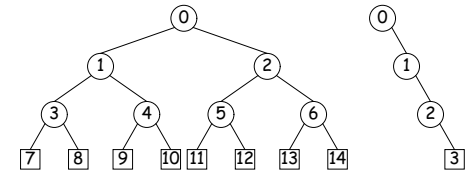
```

void doLevel (Tree T, int lev) {
    if (lev == 0)
        visit T
    else
        for each non-null child, C, of T {
            doLevel (C, lev-1);
        }
}
  
```

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Iterative Deepening Time?



- Let h be height, N be # of nodes.
- Count # edges traversed (i.e. # of calls, not counting null nodes).
- First (full) tree: 1 for level 0, 3 for level 1, 7 for level 2, 15 for level 3.
- Or in general $(2^1-1) + (2^2-1) + \dots + (2^{h+1}-1) = 2^{h+2} - h \in \Theta(N)$, since $N = 2^{h+1} - 1$ for this tree.
- Second (right leaning) tree: 1 for level 0, 2 for level 1, 3 for level 2.
- Or in general $(h+1)(h+2)/2 = N(N+1)/2 \in \Theta(N^2)$, since $N = h+1$ for this kind of tree.

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Iterative Traversals

- Tree recursion conceals data: a *stack* of nodes (all the T arguments) and a little extra information.
- Can make the data explicit, e.g.:

```

void preorderTraverse2 (Tree T, Action whatToDo) {
    Stack s = new Stack ();
    s.push (T);
    while (! s.isEmpty ()) {
        Tree node = (Tree) s.pop ();
        if (node == null)
            continue;
        whatToDo.action (node);
        for (int i = node.maxChild ()-1; i >= 0; i --)
            s.push (node.child (i));
    }
}
  
```

- To do a breadth-first traversal, use a queue instead of a stack, replace push with add, and pop with removeFirst.
- Makes breadth-first traversal worst-case linear time in all cases, but also linear space for "bushy" trees.

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Iterators for Trees

- Frankly, iterators are not terribly convenient on trees.
- But can use ideas from iterative methods.

```

class PreorderTreeIterator implements Iterator {
    private Stack s = new Stack ();

    public PreorderTreeIterator (Tree T) {
        s.push (T);
    }

    public boolean hasNext () {
        return ! s.isEmpty ();
    }

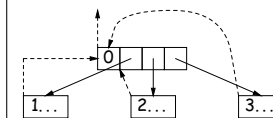
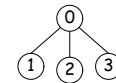
    public Object next () {
        Tree result = (Tree) s.pop ();
        for (int i = result.maxChild ()-1; i >= 0; i --)
            s.push (result.child (i));
        return result.label ();
    }

    void remove () { // I'm lazy.
        throw new UnsupportedOperationException ();
    }
}
  
```

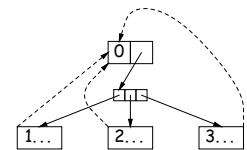
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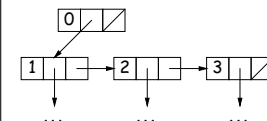
Representation Choices



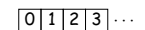
(a) Embedded child pointers
(+ optional parent pointers)



(b) Array of child pointers
(+ optional parent pointers)



(c) child/sibling pointers



(d) pre-order array
(complete trees)

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