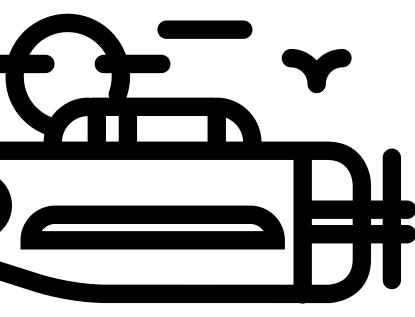
In Search of an Understandable Consensus Algorithm



What is Raft

Another protocol for building replicated state machines.

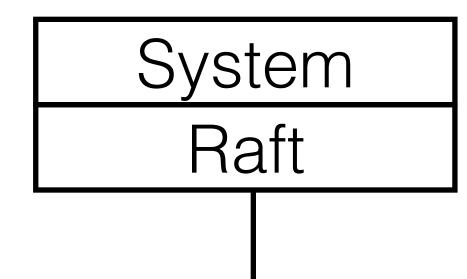
What is Raft

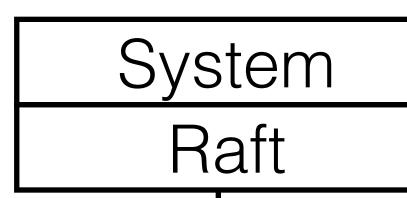
- Another protocol for building replicated state machines.
- Different level of abstraction compared to Paxos made Simple
 - Mainly describes a protocol to keep logs consistent.
 - Describes something closer to view-stamped replication than Paxos.

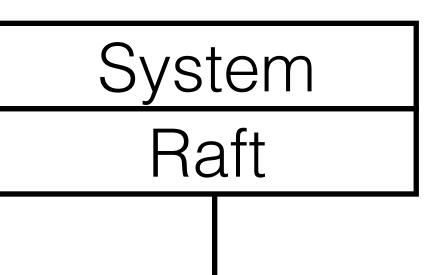
What is Raft

- Another protocol for building replicated state machines.
- Different level of abstraction compared to Paxos made Simple
 - Mainly describes a protocol to keep logs consistent.
 - Describes something closer to view-stamped replication than Paxos.
- Claim: Easier to understand (really?)
 - Will return to this later.

