CS152
Computer Architecture and Engineering
Lecture 10: Designing Single Cycle Control

September 24, 1997

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lecture slides: http://www-inst.eecs.berkeley.edu/~cs152/

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Recap: Summary from last time

- 5 steps to design a processor
  - 1. Analyze instruction set => datapath requirements
  - 2. Select set of datapath components & establish clock methodology
  - 3. Assemble datapath meeting the requirements
  - 4. Analyze implementation of each instruction to determine setting of control points that effects the register transfer.
  - 5. Assemble the control logic
- MIPS makes it easier
  - Instructions same size
  - Source registers always in same place
  - Immediates same size, location
  - Operations always on registers/immediates
- Single cycle datapath => CPI=1, CCT => long
- Next time: implementing control

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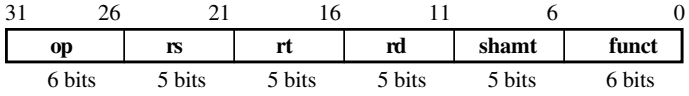
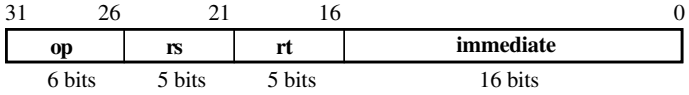
Recap: The MIPS Instruction Formats

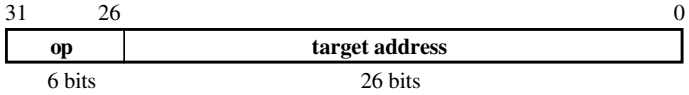
- All MIPS instructions are 32 bits long. The three instruction formats:
  - R-type



  - I-type



  - J-type



- The different fields are:
  - op: operation of the instruction
  - rs, rt, rd: the source and destination registers specifier
  - shamt: shift amount
  - funct: selects the variant of the operation in the "op" field
  - address / immediate: address offset or immediate value
  - target address: target address of the jump instruction

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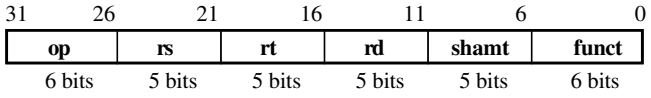
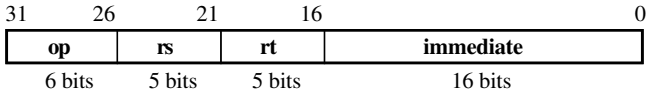
Recap: The MIPS Subset

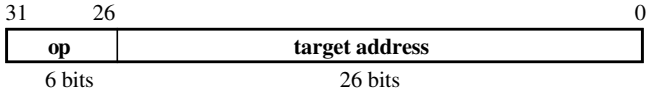
- ADD and subtract
  - add rd, rs, rt
  - sub rd, rs, rt



- OR Imm:
  - ori rt, rs, imm16



- LOAD and STORE
  - lw rt, rs, imm16
  - sw rt, rs, imm16
- BRANCH:
  - beq rs, rt, imm16
- JUMP:
  - j target



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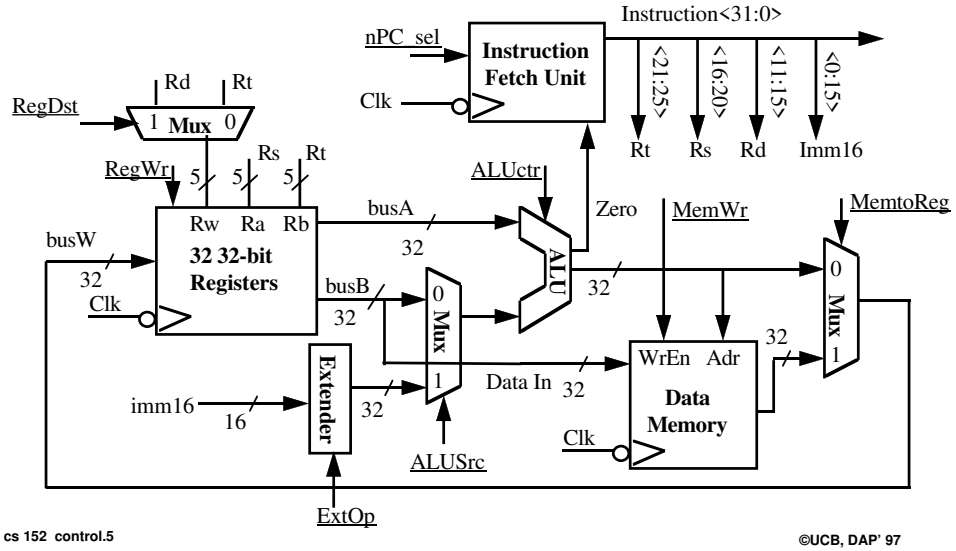
Recap: A Single Cycle Datapath

