

Statistical NLP Spring 2007



Lecture 15: Parsing II

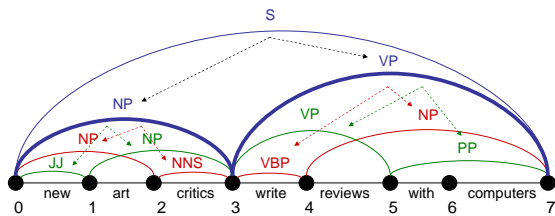
Dan Klein – UC Berkeley

A Recursive Parser

```
bestScore(X,i,j,s)
  if (j = i+1)
    return tagScore(X,s[i])
  else
    return max score(X->YZ) *
              bestScore(Y,i,k) *
              bestScore(Z,k,j)
```

- Will this parser work?
- Why or why not?
- Memory requirements?

An Example



A Memoized Parser

- One small change:

```
bestScore(X,i,j,s)
  if (scores[X][i][j] == null)
    if (j = i+1)
      score = tagScore(X,s[i])
    else
      score = max score(X->YZ) *
              bestScore(Y,i,k) *
              bestScore(Z,k,j)
    scores[X][i][j] = score
  return scores[X][i][j]
```

Memory: Theory

- How much memory does this require?
 - Have to store the score cache
 - Cache size: $|\text{symbols}| * n^2$ doubles
 - For the plain treebank grammar:
 - $X \sim 20K, n = 40, \text{double} \sim 8 \text{ bytes} = \sim 256MB$
 - Big, but workable.
- What about sparsity?

Time: Theory

- How much time will it take to parse?
 - Have to fill each cache element (at worst)
 - Each time the cache fails, we have to:
 - Iterate over each rule $X \rightarrow YZ$ and split point k
 - Do constant work for the recursive calls
 - Total time: $|\text{rules}| * n^3$
 - Cubic time
 - Something like 5 sec for an unoptimized parse of a 20 word sentences

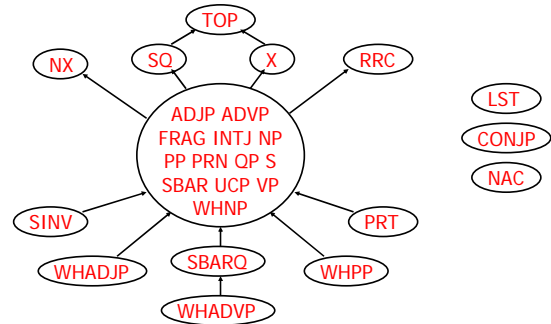
Unary Rules

- Unary rules?

```

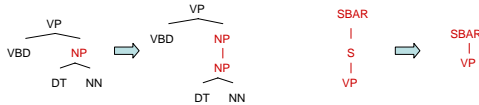
bestScore(X,i,j,s)
  if (j = i+1)
    return tagScore(X,s[i])
  else
    return max max score(X->YZ) *
                bestScore(Y,i,k) *
                bestScore(Z,k,j)
                max score(X->Y) *
                bestScore(Y,i,j)
    
```

Same-Span Reachability



CNF + Unary Closure

- We need unaries to be non- ϵ
 - Can address by pre-calculating the *unary closure*
 - Rather than having zero or more unaries, always have exactly one



- Alternate unary and binary layers
- Reconstruct unary chains afterwards

Alternating Layers

```

bestScoreB(X,i,j,s)
  return max max score(X->YZ) *
              bestScoreU(Y,i,k) *
              bestScoreU(Z,k,j)
    
```

```

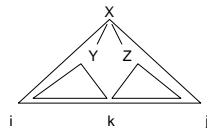
bestScoreU(X,i,j,s)
  if (j = i+1)
    return tagScore(X,s[i])
  else
    return max max score(X->Y) *
                bestScoreB(Y,i,j)
    
```

A Bottom-Up Parser (CKY)

- Can also organize things bottom-up

```

bestScore(s)
  for (i : [0,n-1])
    for (X : tags[s[i]])
      score[X][i][i+1] =
        tagScore(X,s[i])
    for (diff : [2,n])
      for (i : [0,n-diff])
        j = i + diff
        for (X->YZ : rule)
          for (k : [i+1, j-1])
            score[X][i][j] = max score[X][i][k],
                                score[X][k][j],
                                score(X->YZ) *
                                score[Y][i][k] *
                                score[Z][k][j]
    
```



Efficient CKY

- Lots of tricks to make CKY efficient
 - Most of them are little engineering details:
 - E.g., first choose k, then enumerate through the Y:[i,k] which are non-zero, then loop through rules by left child.
 - Optimal layout of the dynamic program depends on grammar, input, even system details.
 - Another kind is more critical:
 - Many X:[i,j] can be suppressed on the basis of the input string
 - We'll see this next class as figures-of-merit or A* heuristics

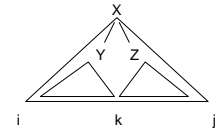
Memory: Practice

- Memory:
 - Still requires memory to hold the score table
- Pruning:
 - score[X][i][j] can get too large (when?)
 - can instead keep beams scores[i][j] which only record scores for the top K symbols found to date for the span [i,j]

Time: Theory

- How much time will it take to parse?