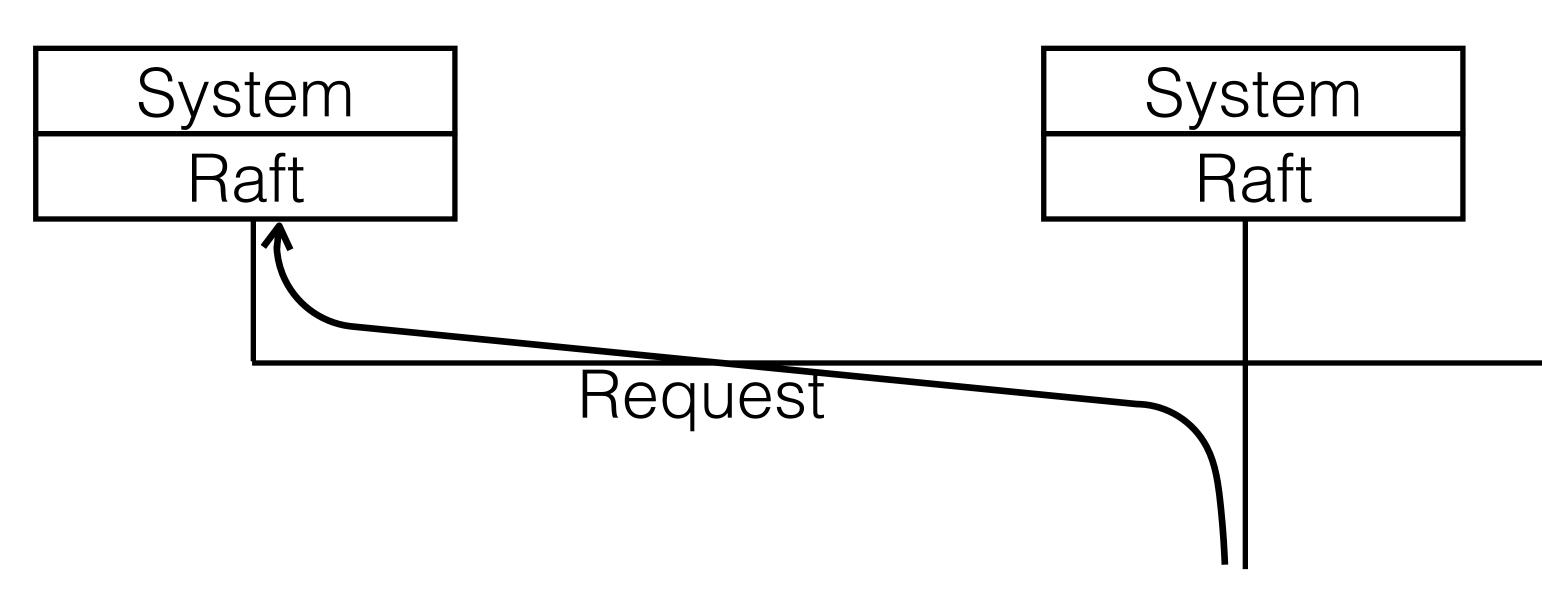
System Model

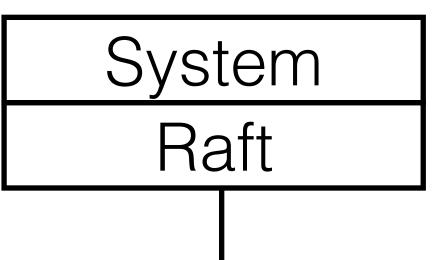


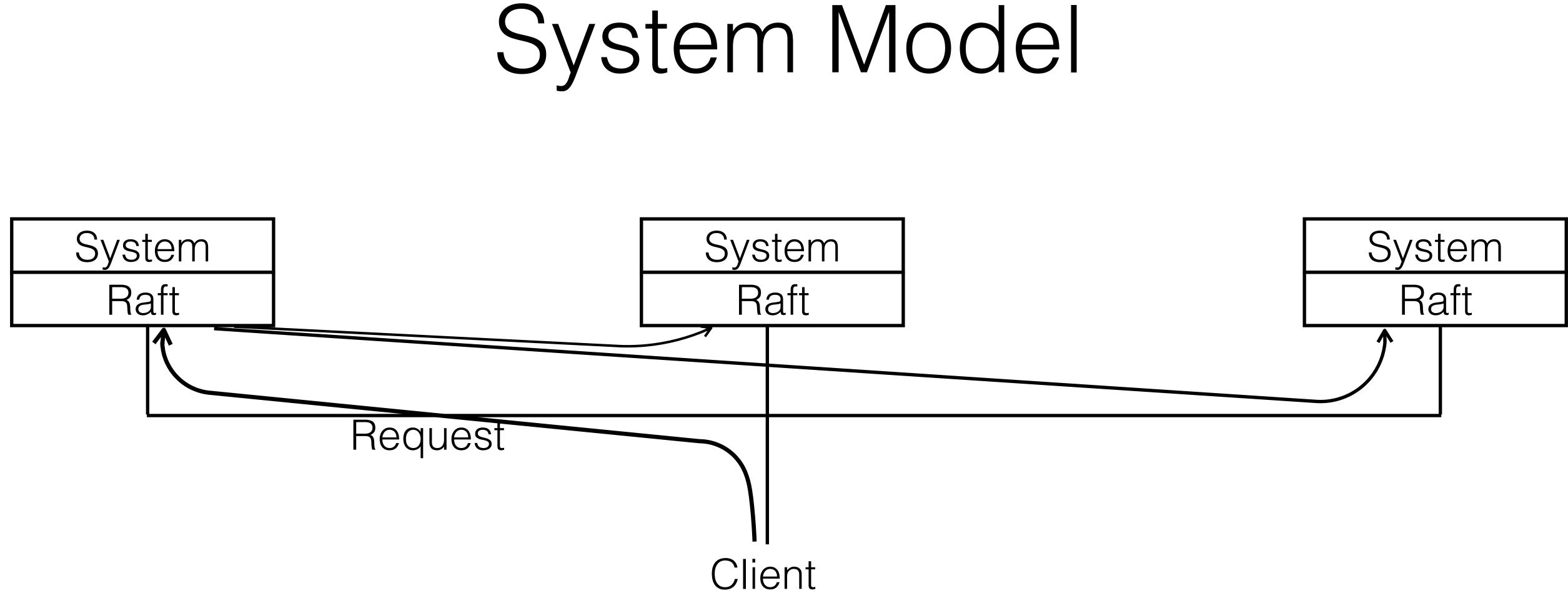


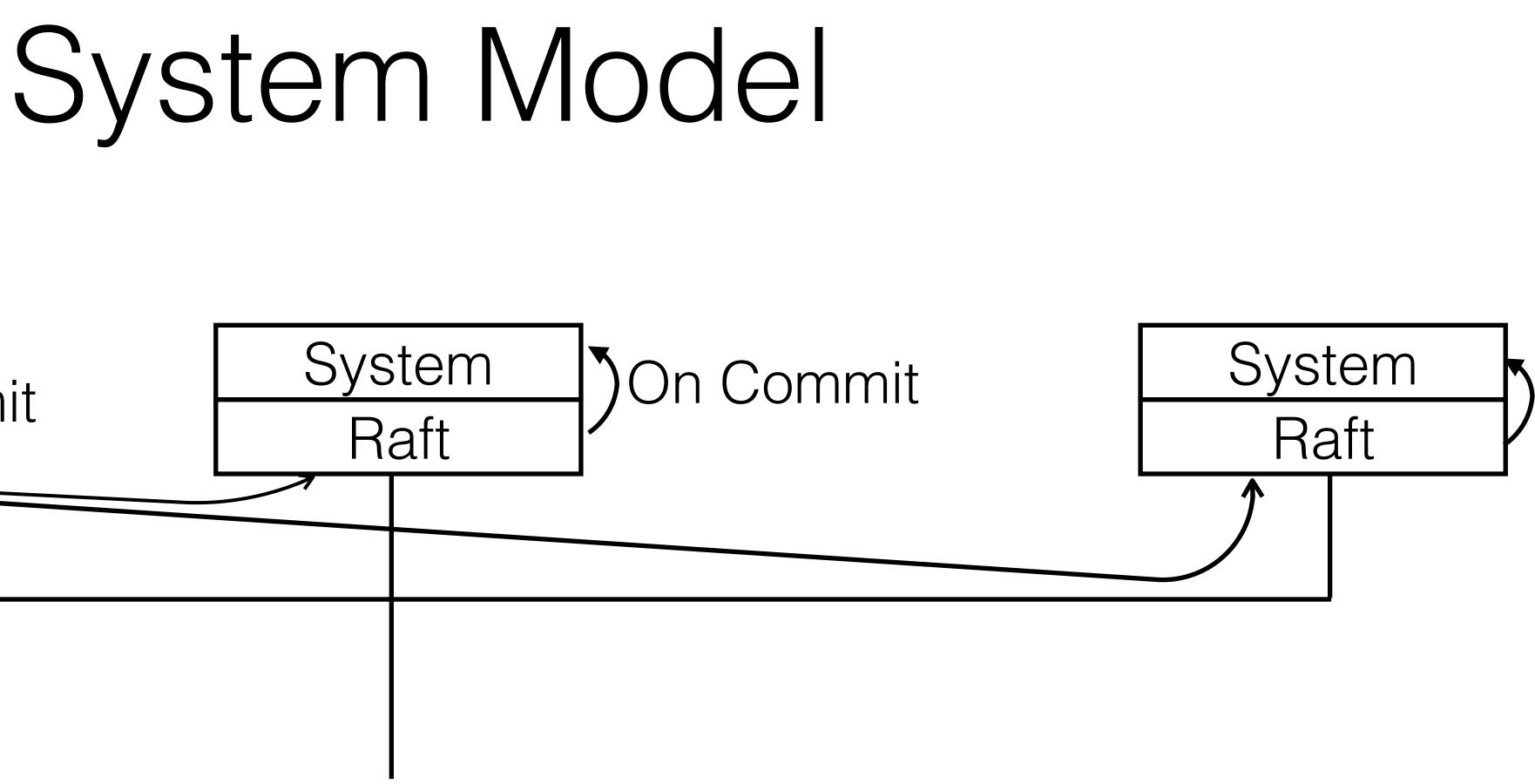


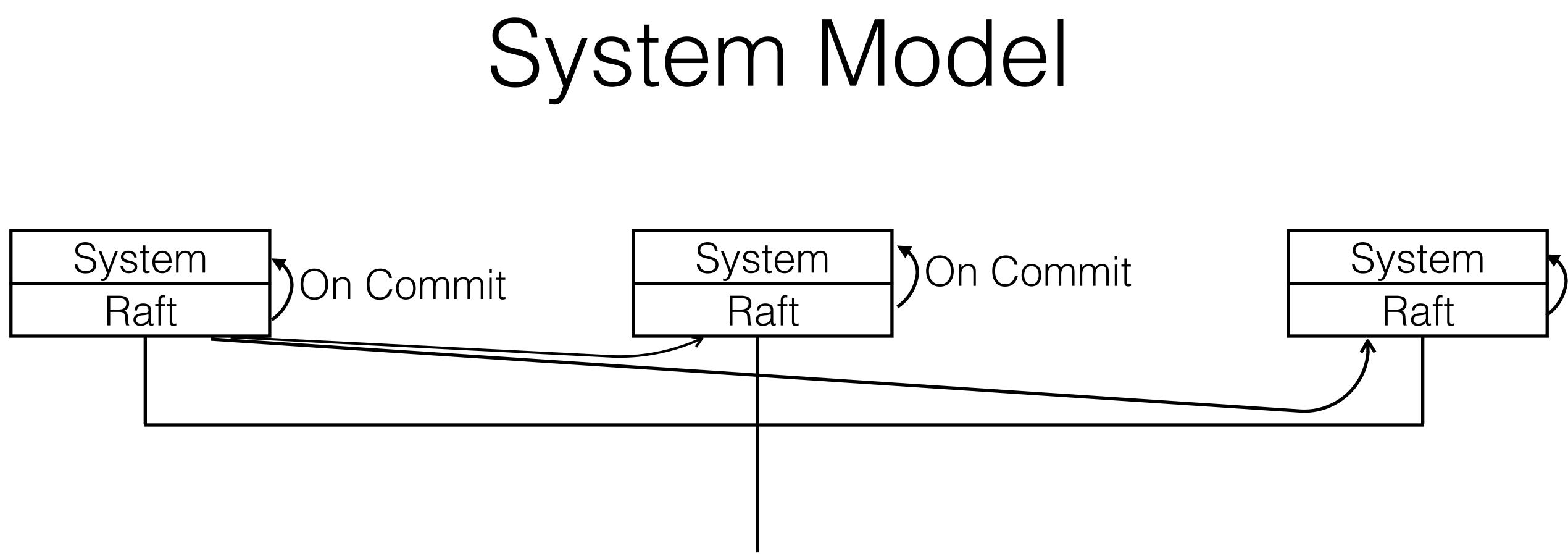


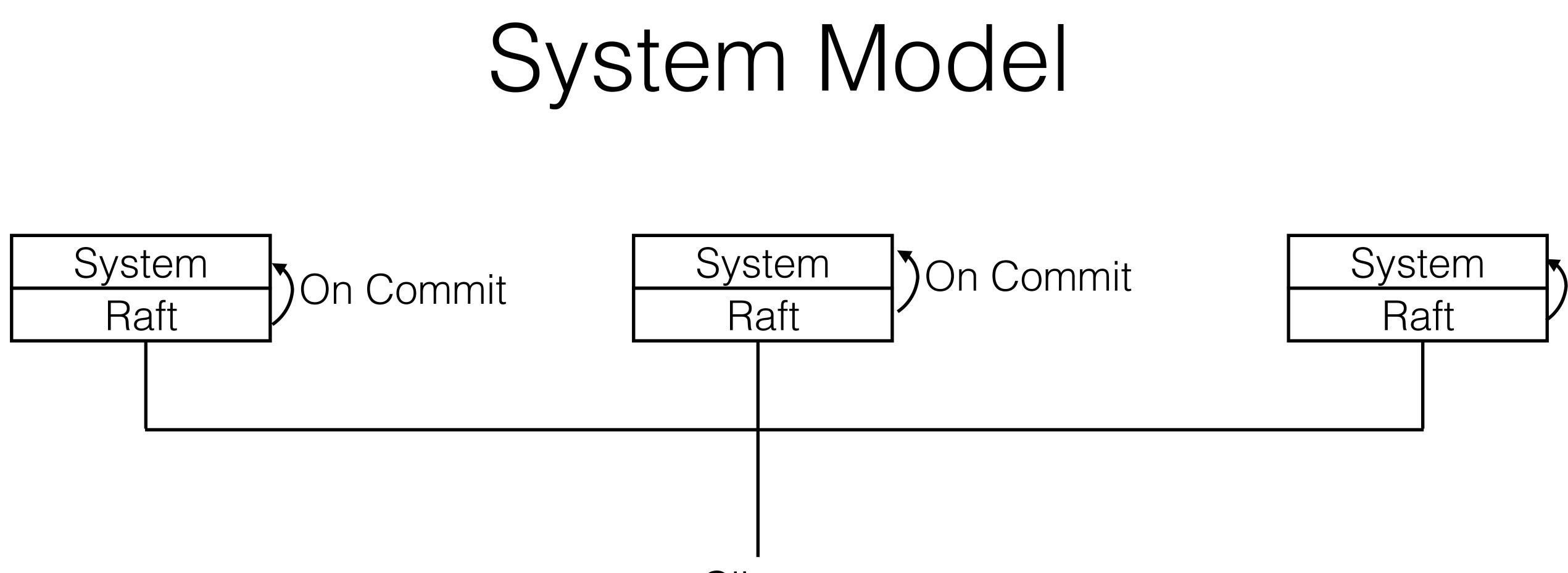












Safety Requirements

- •
- State Machine Safety: all server's logs agree on committed entries.