- We have everything except control signals (underline)
  - Today's lecture will show you how to generate the control signals



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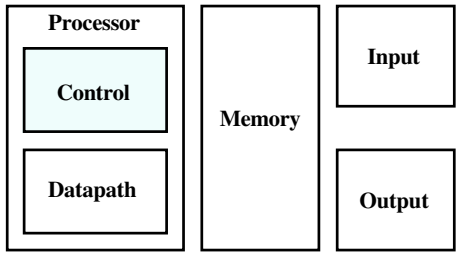
The Big Picture: Where are We Now?

- The Five Classic Components of a Computer



- Today's Topic: Designing the Control for the Single Cycle Datapath

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Outline of Today's Lecture

- Recap and Introduction (10 minutes)
- Control for Register-Register & Or Immediate instructions (10 minutes)
- Questions and Administrative Matters (5 minutes)
- Control signals for Load, Store, Branch, & Jump (15 minutes)
- Building a local controller: ALU Control (10 minutes)
- Break (5 minutes)
- The main controller (20 minutes)
- Summary (5 minutes)

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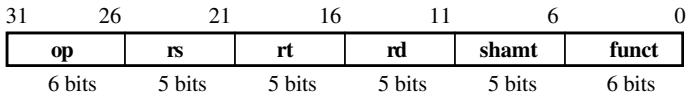
RTL: The `Add` Instruction



- add rd, rs, rt

• mem[PC]	Fetch the instruction from memory
• R[rd] <- R[rs] + R[rt]	The actual operation
• PC <- PC + 4	Calculate the next instruction's address

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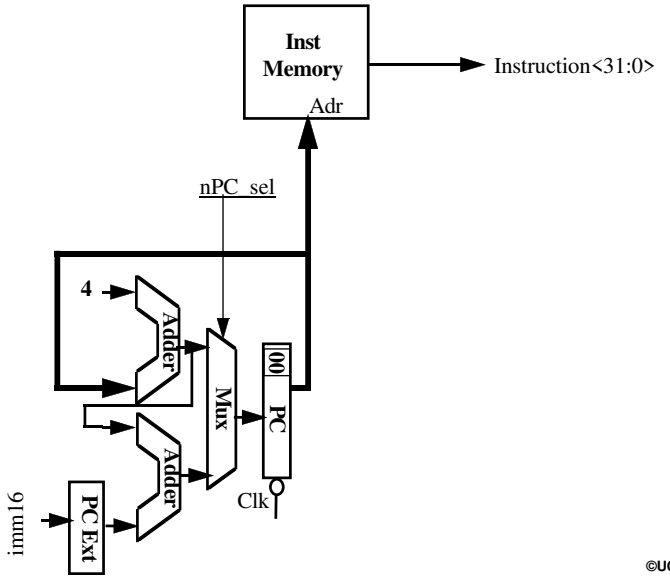
Instruction Fetch Unit at the Beginning of Add

- Fetch the instruction from Instruction memory: Instruction <- mem[PC]
  - This is the same for all instructions



