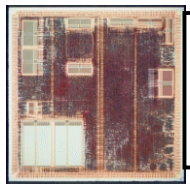
CS162
Operating Systems and Systems Programming
Lecture 1

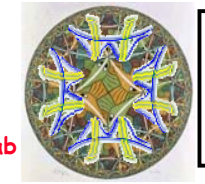
What is an Operating System?

August 27th, 2008
Prof. John Kubiatowicz
http://inst.eecs.berkeley.edu/~cs162

Who am I?

- Professor John Kubiatowicz (Prof "Kubi")
  - Background in Hardware Design
    - » Alewife project at MIT
    - » Designed CMMU, Modified SPAR C processor
    - » Helped to write operating system
  - Background in Operating Systems
    - » Worked for Project Athena (MIT)
    - » OS Developer (device drivers, network file systems)
    - » Worked on Clustered High-Availability systems (CLAM Associates)
    - » OS lead researcher for the new Berkeley PARLab (Tessellation OS). More later.
  - Peer-to-Peer
    - » OceanStore project – Store your data for 1000 years
    - » Tapestry and Bamboo – Find you data around globe
  - Quantum Computing
    - » Well, this is just cool, but probably not apropos


Alewife


Tessellation

OceanStore

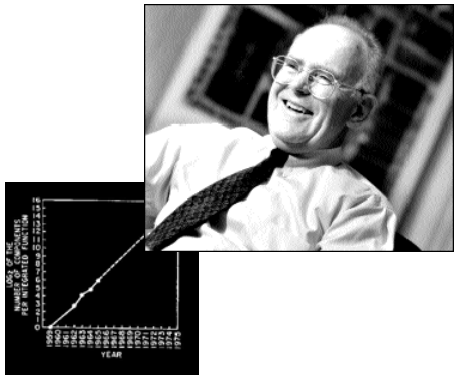
Goals for Today

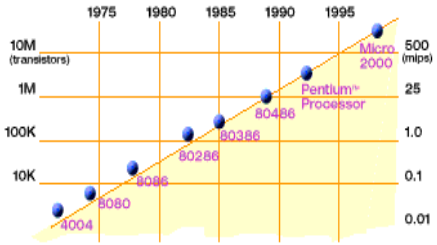
- What is an Operating System?
  - And – what is it not?
- Examples of Operating Systems design
- Why study Operating Systems?
- Oh, and "How does this class operate?"

Interactive is important!
Ask Questions!

Note: Some slides and/or pictures in the following are adapted from slides ©2005 Silberschatz, Galvin, and Gagne. Slides courtesy of Kubiatowicz, AJ Shankar, George Necula, Alex Aiken, Eric Brewer, Ras Bodik, Ion Stoica, Doug Tygar, and David Wagner.

Technology Trends: Moore's Law





2X transistors/Chip Every 1.5 years
Called "Moore's Law"

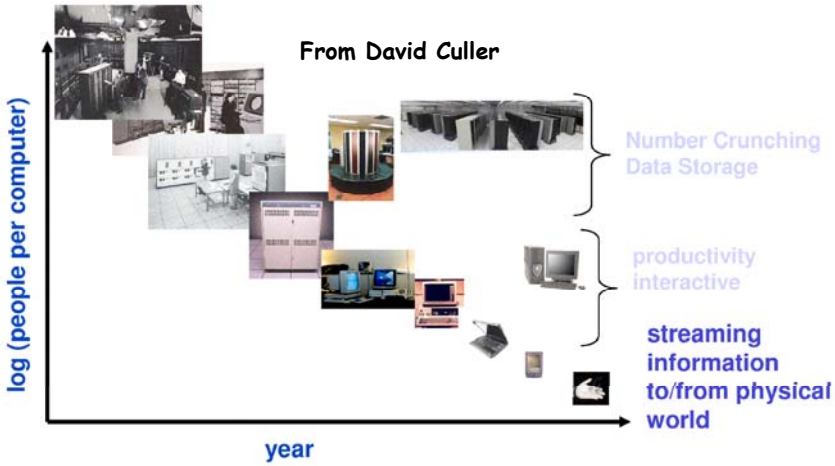
Gordon Moore (co-founder of Intel) predicted in 1965 that the transistor density of semiconductor chips would double roughly every 18 months.

Microprocessors have become smaller, denser, and more powerful.

