**CS162** Application **Operating Systems and** - Virtual Machine Interface Systems Programming **Operating System** Lecture 2 **Physical Machine Interface** Hardware History of the World Parts 1-5• Software Engineering Problem: **Operating Systems Structures** - Turn hardware/software guirks  $\Rightarrow$ what programmers want/need September 3<sup>rd</sup>, 2008 - Optimize for convenience, utilization, security, reliability, etc... Prof. John Kubiatowicz • For Any OS area (e.g. file systems, virtual memory, http://inst.eecs.berkeley.edu/~cs162 networking, scheduling): - What's the hardware interface? (physical reality) - What's the application interface? (nicer abstraction) 9/03/08 Kubiatowicz CS162 ©UCB Fall 2008 Lec 2.2

# Goals for Today

- Finish Protection Example
- History of Operating Systems
  - Really a history of resource-driven choices
- Operating Systems Structures
- Operating Systems Organizations
- Abstractions and layering

Note: Some slides and/or pictures in the following are adapted from slides ©2005 Silberschatz, Galvin, and Gagne. Many slides generated from lecture notes by Joseph.

# **Example: Protecting Processes from Each Other**

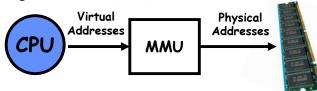
**Review: Virtual Machine Abstraction** 

- Problem: Run multiple applications in such a way that they are protected from one another
- Goal:
  - Keep User Programs from Crashing OS
  - Keep User Programs from Crashing each other
  - [Keep Parts of OS from crashing other parts?]
- (Some of the required) Mechanisms:
  - Address Translation
  - Dual Mode Operation
- Simple Policy:
  - Programs are not allowed to read/write memory of other Programs or of Operating System

Lec 2.3

### **Example: Address Translation**

- Address Space
  - A group of memory addresses usable by something
  - Each program (process) and kernel has potentially different address spaces.
- Address Translation:
  - Translate from Virtual Addresses (emitted by CPU) into Physical Addresses (of memory)
  - Mapping *often* performed in Hardware by Memory Management Unit (MMU)

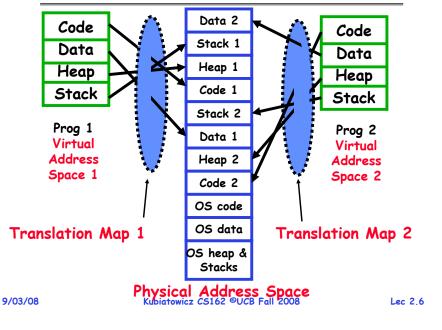


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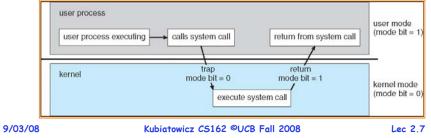
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### Example: Example of Address Translation



### Example: Dual Mode Operation

- Hardware provides at least two modes:
  - "Kernel" mode (or "supervisor" or "protected")
  - "User" mode: Normal programs executed
- Some instructions/ops prohibited in user mode:
  - Example: cannot modify page tables in user mode
    » Attempt to modify ⇒ Exception generated
- $\boldsymbol{\cdot}$  Transitions from user mode to kernel mode:
  - System Calls, Interrupts, Other exceptions



### **UNIX System Structure**

User Mode		Applications	(the users)				
User Mode		Standard Libs shells and commands compilers and interpreters system libraries					
	ſ	system	n-call interface to the ke	ernel			
Kernel Mode	Kernel	signals terminal handling character I/O system terminal drivers	file system swapping block I/O system disk and tape drivers	CPU scheduling page replacement demand paging virtual memory			
		kernel interface to the hardware					
Hardware		terminal controllers terminals	device controllers disks and tapes	memory controllers physical memory			