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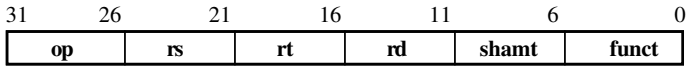
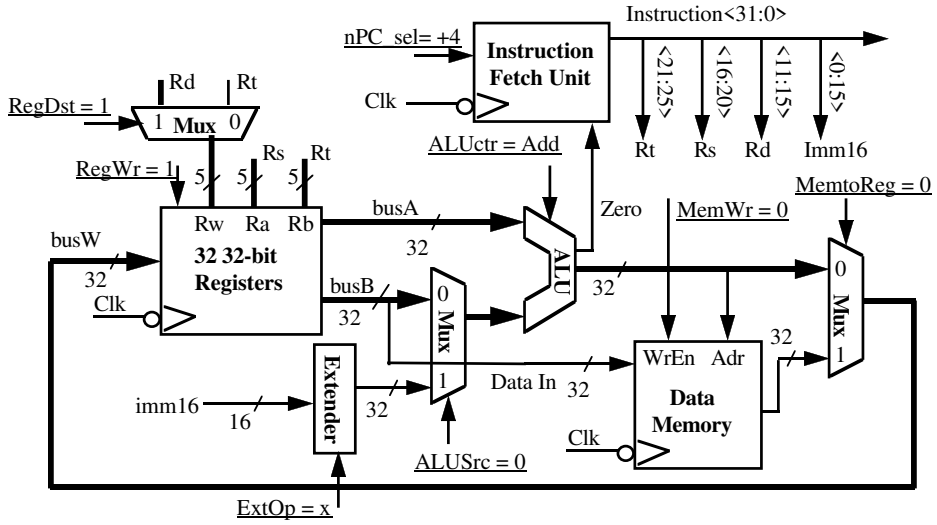
The Single Cycle Datapath during Add



- R[rd] <- R[rs] + R[rt]



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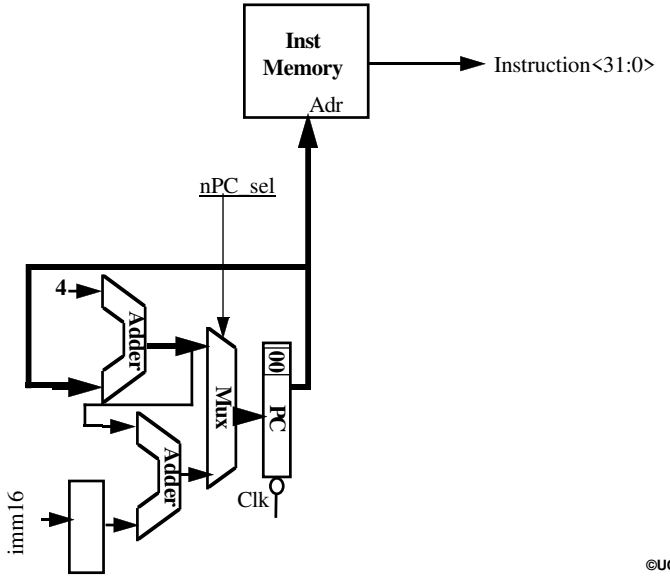
Instruction Fetch Unit at the End of Add

- PC <- PC + 4
  - This is the same for all instructions except: Branch and Jump



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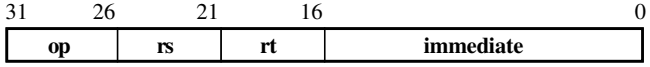
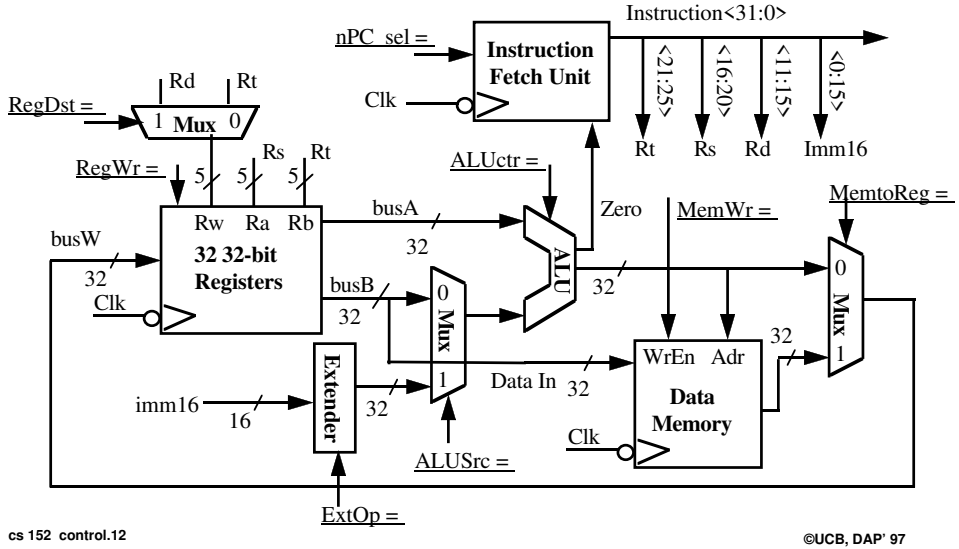
The Single Cycle Datapath during Or Immediate