Societal Scale Information Systems

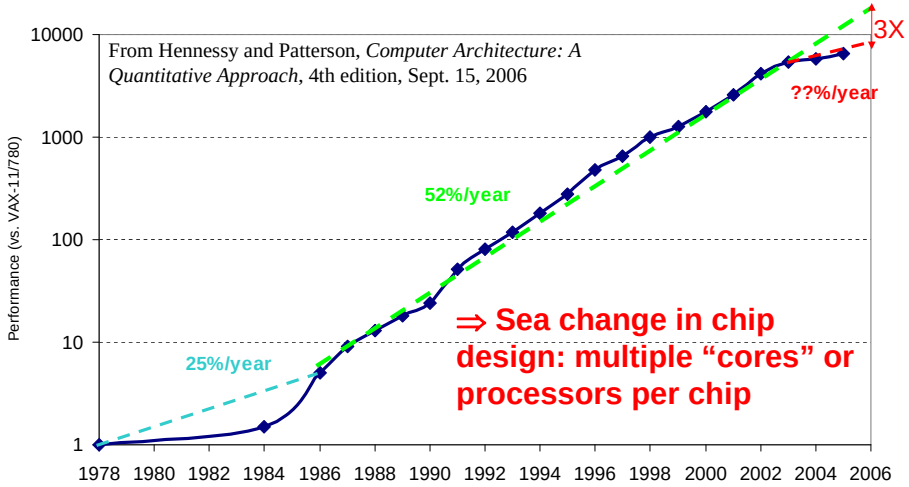
- The world is a large parallel system
  - Microprocessors in everything
  - Vast infrastructure behind them

People-to-Computer Ratio Over Time



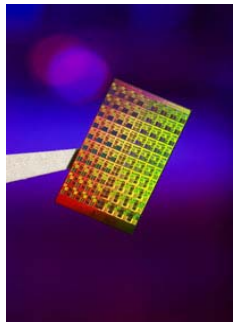
- Today: Multiple CPUs/person!
  - Approaching 100s?

New Challenge: Slowdown in Joy's law of Performance



- VAX : 25%/year 1978 to 1986
- RISC + x86: 52%/year 1986 to 2002
- RISC + x86: ??%/year 2002 to present

ManyCore Chips: The future is here

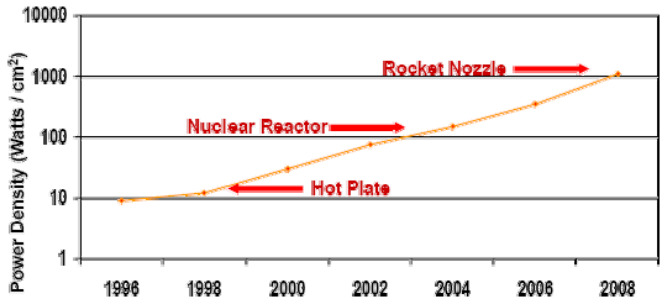


- Intel 80-core multicore chip (Feb 2007)
  - 80 simple cores
  - Two floating point engines /core
  - Mesh-like "network-on-a-chip"
  - 100 million transistors
  - 65nm feature size

Frequency	Voltage	Power	Bandwidth	Performance
3.16 GHz	0.95 V	62W	1.62 Terabits/s	1.01 Teraflops
5.1 GHz	1.2 V	175W	2.61 Terabits/s	1.63 Teraflops
5.7 GHz	1.35 V	265W	2.92 Terabits/s	1.81 Teraflops

- "ManyCore" refers to many processors/chip
  - 64? 128? Hard to say exact boundary
- How to program these?
  - Use 2 CPUs for video/audio
  - Use 1 for word processor, 1 for browser
  - 76 for virus checking???
- Parallelism must be exploited at all levels