### Moore's Law Change Drives OS Change

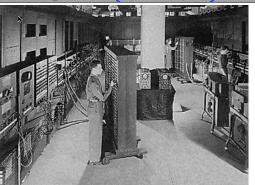
	1981	2006	Factor
CPU MHz,	10	3200x4	1,280
Cycles/inst	3—10	0.25-0.5	6—40
DRAM capacity	128KB	4GB	32,768
Disk capacity	10MB	1TB	100,000
Net bandwidth	9600 b/s	1 Gb/s	110,000
# addr bits	16	32	2
#users/machine	10s	≤ <b>1</b>	≤ <b>0</b> .1
Price	\$25,000	\$4,000	0.2

# Typical academic computer 1981 vs 2006

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- "The machine designed by Drs. Eckert and Mauchly was a monstrosity. When it was finished, the ENIAC filled an entire room, weighed thirty tons, and consumed two hundred kilowatts of power."
- http://ei.cs.vt.edu/~history/ENIAC.Richey.HTML

### Moore's law effects

- Nothing like this in any other area of business • Transportation in over 200 years: - 2 orders of magnitude from horseback @10mph to Concorde @1000mph - Computers do this every decade (at least until 2002)! • What does this mean for us? - Techniques have to vary over time to adapt to changing tradeoffs • I place a lot more emphasis on principles - The key concepts underlying computer systems - Less emphasis on facts that are likely to change over the next few years...  $\cdot$  Let's examine the way changes in \$/MIP has radically changed how OS's work 9/03/08 Kubiatowicz CS162 ©UCB Fall 2008 Lec 2,10 History Phase 1 (1948-1970) Hardware Expensive, Humans Cheap • When computers cost millions of \$'s, optimize for more efficient use of the hardware!
  - User at console: one user at a time

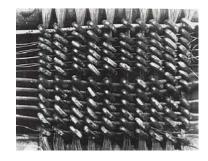
- Lack of interaction between user and computer

- Batch monitor: load program, run, print
- Optimize to better use hardware
  - When user thinking at console, computer idle  $\Rightarrow$  BAD!
  - Feed computer batches and make users wait
  - Autograder for this course is similar
- No protection: what if batch program has bug?

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Lec 2.9

### Core Memories (1950s & 60s)



The first magnetic core memory, from the IBM 405 Alphabetical Accounting Machine.

- · Core Memory stored data as magnetization in iron rings
  - Iron "cores" woven into a 2-dimensional mesh of wires
  - Origin of the term "Dump Core"
  - Rumor that IBM consulted Life Saver company
- See: http://www.columbia.edu/acis/history/core.html

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### A Multics System (Circa 1976)



- The 6180 at MIT IPC, skin doors open, circa 1976:
  - "We usually ran the machine with doors open so the operators could see the AQ register display, which gave you an idea of the machine load, and for convenient access to the EXECUTE button, which the operator would push to enter BOS if the machine crashed "

 http://www.multicians.org/multics-stories.html Kubiatowicz CS162 ©UCB Fall 2008 9/03/08 Lec 2,15

### History Phase $1\frac{1}{2}$ (late 60s/early 70s)

- Data channels, Interrupts: overlap I/O and compute - DMA - Direct Memory Access for I/O devices - I/O can be completed asynchronously • Multiprogramming: several programs run simultaneously - Small jobs not delayed by large jobs - More overlap between I/O and CPU
  - Need memory protection between programs and/or OS
- Complexity gets out of hand:
  - Multics: announced in 1963, ran in 1969
    - » 1777 people "contributed to Multics" (30-40 core dev)
    - » Turing award lecture from Fernando Corbató (key researcher): "On building systems that will fail"
  - OS 360: released with 1000 known bugs (APARs) » "Anomalous Program Activity Report"

### • OS finally becomes an important science:

- How to deal with complexity???
- UNIX based on Multics, but vastly simplified

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### Early Disk History Model 3370 Model 3340 hard disk 1973 1979 7.7 1.7 2,300 140 1973: 1979: 1.7 Mbit/sq. in 7.7 Mbit/sq. in 140 MBytes 2,300 MBytes