- R[rt] <- R[rs] or ZeroExt[Imm16]



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Questions and Administrative Matters

- Midterm next Wednesday 10/8/97:
  - 5:30pm to 8:30pm, 306 Soda
  - No class on that day
- Midterm reminders:
  - Pencil, calculator, two 8.5" x 11" pages of handwritten notes
  - Sit in every other chair, every other row (odd row & odd seat)
- Meet at LaVal's pizza after the midterm
  - Need a headcount. How many are definitely coming?

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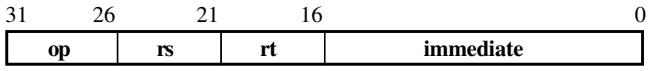
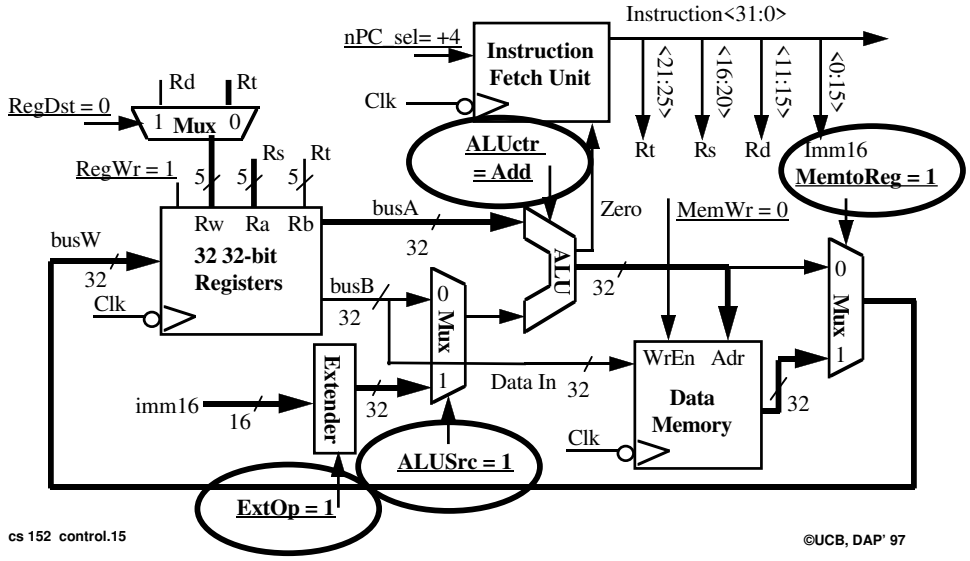
The Single Cycle Datapath during Load