Another Challenge: Power Density

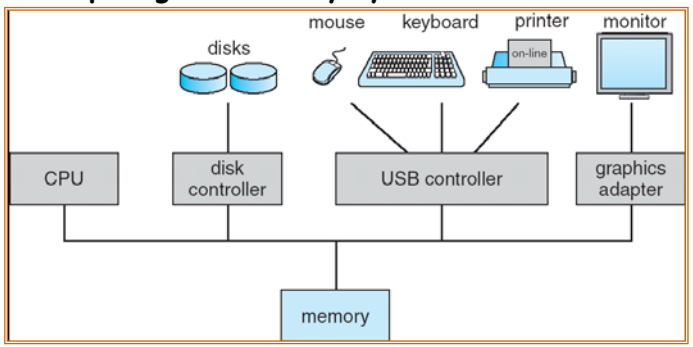


Power Density Becomes Too High to Cool Chips Inexpensively

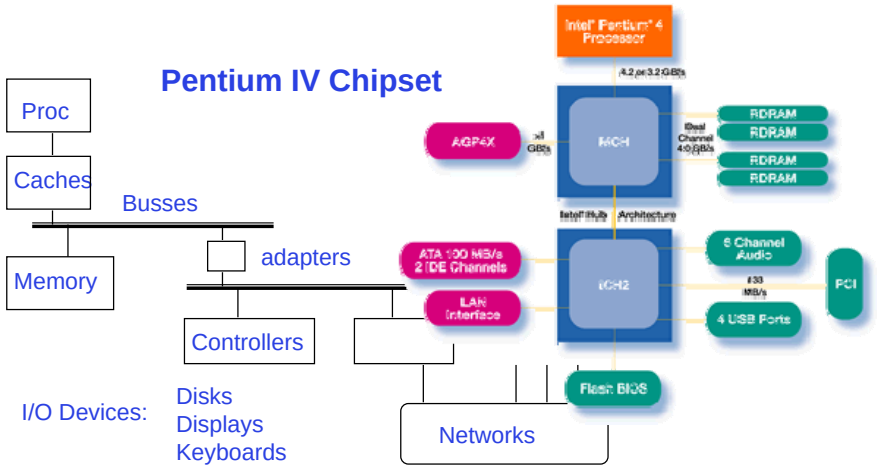
- Moore's Law Extrapolation
  - Potential power density reaching amazing levels!
- Flip side: Battery life very important
  - Moore's law can yield more functionality at equivalent (or less) total energy consumption

Computer System Organization

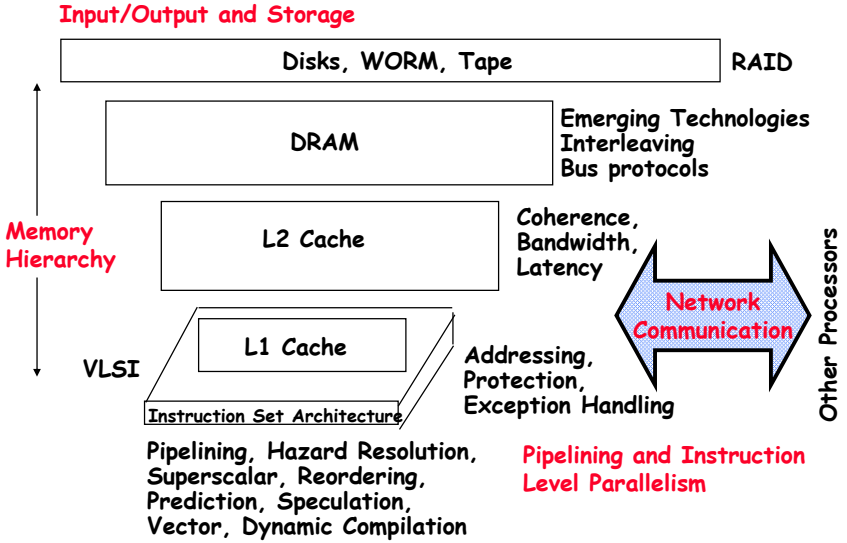
- Computer-system operation
  - One or more CPUs, device controllers connect through common bus providing access to shared memory
  - Concurrent execution of CPUs and devices competing for memory cycles



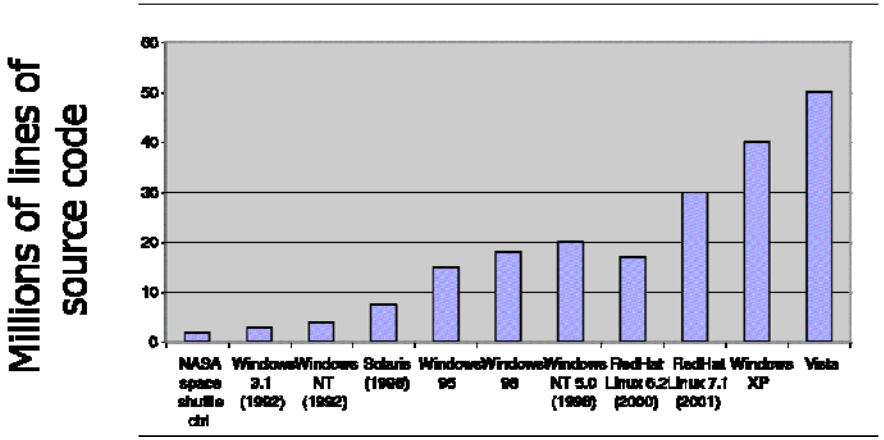
Functionality comes with great complexity!