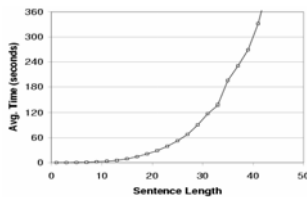
- For each diff ($\leq n$)
 - For each i ($\leq n$)
 - For each rule $X \rightarrow YZ$
 - For each split point k
 - Do constant work



- Total time: $|\text{rules}| * n^3$

Runtime: Practice

- Parsing with the vanilla treebank grammar:



~ 20K Rules
(not an optimized parser!)
Observed exponent:
3.6

- Why's it worse in practice?
 - Longer sentences "unlock" more of the grammar
 - All kinds of systems issues don't scale

Rule State Reachability

Example: NP CC •



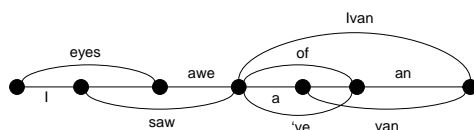
Example: NP CC NP •



- Many states are more likely to match larger spans!

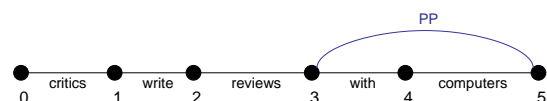
(Speech) Lattices

- There was nothing magical about words spanning exactly one position.
- When working with speech, we generally don't know how many words there are, or where they break.
- We can represent the possibilities as a lattice and parse these just as easily.



A Simple Chart Parser

- Chart parsers are sparse dynamic programs
- Ingredients:
 - Nodes: positions between words
 - Edges: spans of words with labels, represent the set of trees over those words rooted at x
 - A chart: records which edges we've built
 - An agenda: a holding pen for edges (a queue)
- We're going to figure out:
 - What edges can we build?
 - All the ways we built them.



Word Edges

- An edge found for the first time is called discovered. Edges go into the agenda on discovery.
- To initialize, we discover all word edges.

AGENDA

critics[0,1], write[1,2], reviews[2,3], with[3,4], computers[4,5]

CHART [EMPTY]



Unary Projection

- When we pop an word edge off the agenda, we check the lexicon to see what tag edges we can build from it

critics[0,1] write[1,2] reviews[2,3] with[3,4] computers[4,5]
 NNS[0,1] VBP[1,2] NNS[2,3] IN[3,4] NNS[3,4]



The "Fundamental Rule"

- When we pop edges off of the agenda:
 - Check for unary projections (NNS → critics, NP → NNS)

$Y[i,j]$ with $X \rightarrow Y$ forms $X[i,j]$

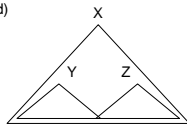
- Combine with edges already in our chart (this is sometimes called the fundamental rule)

$Y[i,j]$ and $Z[j,k]$ with $X \rightarrow YZ$ form $X[i,k]$

- Enqueue resulting edges (if newly discovered)
- Record backtraces (called traversals)
- Stick the popped edge in the chart

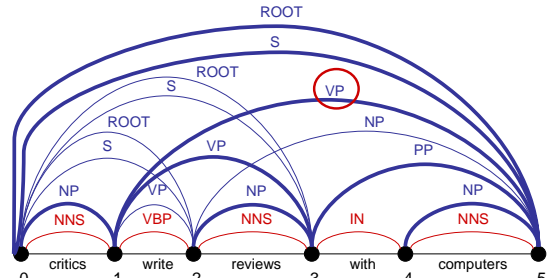
- Queries a chart must support:

- Is edge $X[i,j]$ in the chart?
- What edges with label Y end at position j ?
- What edges with label Z start at position i ?



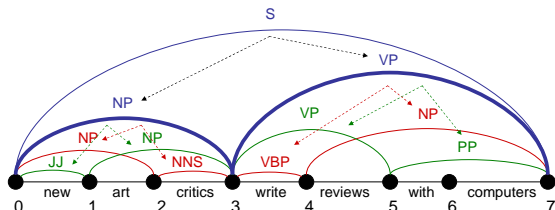
An Example

NNS[0,1] VBP[1,2] NNS[2,3] IN[3,4] NNS[3,4] NP[0,1] VP[1,2] NP[2,3] NP[4,5] S[0,2]
 VP[1,3] PP[3,5] ROOT[0,2] S[0,3] VP[1,5] NP[2,5] ROOT[0,3] S[0,5] ROOT[0,5]



Exploiting Substructure

- Each edge records all the ways it was built (locally)
 - Can recursively extract trees
 - A chart may represent too many parses to enumerate (how many?)



Order Independence

- A nice property:
 - It doesn't matter what policy we use to order the agenda (FIFO, LIFO, random).
 - Why? Invariant: before popping an edge:
 - Any edge $X[i,j]$ that can be directly built from chart edges and a single grammar rule is either in the chart or in the agenda.
 - Convince yourselves this invariant holds!
 - This will not be true once we get weighted parsers.

Empty Elements

- Sometimes we want to posit nodes in a parse tree that don't contain any pronounced words:

I want John to parse this sentence

I want [] to parse this sentence

- **These are easy to add to our chart parser!**
 - For each position i , add the "word" edge $\epsilon:[i,i]$
 - Add rules like $NP \rightarrow \epsilon$ to the grammar
 - That's it!