- R[rt] <- Data Memory {R[rs] + SignExt[imm16]}



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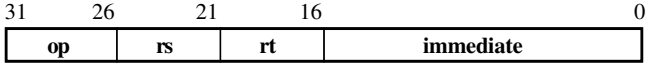
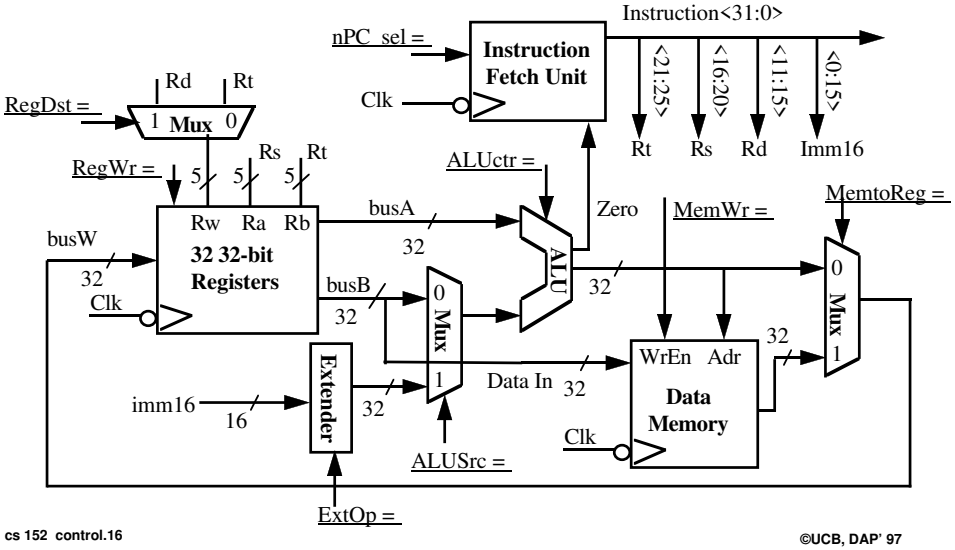
The Single Cycle Datapath during Store



- Data Memory {R[rs] + SignExt[imm16]} <- R[rt]



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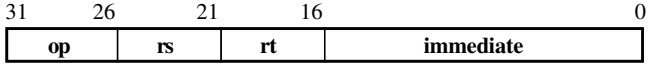
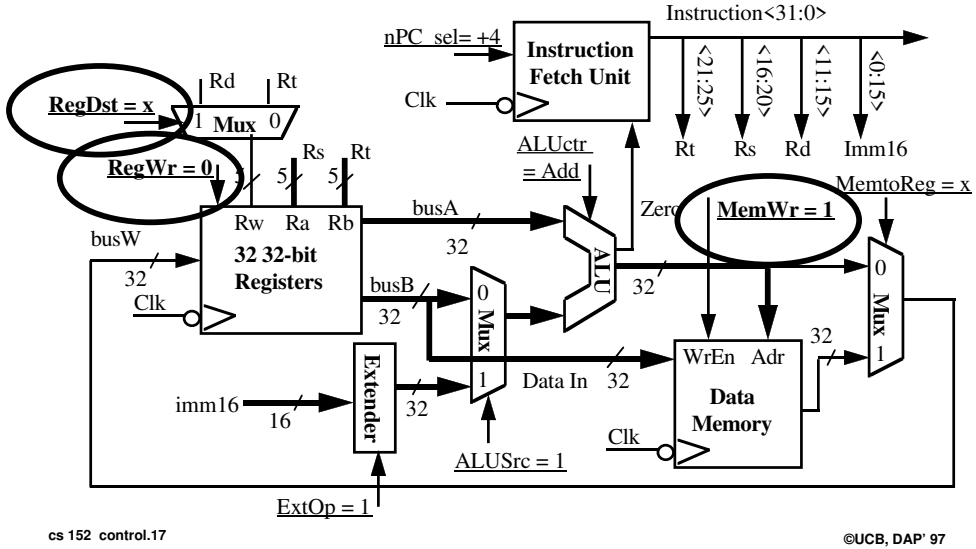
The Single Cycle Datapath during Store



- Data Memory {R[rs] + SignExt[imm16]} <- R[rt]



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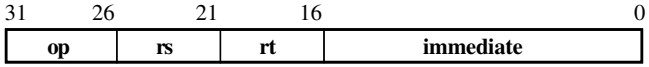
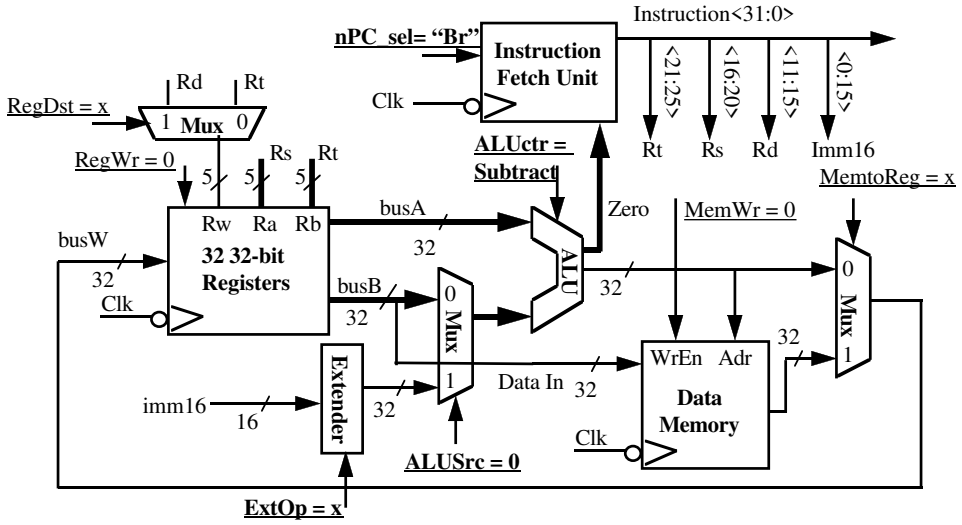
The Single Cycle Datapath during Branch