Sample of Computer Architecture Topics

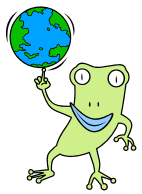


Increasing Software Complexity



From MIT's 6.033 course

Example: Some Mars Rover ("Pathfinder") Requirements



- Pathfinder hardware limitations/complexity:
  - 20Mhz processor, 128MB of DRAM, VxWorks OS
  - cameras, scientific instruments, batteries, solar panels, and locomotion equipment
  - Many independent processes work together
- Can't hit reset button very easily!
  - Must reboot itself if necessary
  - Always able to receive commands from Earth
- Individual Programs must not interfere
  - Suppose the MUT (Martian Universal Translator Module) buggy
  - Better not crash antenna positioning software!
- Further, all software may crash occasionally
  - Automatic restart with diagnostics sent to Earth
  - Periodic checkpoint of results saved?
- Certain functions time critical:
  - Need to stop before hitting something
  - Must track orbit of Earth for communication

How do we tame complexity?

- Every piece of computer hardware different
  - Different CPU
    - » Pentium, PowerPC, ColdFire, ARM, MIPS
  - Different amounts of memory, disk, …
  - Different types of devices
    - » Mice, Keyboards, Sensors, Cameras, Fingerprint readers
  - Different networking environment
    - » Cable, DSL, Wireless, Firewalls,…
- Questions:
  - Does the programmer need to write a single program that performs many independent activities?
  - Does every program have to be altered for every piece of hardware?
  - Does a faulty program crash everything?
  - Does every program have access to all hardware?

OS Tool: Virtual Machine Abstraction

Application

Virtual Machine Interface

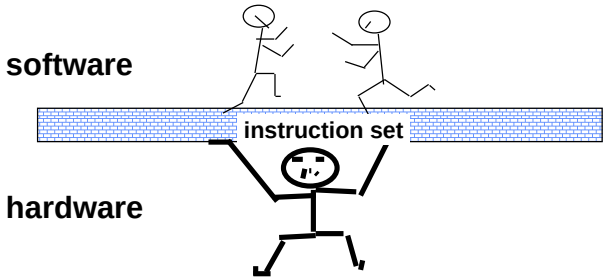
Operating System

Physical Machine Interface

Hardware

- Software Engineering Problem:
  - Turn hardware/software quirks ⇒ what programmers want/need
  - Optimize for convenience, utilization, security, reliability, etc…
- For Any OS area (e.g. file systems, virtual memory, networking, scheduling):
  - What's the hardware interface? (physical reality)
  - What's the application interface? (nicer abstraction)

Interfaces Provide Important Boundaries



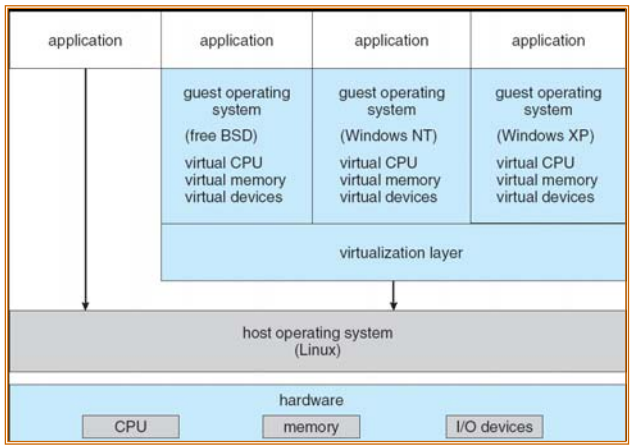
- Why do interfaces look the way that they do?
  - History, Functionality, Stupidity, Bugs, Management
  - CS152 ⇒ Machine interface
  - CS160 ⇒ Human interface
  - CS169 ⇒ Software engineering/management
- Should responsibilities be pushed across boundaries?
  - RISC architectures, Graphical Pipeline Architectures

Virtual Machines

- Software emulation of an abstract machine
  - Make it look like hardware has features you want
  - Programs from one hardware & OS on another one
- Programming simplicity
  - Each process thinks it has all memory/CPU time
  - Each process thinks it owns all devices
  - Different Devices appear to have same interface
  - Device Interfaces more powerful than raw hardware
    - » Bitmapped display ⇒ windowing system
    - » Ethernet card ⇒ reliable, ordered, networking (TCP/IP)
- Fault Isolation
  - Processes unable to directly impact other processes
  - Bugs cannot crash whole machine
- Protection and Portability
  - Java interface safe and stable across many platforms

Virtual Machines (con't): Layers of OSs

- Useful for OS development
  - When OS crashes, restricted to one VM
  - Can aid testing programs on other OSs



Course Administration

- Instructor: John Kubiatowicz (kubitron@cs.berkeley.edu)
  673 Soda Hall
  Office Hours(Tentative): M/W 2:30pm-3:30pm
- TAs: Tony Huang (cs162-ta@cory)
  Jon Whiteaker (cs162-tb@cory)
  Andrey Ermolinskiy (cs162-tc@cory)
- Labs: Second floor of Soda Hall
- Website: http://inst.eecs.berkeley.edu/~cs162
  Mirror: http://www.cs.berkeley.edu/~kubitron/cs162
- Webcast: http://webcast.berkeley.edu/courses/index.php
- Newsgroup: ucb.class.cs162 (use authnews.berkeley.edu)
- Course Email: cs162@cory.cs.berkeley.edu
- Reader: TBA (Stay tuned!)