Log Matching: if the kth log entry for 2 logs match, so do all previous entries.



Leader Election + Log Replication



Leader Election

Same limitations as consensus

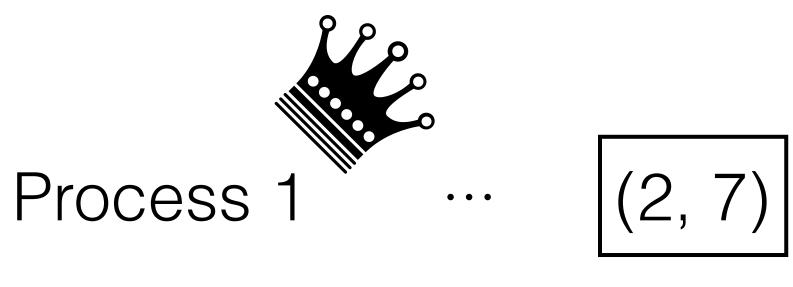
Leader Election + Log Replication



Same limitations as consensus

Leader Election + Log Replication



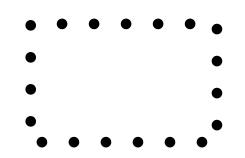


Process 2 \cdots (2, 7):

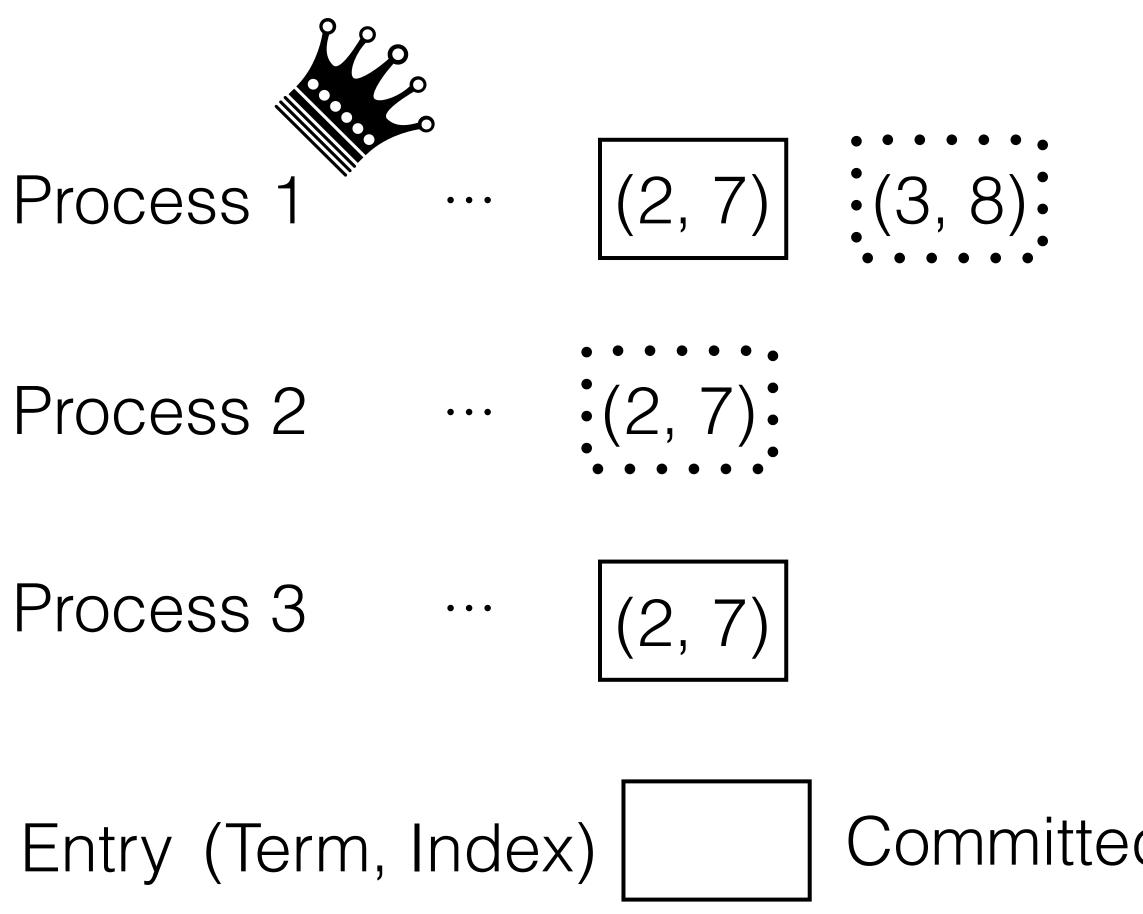
Process 3 …

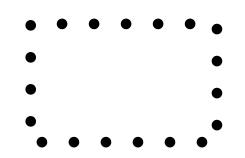
(2, 7)

Entry (Term, Index)