Contrast: Seagate 1TB, 164 GB/SQ in,  $3\frac{1}{2}$  in disk, **4** platters

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### Administrivia

- Cs162-xx accounts:
  - Make sure you got an account form
    - » We have more forms for those of you who didn't get one
  - If you haven't logged in yet, you need to do so
- Nachos readers:
  - TBA: Will be down at Copy Central on Hearst
  - Will include lectures and printouts of all of the code
- Video "Screencast" archives available off lectures page
  - Just click on the title of a lecture for webcast
  - Only works for lectures that I have already given!
- No slip days on first design document for each phase
  - Need to get design reviews in on time
- Don't know Java well?
  - Talk CS 9G self-paced Java course

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History Phase 2 (1970 – 1985) Hardware Cheaper, Humans Expensive

- Computers available for tens of thousands of dollars instead of millions
- OS Technology maturing/stabilizing
- Interactive timesharing:
  - Use cheap terminals (~\$1000) to let multiple users interact with the system at the same time
  - Sacrifice CPU time to get better response time
  - Users do debugging, editing, and email online
- Problem: Thrashing
  - Performance very non-linear response with load
  - Thrashing caused by many factors including
    - » Swapping, queueing



### Users

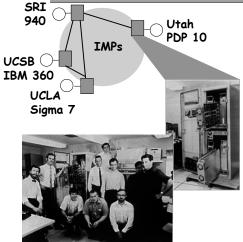
Lec 2,19

### Administriva: Time for Project Signup

- Project Signup: Watch "Group/Section Assignment Link"
  - 4-5 members to a group
    - » Everyone in group must be able to *actually* attend same section
    - » The sections assigned to you by Telebears are temporary!
  - Only submit once per group!
    - » Everyone in group must have logged into their cs162-xx accounts once before you register the group
    - » Make sure that you select at least 2 potential sections
    - » Due tomorrow: Thursday 9/4 by 11:59pm
- Sections:
  - Watch for section assignments next Monday
  - Attend new sections next week

Section	Time	Location	TA
101	Tu 11:00-12:00P	B56 Hildebrand	Andrey Ermolinskiy
102	Tu 1:00-2:00P	B56 Hildebrand	Jon Whiteaker
103	Tu 2:00-3:00P	87 Evans	Andrey Ermolinskiy
104	W 11:00-12:00P	87 Evans	Tony Huang
105	W 2:00-3:00P	3 Evans	Jon Whiteaker

# The ARPANet (1968-1970's)

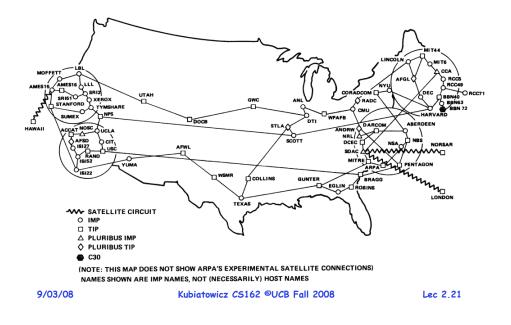


BBN team that implemented the interface message processor



- Paul Baran
  - RAND Corp, early 1960s
  - Communications networks that would survive a major enemy attack
- ARPANet: Research vehicle for "Resource Sharing Computer Networks"
  - 2 September 1969: UCLA first node on the ARPANet
- December 1969: 4 nodes connected by 56 kbps phone lines
- 1971: First Email
- 1970's: <100 computers

### **ARPANET GEOGRAPHIC MAP, OCTOBER 1980**



### **ARPANet Evolves into Internet**

• First E-mail SPAM message: 1 May 1978 12:33 EDT

- 80-83: TCP/IP, DNS; ARPANET and MILNET split
- 85-86: NSF builds NSFNET as backbone, links 6 Supercomputer centers, 1.5 Mbps, 10,000 computers
- 87-90: link regional networks, NSI (NASA), ESNet (DOE), DARTnet, TWBNet (DARPA), 100,000 computers