Class Schedule

- Class Time: M/W 4:00-5:30 PM, 277 Cory Hall
  - Please come to class. Lecture notes do not have everything in them. The best part of class is the interaction!
  - Also: 5% of the grade is from class participation
- Sections:
  - Important information is in the sections
  - The sections assigned to you by Telebears are temporary!
  - Every member of a project group must be in same section
  - No sections this week

Section	Time	Location	TA
101	Tu 11:00A-12:00P	B51 Hildebrand	TBA
102	Tu 1:00P-2:00P	B51 Hildebrand	TBA
103	Tu 2:00P-3:00P	81 Evans	TBA
104	W 11:00P-12:00P	81 Evans	TBA
105	W 2:00P-3:00P	3 Evans	TBA

Textbook

- Text: Operating Systems Concepts, 8th Edition Silbershatz, Galvin, Gagne
- Online supplements
  - See "Information" link on course website
  - Includes Appendices, sample problems, etc
- Question: need 8th edition?
  - No, but has new material that we may cover
  - Completely reorganized
  - Will try to give readings from both the 7th and 8th editions on the lecture page

Topic Coverage

Textbook: Silberschatz, Galvin, and Gagne, *Operating Systems Concepts*, 8th Ed., 2008

- 1 week: Fundamentals (Operating Systems Structures)
- 1.5 weeks: Process Control and Threads
- 2.5 weeks: Synchronization and scheduling
- 2 week: Protection,Address translation, Caching
- 1 week: Demand Paging
- 1 week: File Systems
- 2.5 weeks: Networking and Distributed Systems
- 1 week: Protection and Security
- ??: Advanced topics

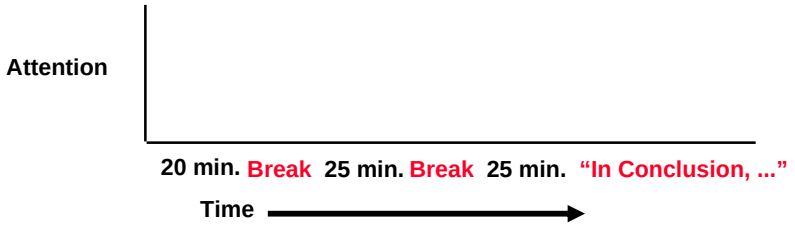
Grading

- Rough Grade Breakdown
  - Two Midterms: 15% each
    One Final: 15%
    Four Projects: 50% (i.e. 12.5% each)
    Participation: 5%
- Four Projects:
  - Phase I: Build a thread system
  - Phase II: Implement Multithreading
  - Phase III: Caching and Virtual Memory
  - Phase IV: Parallel and Distributed Systems
- Late Policy:
  - Each group has 5 "slip" days.
  - For Projects, slip days deducted from *all* partners
  - 10% off per day after slip days exhausted

Group Project Simulates Industrial Environment

- Project teams have 4 or 5 members in same discussion section
  - Must work in groups in "the real world"
- Communicate with colleagues (team members)
  - Communication problems are natural
  - What have you done?
  - What answers you need from others?
  - You must document your work!!!
  - Everyone must keep an on-line notebook
- Communicate with supervisor (TAs)
  - How is the team's plan?
  - Short progress reports are required:
    - » What is the team's game plan?
    - » What is each member's responsibility?

Typical Lecture Format



- 1-Minute Review
- 20-Minute Lecture
- 5- Minute Administrative Matters
- 25-Minute Lecture
- 5-Minute Break (water, stretch)
- 25-Minute Lecture
- Instructor will come to class early & stay after to answer questions

Lecture Goal

Interactive!!!

Computing Facilities

- Every student who is enrolled should get an account form at end of lecture
  - Gives you an account of form cs162-xx@cory
  - This account is required
    - » Most of your debugging can be done on other EECS accounts, however…
    - » All of the final runs must be done on your cs162-xx account and must run on the x86 Solaris machines
- Make sure to log into your new account this week and fill out the questions
- Project Information:
  - See the "Projects and Nachos" link off the course home page
- Newsgroup (ucb.class.cs162):
  - Read this regularly!