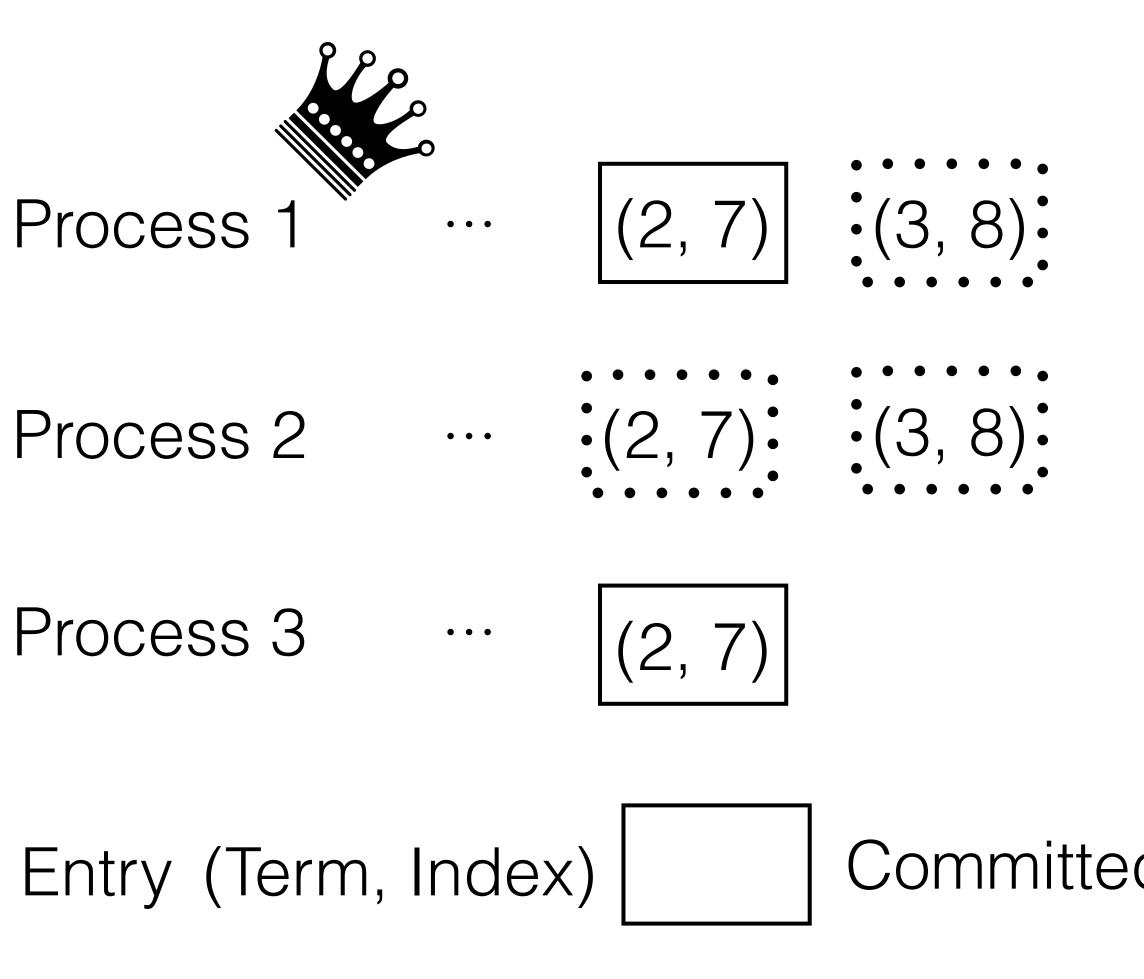


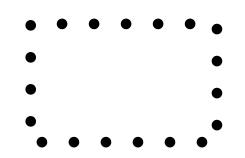




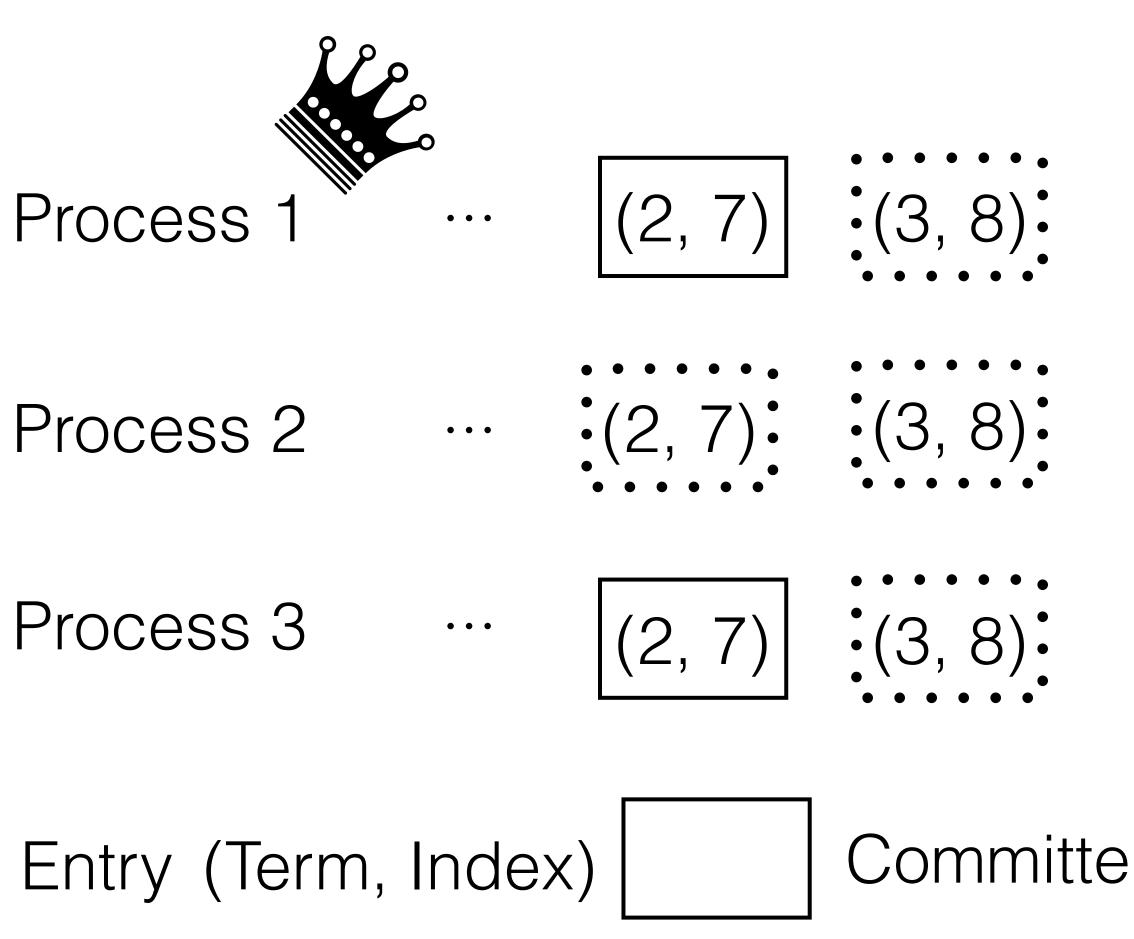


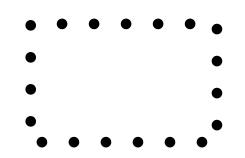




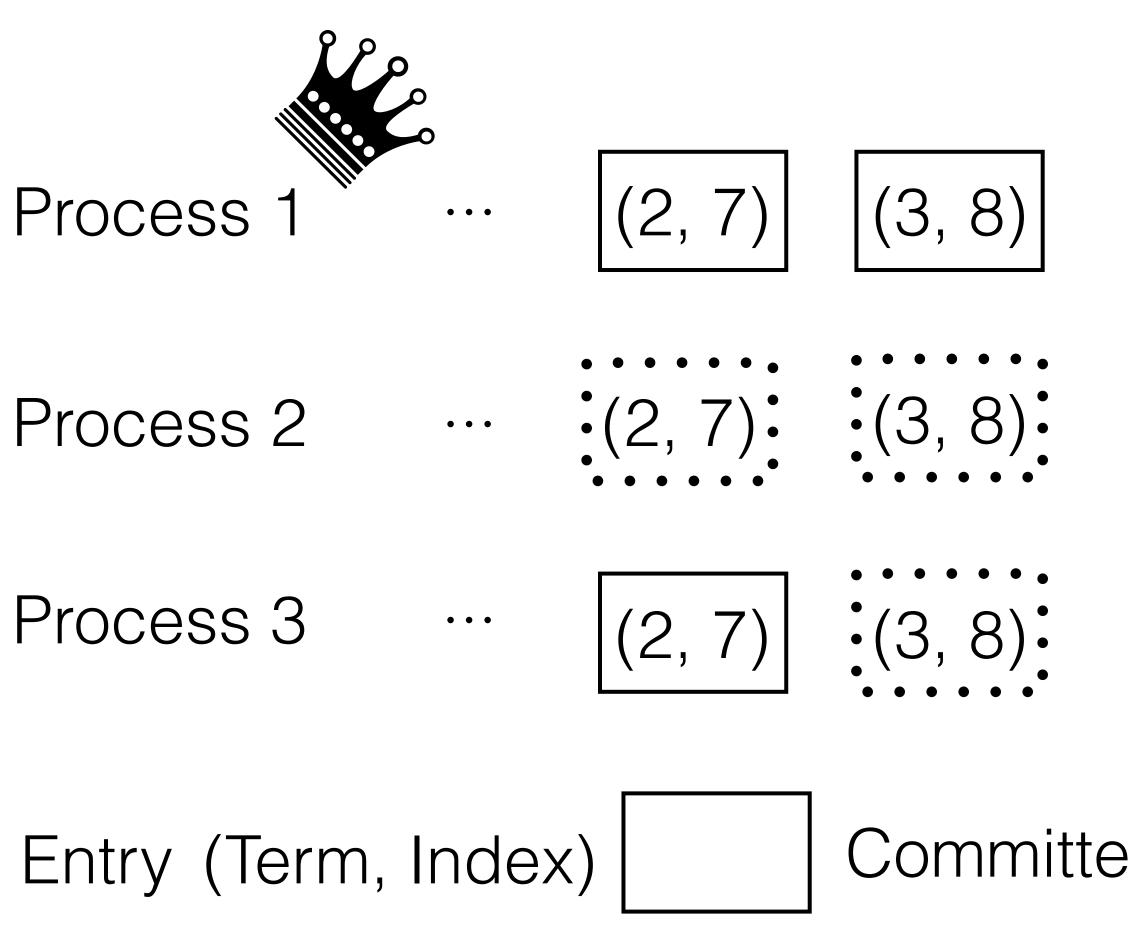


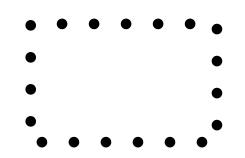