° if (R[rs] - R[rt] == 0) then Zero <- 1 ; else Zero <- 0



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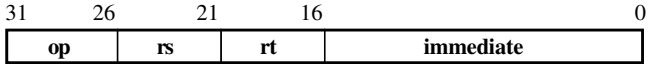
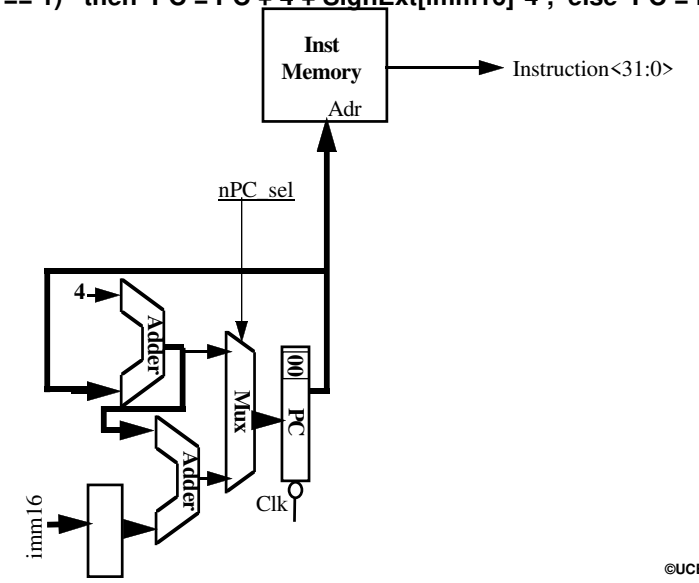
Instruction Fetch Unit at the End of Branch



° if (Zero == 1) then PC = PC + 4 + SignExt[imm16]*4 ; else PC = PC + 4



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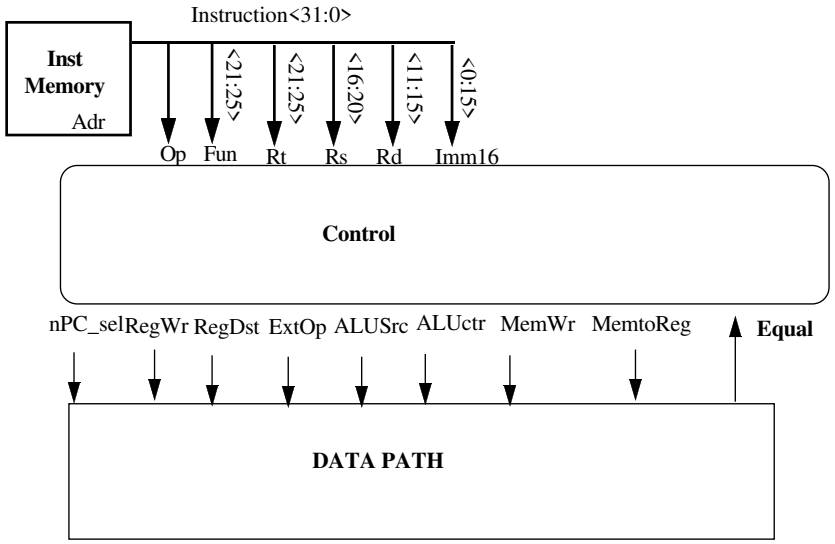
Step 4: Given Datapath: RTL -> Control