Academic Dishonesty Policy

- Copying all or part of another person's work, or using reference material not specifically allowed, are forms of cheating and will not be tolerated. A student involved in an incident of cheating will be notified by the instructor and the following policy will apply:

http://www.eecs.berkeley.edu/Policies/acad.dis.shtml

- The instructor may take actions such as:
  - require repetition of the subject work,
  - assign an F grade or a 'zero' grade to the subject work,
  - for serious offenses, assign an F grade for the course.
- The instructor must inform the student and the Department Chair in writing of the incident, the action taken, if any, and the student's right to appeal to the Chair of the Department Grievance Committee or to the Director of the Office of Student Conduct.
- The Office of Student Conduct may choose to conduct a formal hearing on the incident and to assess a penalty for misconduct.
- The Department will recommend that students involved in a second incident of cheating be dismissed from the University.

What does an Operating System do?

- Silerschatz and Gavin:
  "An OS is Similar to a government"
  - Begs the question: does a government do anything useful by itself?
- Coordinator and Traffic Cop:
  - Manages all resources
  - Settles conflicting requests for resources
  - Prevent errors and improper use of the computer
- Facilitator:
  - Provides facilities that everyone needs
  - Standard Libraries, Windowing systems
  - Make application programming easier, faster, less error-prone
- Some features reflect both tasks:
  - E.g. File system is needed by everyone (Facilitator)
  - But File system must be Protected (Traffic Cop)

What is an Operating System,… Really?

- Most Likely:
  - Memory Management
  - I/O Management
  - CPU Scheduling
  - Communications? (Does Email belong in OS?)
  - Multitasking/multiprogramming?
- What about?
  - File System?
  - Multimedia Support?
  - User Interface?
  - Internet Browser? ☺
- Is this only interesting to Academics??

Operating System Definition (Cont.)

- No universally accepted definition
- "Everything a vendor ships when you order an operating system" is good approximation
  - But varies wildly
- "The one program running at all times on the computer" is the kernel.
  - Everything else is either a system program (ships with the operating system) or an application program

What if we didn't have an Operating System?

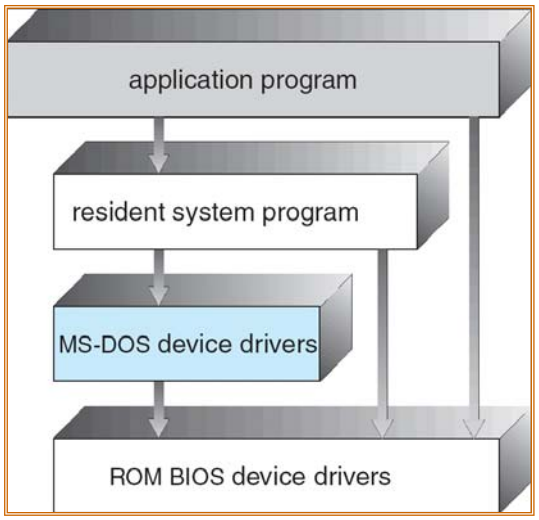
- Source Code⇒Compiler⇒Object Code⇒Hardware
- How do you get object code onto the hardware?
- How do you print out the answer?
- Once upon a time, had to Toggle in program in binary and read out answer from LED's!

Altair 8080

Simple OS: What if only one application?

- Examples:
  - Very early computers
  - Early PCs
  - Embedded controllers (elevators, cars, etc)
- OS becomes just a library of standard services
  - Standard device drivers
  - Interrupt handlers
  - Math libraries

MS-DOS Layer Structure



More thoughts on Simple OS

- What about Cell-phones, Xboxes, etc?
  - Is this organization enough?
- Can OS be encoded in ROM/Flash ROM?
- Does OS have to be software?
  - Can it be Hardware?
  - Custom Chip with predefined behavior
  - Are these even OSs?

More complex OS: Multiple Apps

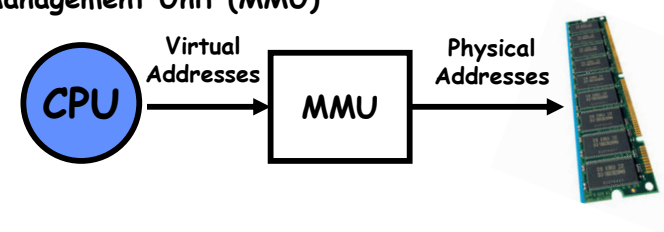
- Full Coordination and Protection
  - Manage interactions between different users
  - Multiple programs running simultaneously
  - Multiplex and protect Hardware Resources
    - » CPU, Memory, I/O devices like disks, printers, etc
- Facilitator
  - Still provides Standard libraries, facilities
- Would this complexity make sense if there were only one application that you cared about?