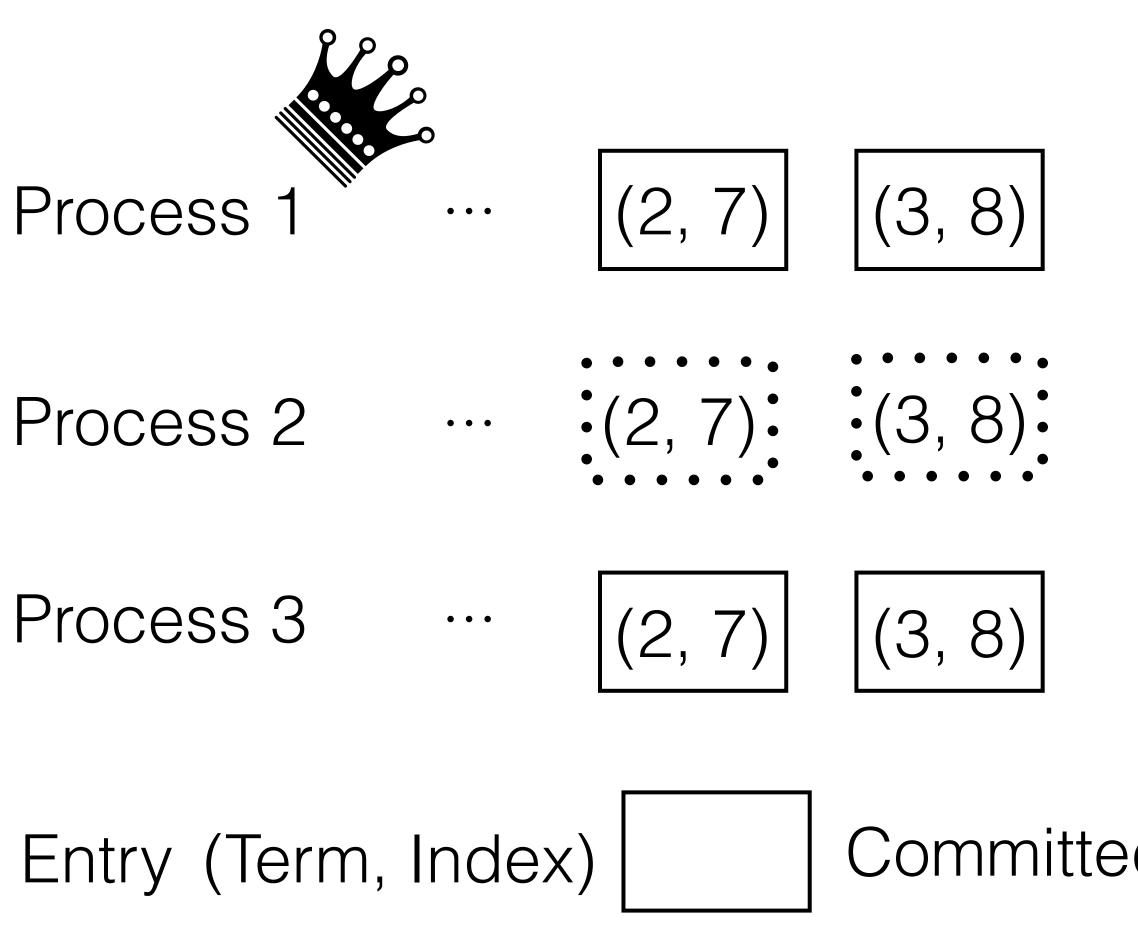


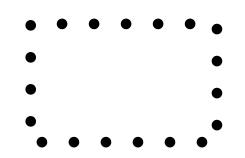




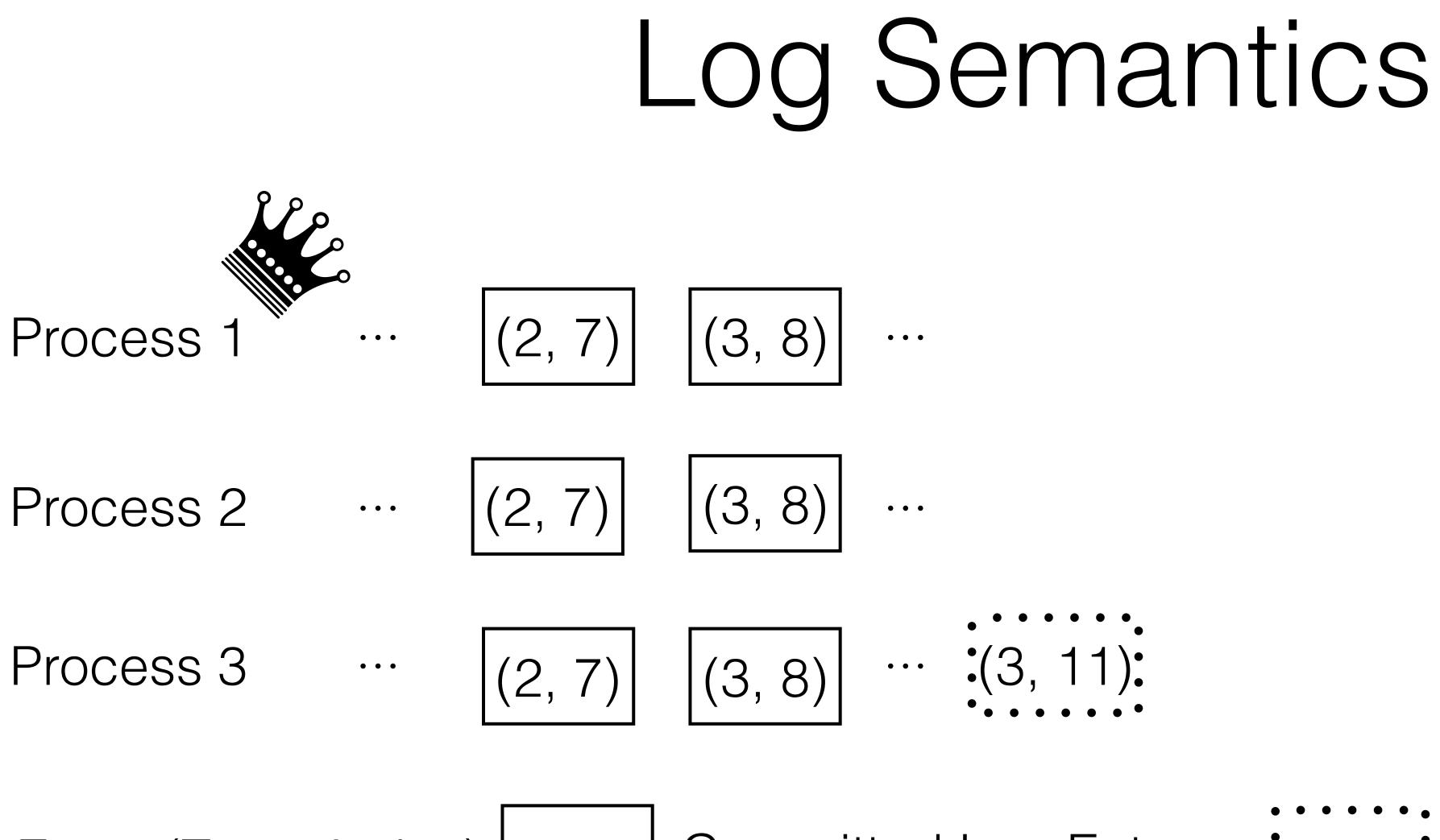






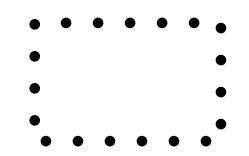






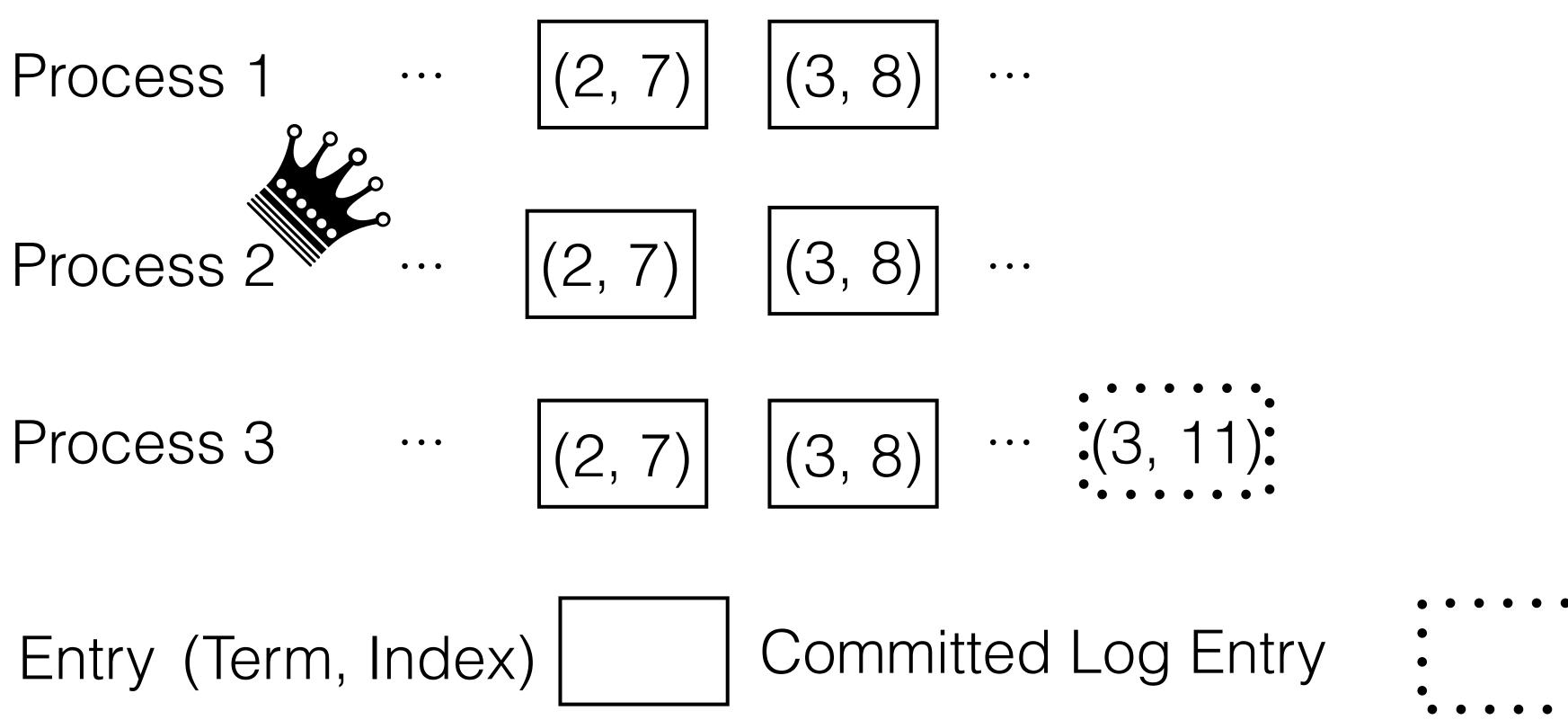
Entry (Term, Index)