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A Summary of Control Signals

inst	Register Transfer	
ADD	R[rd] <– R[rs] + R[rt];	PC <– PC + 4
	ALUsrc = RegB, ALUctr = "add", RegDst = rd, RegWr, nPC_sel = "+4"	
SUB	R[rd] <– R[rs] – R[rt];	PC <– PC + 4
	ALUsrc = RegB, ALUctr = "sub", RegDst = rd, RegWr, nPC_sel = "+4"	
ORi	R[rt] <– R[rs] + zero_ext(Imm16);	PC <– PC + 4
	ALUsrc = Im, Extop = "Z", ALUctr = "or", RegDst = rt, RegWr, nPC_sel = "+4"	
LOAD	R[rt] <– MEM[R[rs] + sign_ext(Imm16)];	PC <– PC + 4
	ALUsrc = Im, Extop = "Sn", ALUctr = "add", MemtoReg, RegDst = rt, RegWr,	nPC_sel = "+4"
STORE	MEM[R[rs] + sign_ext(Imm16)] <– R[rs];	PC <– PC + 4
	ALUsrc = Im, Extop = "Sn", ALUctr = "add", MemWr, nPC_sel = "+4"	
BEQ	if (R[rs] == R[rt]) then PC <– PC + sign_ext(Imm16)] ‖ 00 else PC <– PC + 4	
	nPC_sel = "Br", ALUctr = "sub"	

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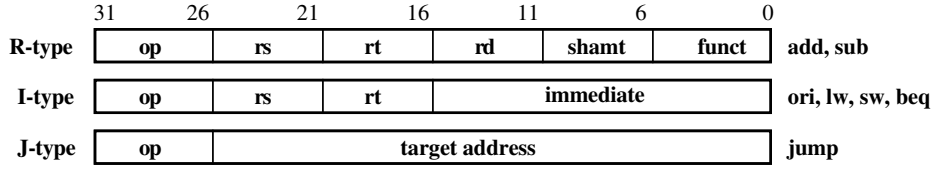
A Summary of the Control Signals

See Appendix A

	func	10 0000	10 0010	We Don't Care :-)				
	op	00 0000	00 0000	00 1101	10 0011	10 1011	00 0100	00 0010
		add	**sub**	**ori**	**lw**	**sw**	**beq**	**jump**
RegDst		1	1	0	0	x	x	x
ALUSrc		0	0	1	1	1	0	x
MemtoReg		0	0	0	1	x	x	x
RegWrite		1	1	1	1	0	0	0
MemWrite		0	0	0	0	1	0	0
nPCsel		0	0	0	0	0	1	0
Jump		0	0	0	0	0	0	1
ExtOp		x	x	0	1	1	x	x
ALUctr<2:0>		Add	Subtract	Or	Add	Add	Subtract	xxx



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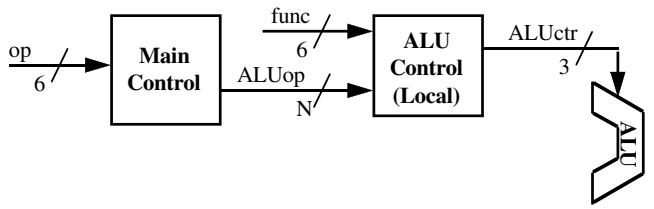
The Concept of Local Decoding

	op	00 0000	00 1101	10 0011	10 1011	00 0100	00 0010
		R-type	**ori**	**lw**	**sw**	**beq**	**jump**
RegDst		1	0	0	x	x	x
ALUSrc		0	1	1	1	0	x
MemtoReg		0	0	1	x	x	x
RegWrite		1	1	1	0	0	0
MemWrite		0	0	0	1	0	0
Branch		0	0	0	0	1	0
Jump		0	0	0	0	0	1
ExtOp		x	0	1	1	x	x
ALUop<N:0>		"R-type"	Or	Add	Add	Subtract	xxx



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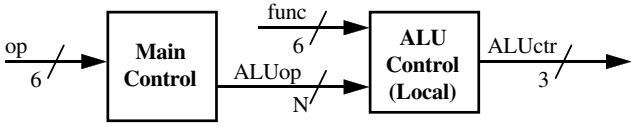
The Encoding of ALUop