S	IRPANet SATNet RNet	TCP/IP	NSFNet	Deregulation & Commercialization WWW	ISP ASP AIP	
1965	19	975	1985	1995		2005
	Net: Satelite ı 2t: Radio Netw					

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### What is a Communication Network? (End-system Centric View)

- Network offers one basic service: move information
  - Bird, fire, messenger, truck, telegraph, telephone, Internet ...
  - Another example, transportation service: move objects
    - » Horse, train, truck, airplane ...
- What distinguish different types of networks?
  - The services they provide
- What distinguish the services?
  - Latency
  - Bandwidth
  - Loss rate
  - Number of end systems
  - Service interface (how to invoke the service?)
  - Others

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» Reliability, unicast vs. multicast, real-time... Kubiatowicz C5162 ©UCB Fall 2008

### What is a Communication Network? (Infrastructure Centric View)

- Communication medium: electron, photon
- Network components:
  - Links carry bits from one place to another (or maybe multiple places): fiber, copper, satellite, ...
  - Interfaces attach devices to links
  - Switches/routers interconnect links: electronic/optic, crossbar/Banyan
  - Hosts communication endpoints: workstations, PDAs, cell phones, toasters
- Protocols rules governing communication between nodes
  - TCP/IP, ATM, MPLS, SONET, Ethernet, X.25
- Applications: Web browser, X Windows, FTP, ...

# Network Components (Examples)

# LinksInterfacesSwitches/routersFibersEthernet cardLarge routerImage: Constant of the second of the second

# Types of Networks

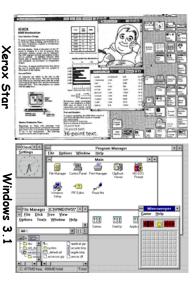
 Geographical distance - Local Area Networks (LAN): Ethernet, Token ring, FDDI - Metropolitan Area Networks (MAN): DQDB, SMDS - Wide Area Networks (WAN): X.25, ATM, frame relav - Caveat: LAN, MAN, WAN may mean different thinas » Service, network technology, networks Information type - Data networks vs. telecommunication networks Application type - Special purpose networks: airline reservation network, banking network, credit card network, telephony - General purpose network: Internet 9/03/08 Kubiatowicz CS162 ©UCB Fall 2008 Lec 2.26