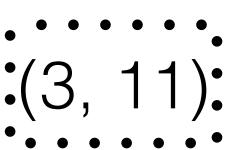
Committed Log Entry

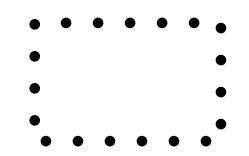


: Uncommitted Log Entry



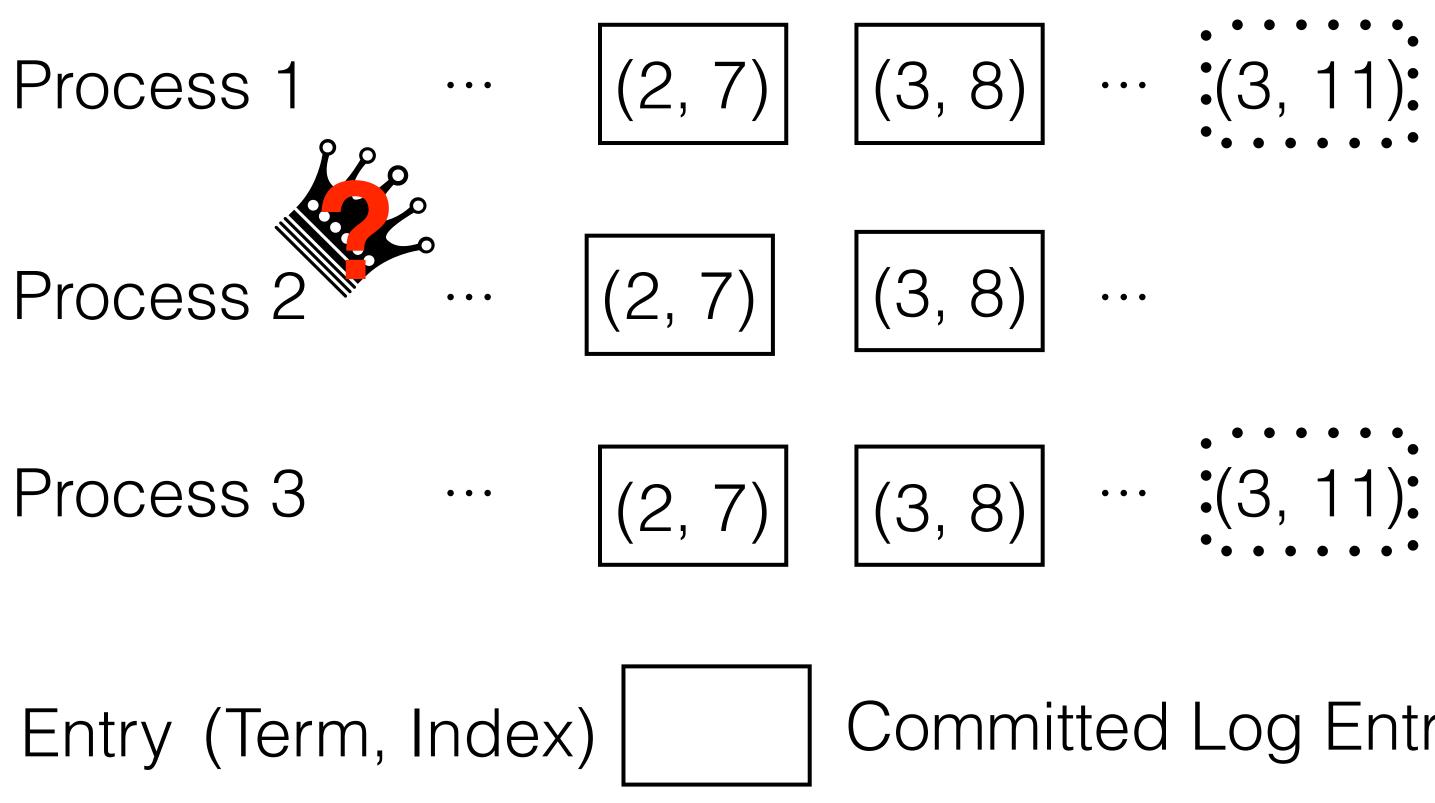


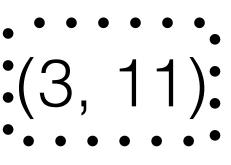


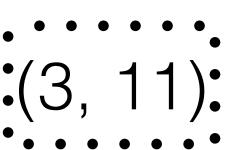


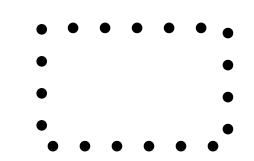
Uncommitted Log Entry



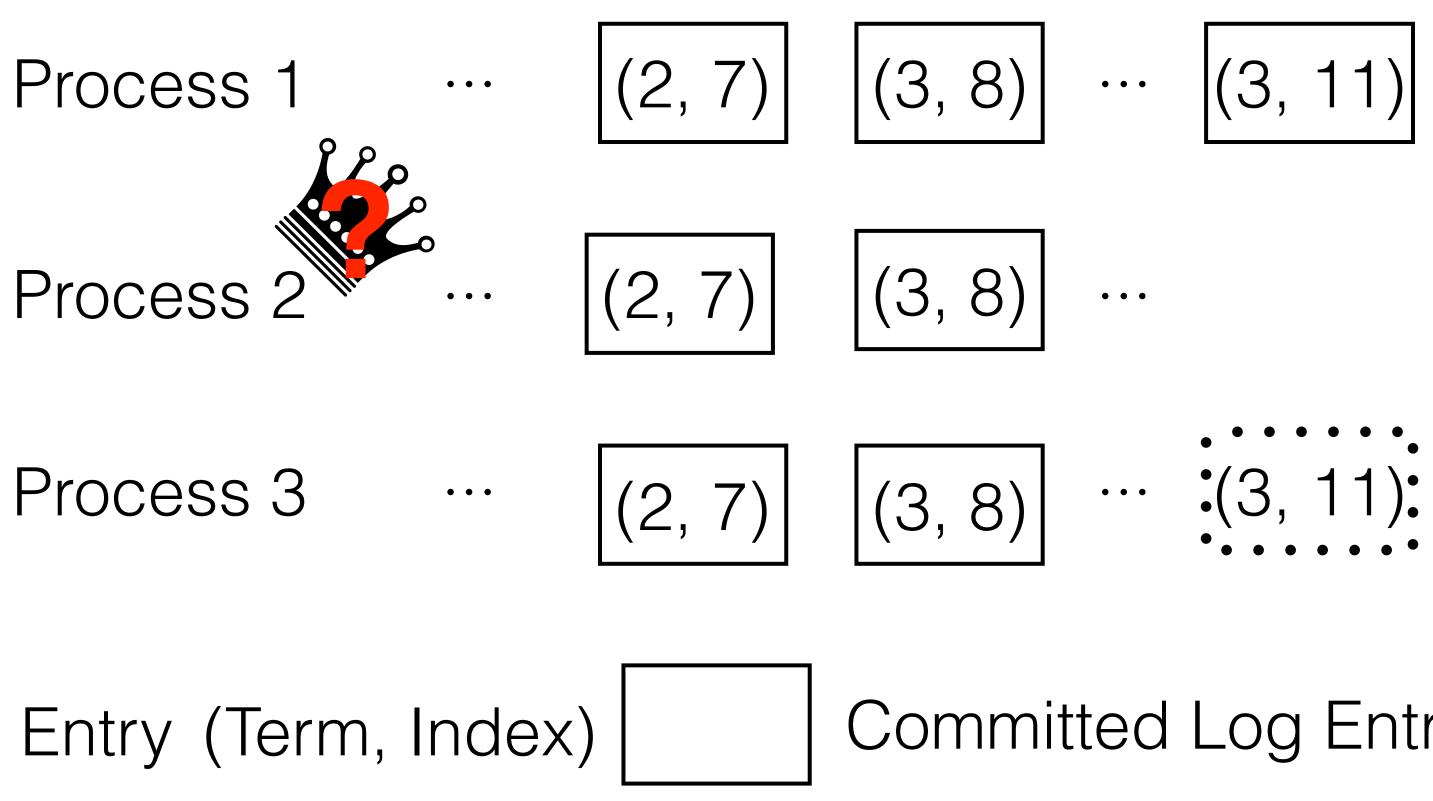


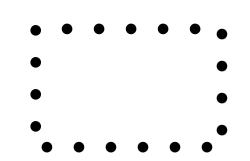




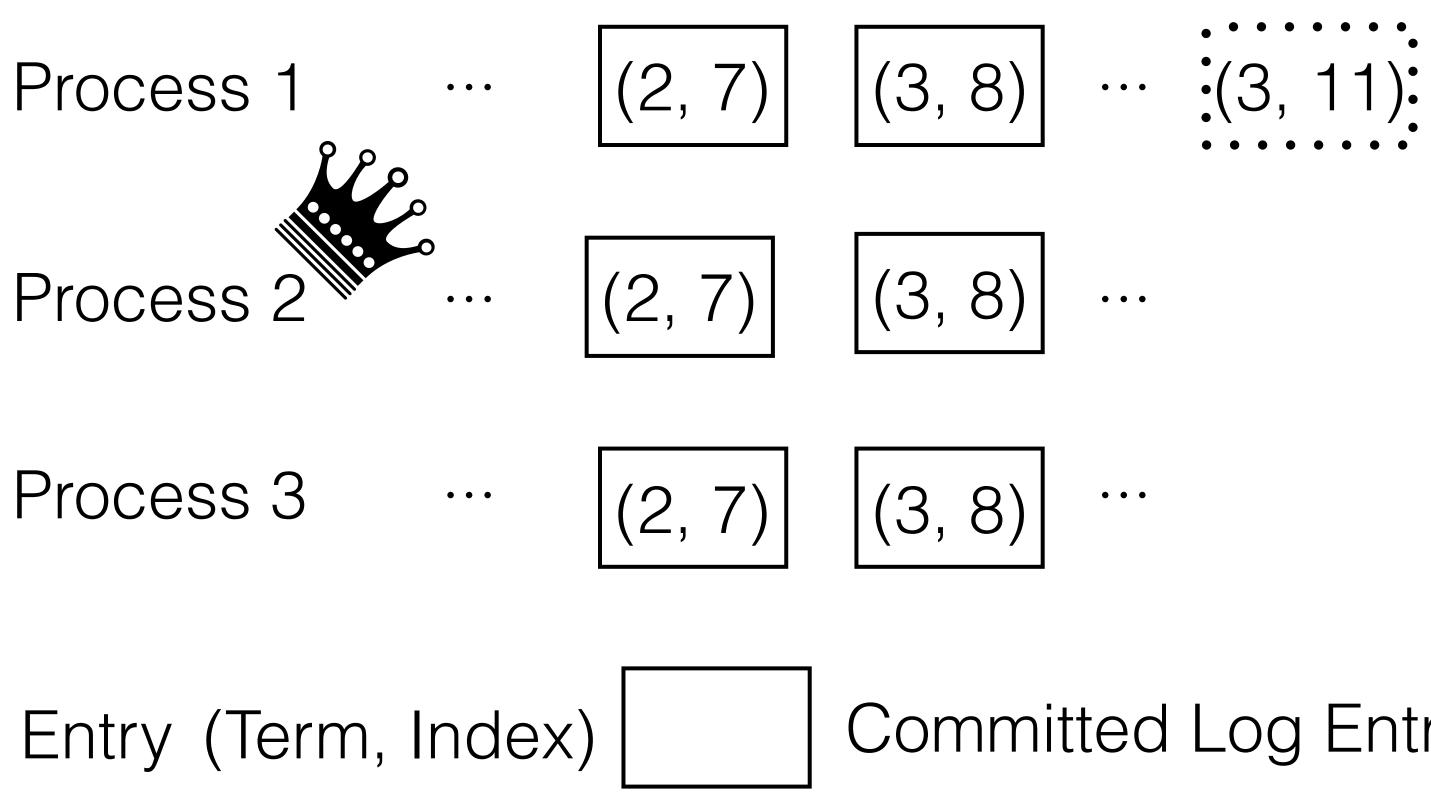


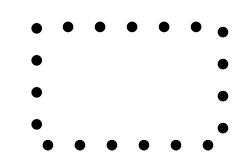




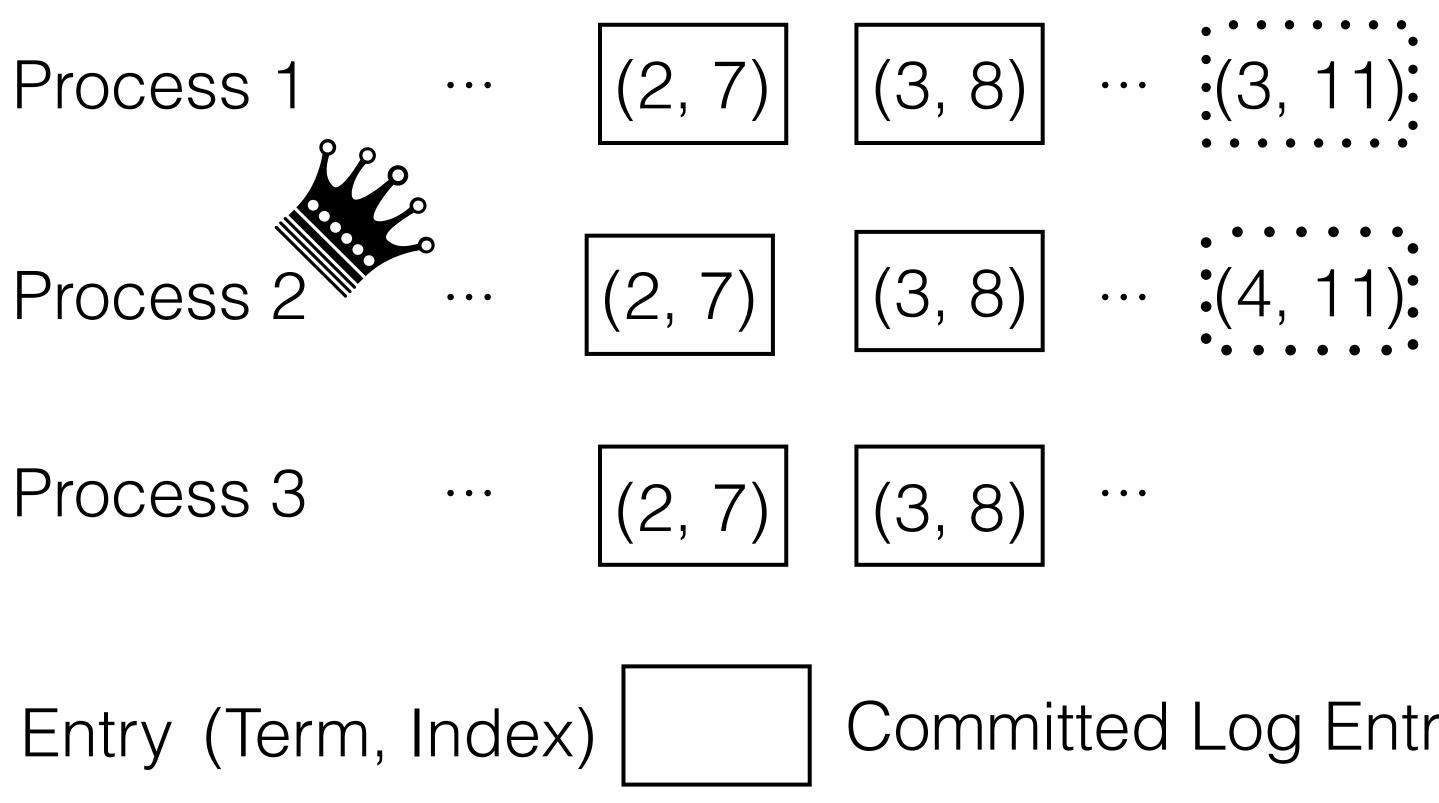


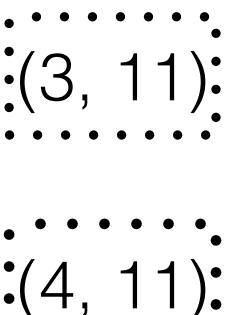


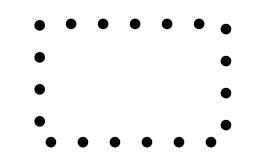




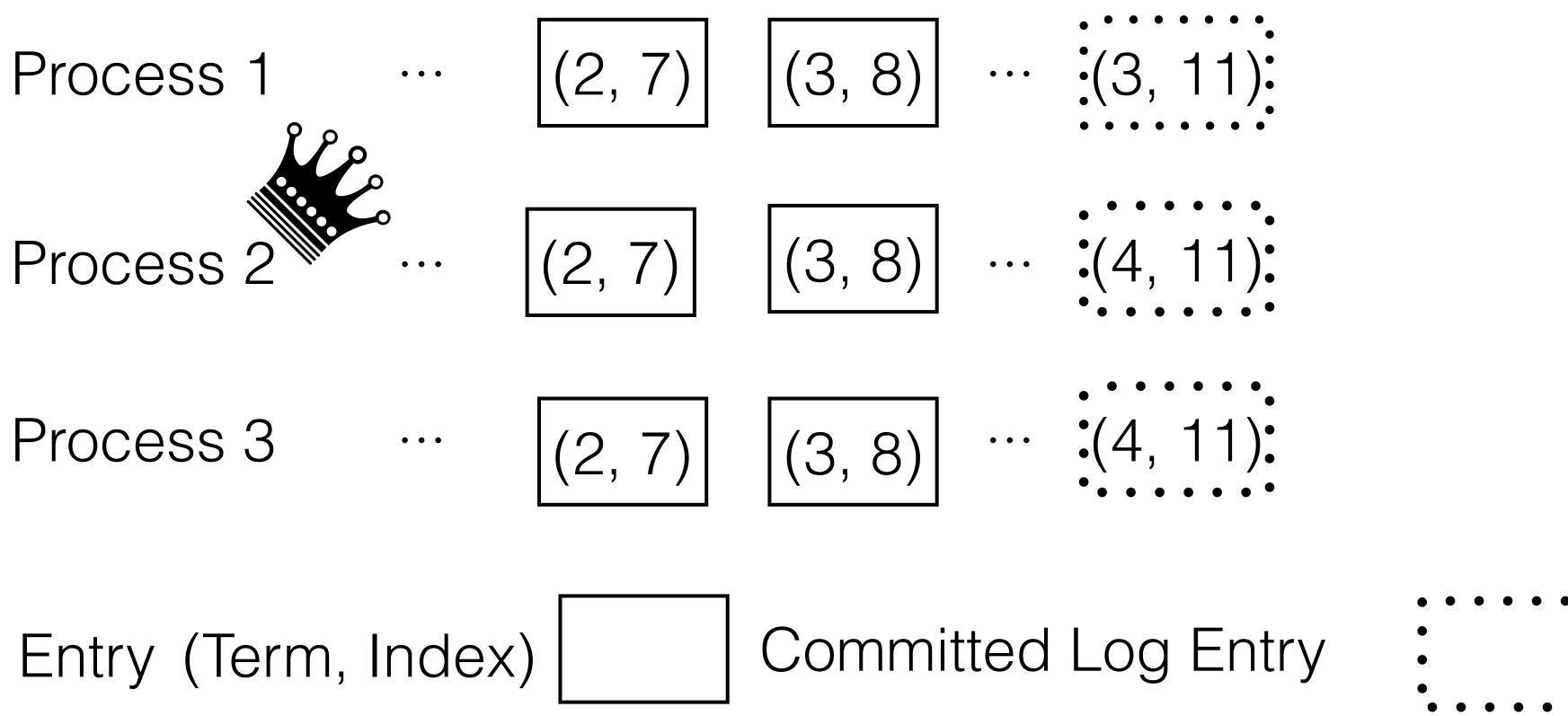