Example: Protecting Processes from Each Other

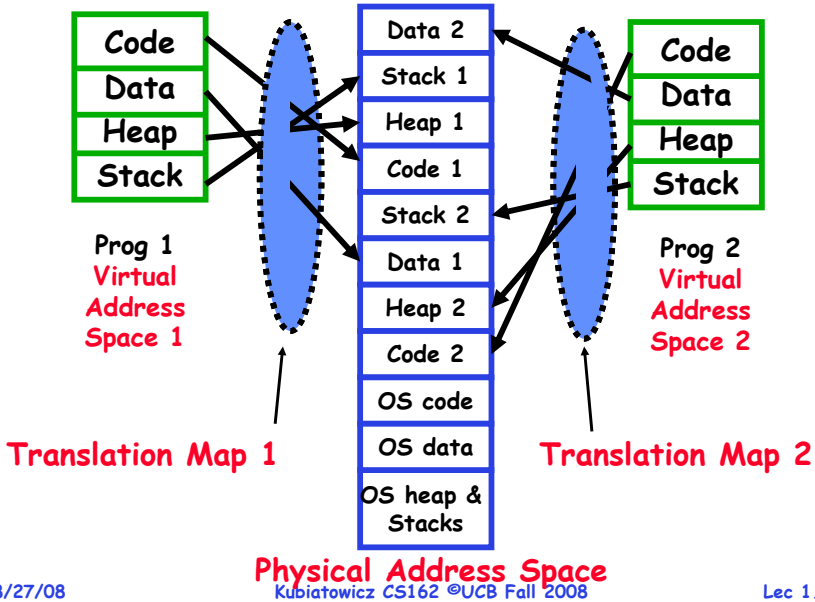
- Problem: Run multiple applications in such a way that they are protected from one another
- Goal:
  - Keep User Programs from Crashing OS
  - Keep User Programs from Crashing each other
  - [Keep Parts of OS from crashing other parts?]
- (Some of the required) Mechanisms:
  - Address Translation
  - Dual Mode Operation
- Simple Policy:
  - Programs are not allowed to read/write memory of other Programs or of Operating System

Address Translation

- Address Space
  - A group of memory addresses usable by something
  - Each program (process) and kernel has potentially different address spaces.
- Address Translation:
  - Translate from Virtual Addresses (emitted by CPU) into Physical Addresses (of memory)
  - Mapping *often* performed in Hardware by Memory Management Unit (MMU)

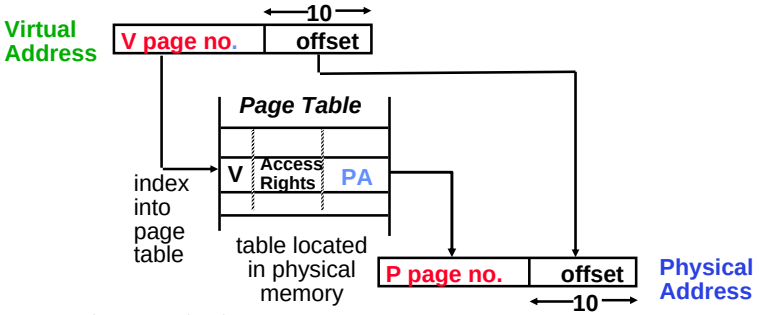


Example of Address Translation



Address Translation Details

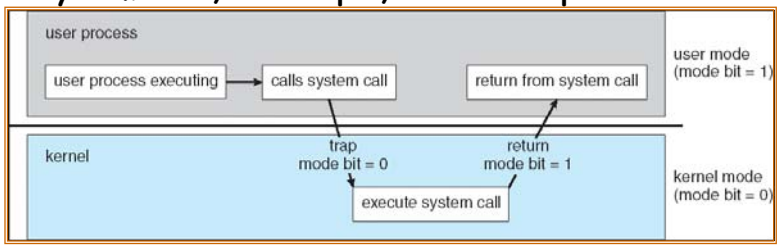
- For now, assume translation happens with table (called a Page Table):