### History Phase 3 (1981— ) Hardware Very Cheap, Humans Very Expensive

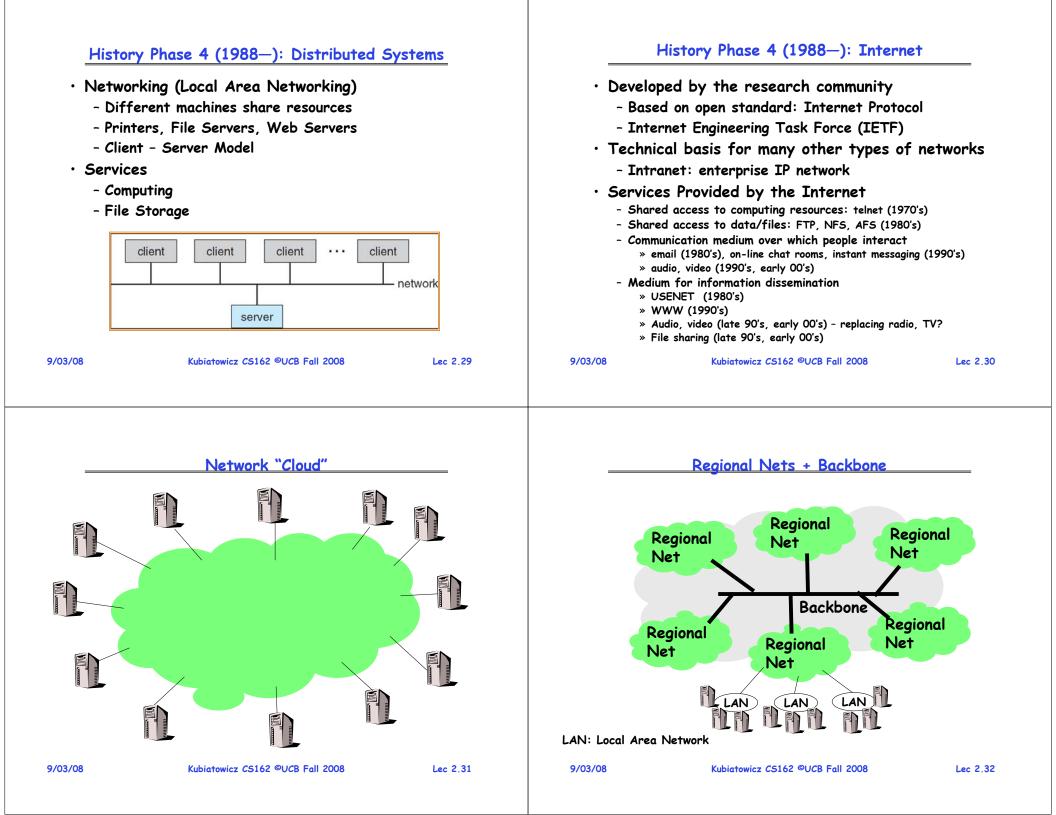
- Computer costs \$1K, Programmer costs \$100K/year
  - If you can make someone 1% more efficient by giving them a computer, it's worth it!
  - Use computers to make people more efficient
- Personal computing:
  - Computers cheap, so give everyone a PC
- Limited Hardware Resources Initially:
  - OS becomes a subroutine library
  - One application at a time (MSDOS, CP/M, ...)
- Eventually PCs become powerful:
  - OS regains all the complexity of a "big" OS
  - multiprogramming, memory protection, etc (NT,OS/2)
- Question: As hardware gets cheaper does need for OS go away?

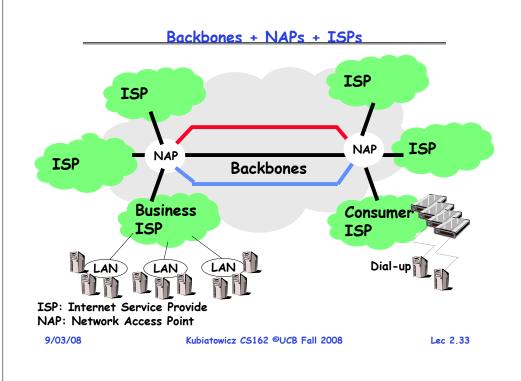
### History Phase 3 (con't) Graphical User Interfaces

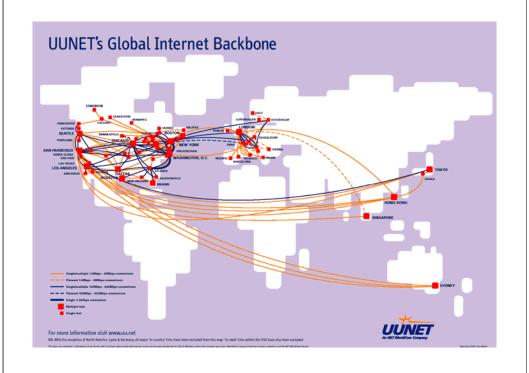
- · CS160  $\Rightarrow$  All about GUIs
- Xerox Star: 1981
  - Originally a research project (Alto)
  - First "mice", "windows"
- Apple Lisa/Machintosh: 1984
  "Look and Feel" suit 1988
- Microsoft Windows:
  - Win 1.0 (1985) Single
  - Win 3.1 (1990) Level
  - Win 95 (1995)
  - Win NT (1993) HAL/Protection
  - Win 2000 (2000) No HAL/
  - Win XP (2001) Full Prot
  - Win Vista (2007)



Lec 2.27



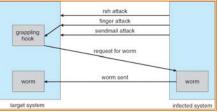




### Computers Inside the Core Head Ends AWAY @home bvad ۸D ISP Cingular Satellite C CPM (de Fixed Wireless Cell Cell Sprint AOL Cell LAN LAN LAN 9/03/08 Subiatowicz CS162 ©UCB Fall 2008 Lec 2.35