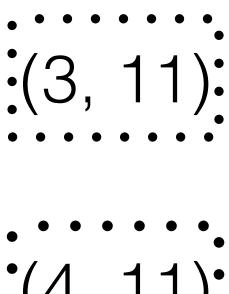


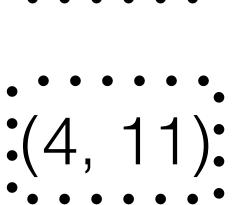


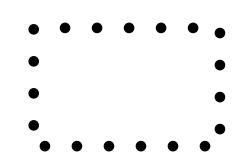






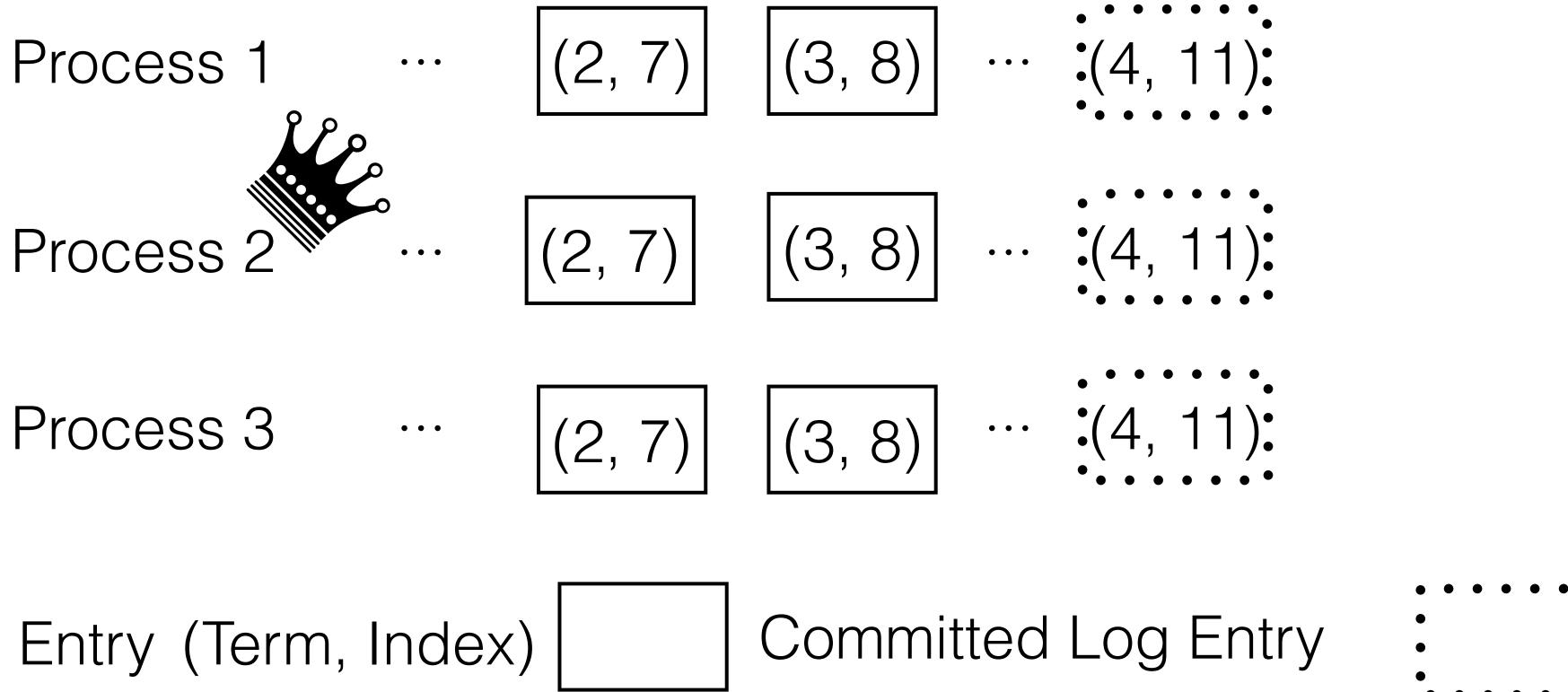


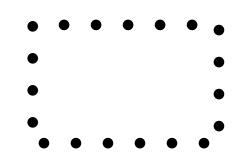




Uncommitted Log Entry







Uncommitted Log Entry



Leader Election

- Each process has a random election timer.
 - On timeout declares itself a candidate and requests votes.
- Any follower votes for a candidate if
 - Requested term > current term
 - Candidate's log is at least as up to date as follower's.

More Understandable?

- What is the point of understandability?
- Easier to show correctness?
 - What is the state of the log at any step?
 - What entries are uncommitted?
 - What entries will survive?
 - Verdi spent a year on proving correctness, uncertain results.

More Understandable?

- Easier to use?
 - Is it that much easier than ZooKeeper, Chubby, etc.?

More Understandable?

- Easier to implement correctly?
 - Lots of implementations, a fair number of bugs.
 - See recent work by Colin Scott and me.

Some More Thoughts

- Raft is more "directly" usable as described in the paper.
 - Higher-level abstraction (RSM), as opposed to an algorithm
 - However, algorithm is (perhaps) easier to fit into different settings.
- Helped by when it was released
 - Distributed systems were hot, "practitioners" weren't looking at old papers.