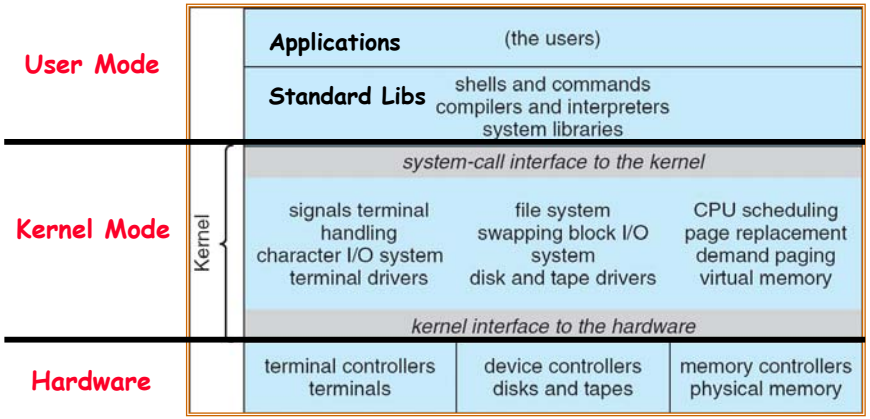
- Translation helps protection:
  - Control translations, control access
  - Should Users be able to change Page Table???

Dual Mode Operation

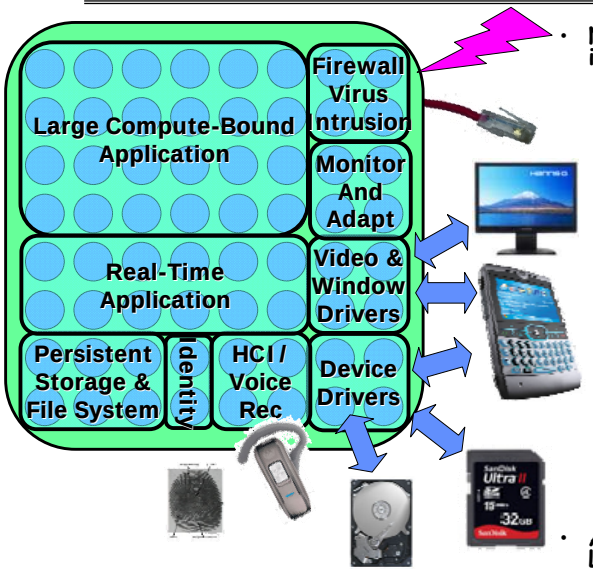
- Hardware provides at least two modes:
  - "Kernel" mode (or "supervisor" or "protected")
  - "User" mode: Normal programs executed
- Some instructions/ops prohibited in user mode:
  - Example: cannot modify page tables in user mode
    - » Attempt to modify ⇒ Exception generated
- Transitions from user mode to kernel mode:
  - System Calls, Interrupts, Other exceptions



UNIX System Structure



New Structures for Multicore chips? Tessellation: The Exploded OS



- Normal Components split into pieces
  - Device drivers (Security/Reliability)
  - Network Services (Performance)
    - » TCP/IP stack
    - » Firewall
    - » Virus Checking
    - » Intrusion Detection
  - Persistent Storage (Performance, Security, Reliability)
  - Monitoring services
    - » Performance counters
    - » Introspection
  - Identity/Environment services (Security)
    - » Biometric, GPS, Possession Tracking
- Applications Given Larger Partitions
  - Freedom to use resources arbitrarily

OS Systems Principles

- OS as illusionist:
  - Make hardware limitations go away
  - Provide illusion of dedicated machine with infinite memory and infinite processors
- OS as government:
  - Protect users from each other
  - Allocate resources efficiently and fairly
- OS as complex system:
  - Constant tension between simplicity and functionality or performance
- OS as history teacher
  - Learn from past
  - Adapt as hardware tradeoffs change

Why Study Operating Systems?

- Learn how to build complex systems:
  - How can you manage complexity for future projects?
- Engineering issues:
  - Why is the web so slow sometimes? Can you fix it?
  - What features should be in the next mars Rover?
  - How do large distributed systems work? (Kazaa, etc)
- Buying and using a personal computer:
  - Why different PCs with same CPU behave differently
  - How to choose a processor (Opteron, Itanium, Celeron, Pentium, Hexium)? [Ok, made last one up]
  - Should you get Windows XP, 2000, Linux, Mac OS …?
  - Why does Microsoft have such a bad name?
- Business issues:
  - Should your division buy thin-clients vs PC?
- Security, viruses, and worms
  - What exposure do you have to worry about?

"In conclusion…"

- Operating systems provide a virtual machine abstraction to handle diverse hardware
- Operating systems coordinate resources and protect users from each other
- Operating systems simplify application development by providing standard services
- Operating systems can provide an array of fault containment, fault tolerance, and fault recovery
- CS162 combines things from many other areas of computer science –
  - Languages, data structures, hardware, and algorithms