### The Morris Internet Worm (1988)

- Internet worm (Self-reproducing)
  - Author Robert Morris, a first-year Cornell grad student
  - Launched close of Workday on November 2, 1988
  - Within a few hours of release, it consumed resources to the point of bringing down infected machines



- Techniques
  - Exploited UNIX networking features (remote access)
  - Bugs in *finger* (buffer overflow) and *sendmail* programs (debug mode allowed remote login)
  - Dictionary lookup-based password cracking
  - Grappling hook program uploaded main worm program

### LoveLetter Virus (May 2000) History Phase 5 (1995—): Mobile Systems • E-mail message with Ubiguitous Mobile Devices **VBScript** (simplified Visual - Laptops, PDAs, phones Basic) - Small, portable, and inexpensive Relies on Windows Scripting Host » Recently twice as many smart phones as PDAs 😡 Reply 🕵 Reply to All 😡 Forward 🚑 🖹 Signature + 🚦 🔻 🕐 - Enabled by default in » Many computers/person! Topy Austin [topy@topyaustin.com] Sept: Thu 04/05/00 10:44 PM Win98/2000 webmaster@tonvaustin.com - Limited capabilities (memory, CPU, power, etc...) Cer User clicks on Subject: ILOVEYOU attachment -> infected! Wireless/Wide Area Networking kindly check the attached LOVELETTER coming from me. - E-mails itself to everyone - Leveraging the infrastructure in Outlook address book - Huge distributed pool of resources extend devices - Replaces some files with a copy of itself - Traditional computers split into pieces. Wireless keyboards/mice, CPU distributed, storage remote - Searches all drives S - Downloads password LOVE-LET. · Peer-to-peer systems (10KB) cracking program - Many devices with equal responsibilities work together 60-80% of US companies - Components of "Operating System" spread across globe infected and 100K European servers Kubiatowicz CS162 ©UCB Fall 2008 Lec 2.37 9/03/08 Kubiatowicz CS162 ©UCB Fall 2008 Lec 2.38 CITRIS's Model: Datacenter is the Computer A Societal Scale Information System Center for Information Technology Research in the • (From Luiz Barroso's talk at RAD Lab 12/11) Interest of Society Google program == Web search, Gmail,... • The Network is the OS Google computer ==

- Functionality spread throughout network



# Aobile, Ubiguitous Systems

### MEM Sensor Nets 9/03/08

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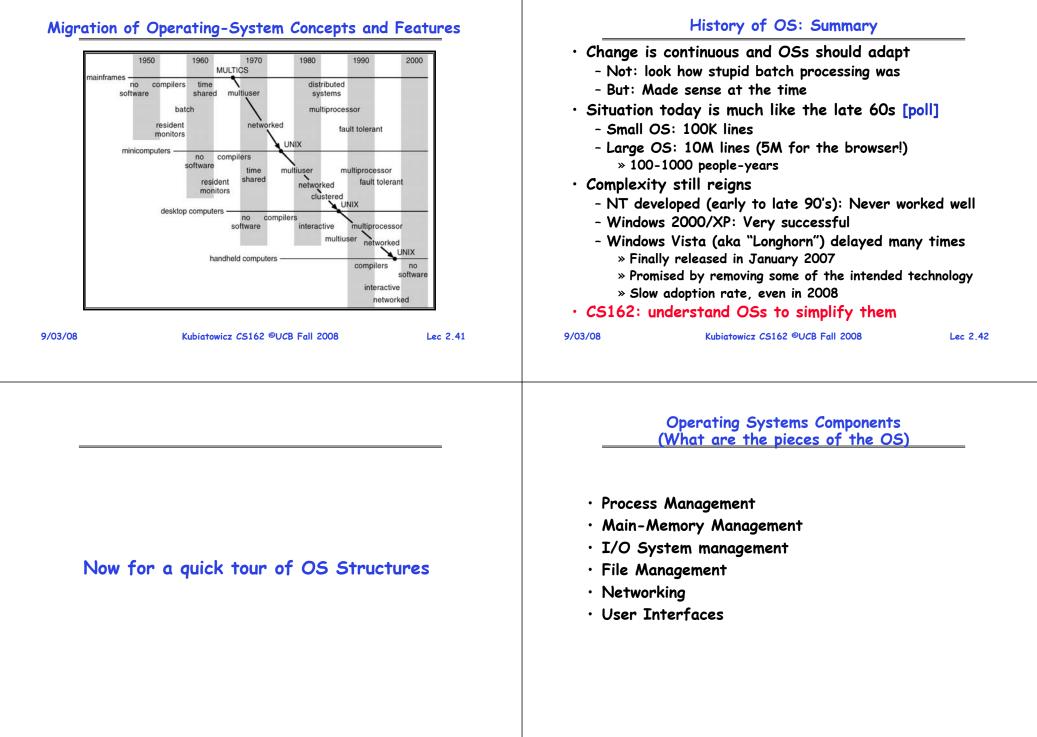
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the next few years

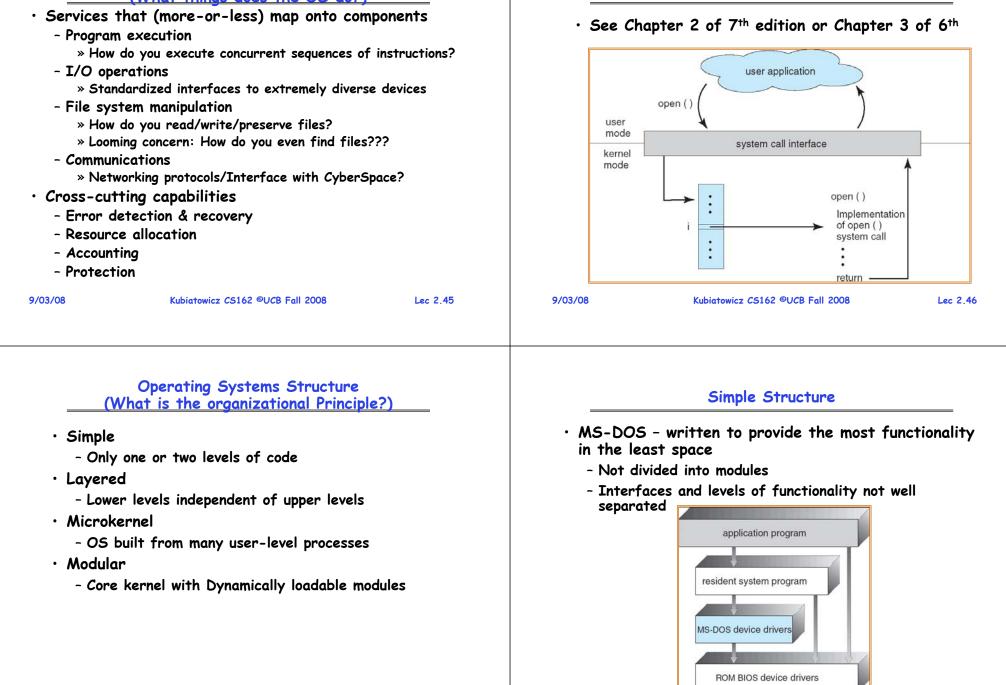
- Thousands of computers, networking, storage

• Warehouse-sized facilities and workloads may be

unusual today but are likely to be more common in



### Operating System Services (What things does the OS do?)



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System Calls (What is the API)

### UNIX: Also "Simple" Structure

### **UNIX System Structure**

Standard Libs compilers and interpreters

**Applications** 

signals terminal

handling

terminal drivers

character I/O system

terminal controllers

terminals

User Mode

Kernel Mode

Hardware

Kernel

(the users)

shells and commands

system libraries

system-call interface to the kernel

file system

swapping block I/O

system

disk and tape drivers

kernel interface to the hardware

device controllers

disks and tapes

**CPU** scheduling

page replacement

demand paging

virtual memory

memory controllers

physical memory

- UNIX limited by hardware functionality
- Original UNIX operating system consists of two separable parts:
  - Systems programs
  - The kernel
    - » Consists of everything below the system-call interface and above the physical hardware
    - » Provides the file system, CPU scheduling, memory management, and other operating-system functions;
    - » Many interacting functions for one level

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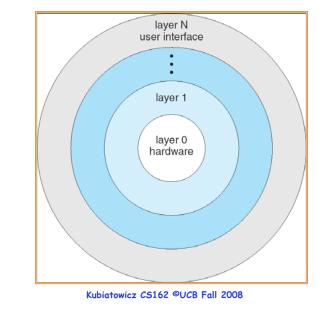
### Layered Structure

- Operating system is divided many layers (levels)
  - Each built on top of lower layers
  - Bottom layer (layer 0) is hardware
  - Highest layer (layer N) is the user interface
- Each layer uses functions (operations) and services of only lower-level layers
  - Advantage: modularity  $\Rightarrow$  Easier debugging/Maintenance
  - Not always possible: Does process scheduler lie above or below virtual memory layer?
    - » Need to reschedule processor while waiting for paging
    - » May need to page in information about tasks
- Important: Machine-dependent vs independent layers
  - Easier migration between platforms
  - Easier evolution of hardware platform
  - Good idea for you as well!



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### Layered Operating System

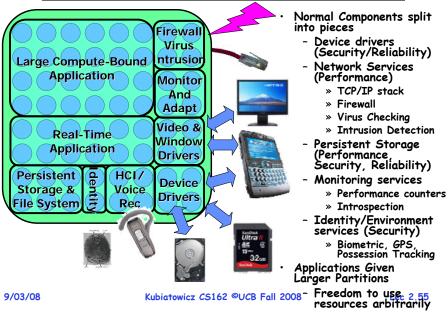


### Microkernel Structure

- Moves as much from the kernel into "*user*" space
  - Small core OS running at kernel level
  - OS Services built from many independent user-level processes
- Communication between modules with message passing
- Benefits:
  - Easier to extend a microkernel
  - Easier to port OS to new architectures
  - More reliable (less code is running in kernel mode)
  - Fault Isolation (parts of kernel protected from other parts)
  - More secure
- Detriments:
  - Performance overhead severe for naïve implementation

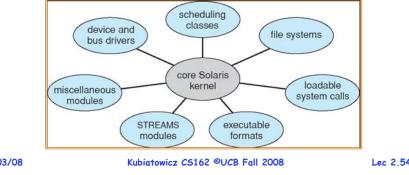
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# Partition Based Structure for Multicore chips?



# Modules-based Structure

- Most modern operating systems implement modules
  - Uses object-oriented approach
  - Each core component is separate
  - Each talks to the others over known interfaces
  - Each is loadable as needed within the kernel
- $\cdot$  Overall, similar to layers but with more flexible



### Implementation Issues (How is the OS implemented?)

- Policy vs. Mechanism
  - Policy: What do you want to do?
  - Mechanism: How are you going to do it?
  - Should be separated, since both change
- Algorithms used
  - Linear, Tree-based, Log Structured, etc...
- Event models used
  - threads vs event loops
- Backward compatability issues
  - Very important for Windows 2000/XP
- System generation/configuration
  - How to make generic OS fit on specific hardware

### Conclusion

- Rapid Change in Hardware Leads to changing OS
  - Batch ⇒ Multiprogramming ⇒ Timeshare ⇒ Graphical UI ⇒ Ubiquitous Devices ⇒ Cyberspace/Metaverse/??
- $\cdot$  OS features migrated from mainframes  $\Rightarrow$  PCs
- Standard Components and Services
  - Process Control
  - Main Memory
  - I/O
  - File System
  - UI
- Policy vs Mechanism
  - Crucial division: not always properly separated!
- Complexity is always out of control
  - However, "Resistance is NOT Useless!"

~	-		

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