- **In this exercise, ALUop has to be 2 bits wide to represent:**
  - **(1) "R-type" instructions**
  - **"I-type" instructions that require the ALU to perform:**
    - **(2) Or, (3) Add, and (4) Subtract**
- **To implement the full MIPS ISA, ALUop has to be 3 bits to represent:**
  - **(1) "R-type" instructions**
  - **"I-type" instructions that require the ALU to perform:**
    - **(2) Or, (3) Add, (4) Subtract, and (5) And (Example: andi)**

	R-type	ori	lw	sw	beq	jump
ALUop (Symbolic)	"R-type"	Or	Add	Add	Subtract	xxx
ALUop<2:0>	1 00	0 10	0 00	0 00	0 01	xxx

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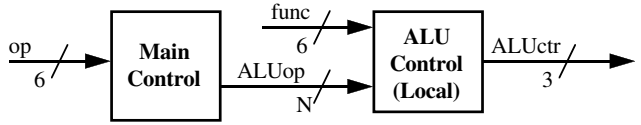
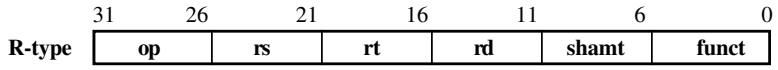
The Decoding of the "func" Field



	R-type	ori	lw	sw	beq	jump
ALUop (Symbolic)	"R-type"	Or	Add	Add	Subtract	xxx
ALUop<2:0>	1 00	0 10	0 00	0 00	0 01	xxx



funct<5:0>	Instruction Operation
10 0000	add
10 0010	subtract
10 0100	and
10 0101	or
10 1010	set-on-less-than

Recall ALU Homework (also P. 286 text):

ALUctr<2:0>	ALU Operation
000	Add
001	Subtract
010	And
110	Or
111	Set-on-less-than

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The Truth Table for ALUctr

ALUop (Symbolic)	R-type	ori	lw	sw	beq
	"R-type"	Or	Add	Add	Subtract
ALUop<2:0>	1 00	0 10	0 00	0 00	0 01

funct<3:0>	Instruction Op.
0000	add
0010	subtract
0100	and
0101	or
1010	set-on-less-than

ALUop			func				ALU Operation	ALUctr		
bit<2>	bit<1>	bit<0>	bit<3>	bit<2>	bit<1>	bit<0>		bit<2>	bit<1>	bit<0>
0	0	0	x	x	x	x	Add	0	1	0
0	x	1	x	x	x	x	Subtract	1	1	0
0	1	x	x	x	x	x	Or	0	0	1
1	x	x	0	0	0	0	Add	0	1	0
1	x	x	0	0	1	0	Subtract	1	1	0
1	x	x	0	1	0	0	And	0	0	0
1	x	x	0	1	0	1	Or	0	0	1
1	x	x	1	0	1	0	Set on <	1	1	1

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Break (5 Minutes)

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The Logic Equation for ALUctr<2>

ALUop			func				ALUctr<2>
bit<2>	bit<1>	bit<0>	bit<3>	bit<2>	bit<1>	bit<0>	
0	x	1	x	x	x	x	1
1	x	x	0	0	1	0	1
1	x	x	1	0	1	0	1

This makes func<3> a don't care

° ALUctr<2> = !ALUop<2> & ALUop<0> +
ALUop<2> & !func<2> & func<1> & !func<0>

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The Logic Equation for ALUctr<1>

ALUop			func				ALUctr<1>
bit<2>	bit<1>	bit<0>	bit<3>	bit<2>	bit<1>	bit<0>	
0	0	0	x	x	x	x	1
0	x	1	x	x	x	x	1
1	x	x	0	0	0	0	1
1	x	x	0	0	1	0	1
1	x	x	1	0	1	0	1

° ALUctr<1> = !ALUop<2> & !ALUop<0> +
ALUop<2> & !func<2> & !func<0>

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The Logic Equation for ALUctr<0>

ALUop			func				
bit<2>	bit<1>	bit<0>	bit<3>	bit<2>	bit<1>	bit<0>	ALUctr<0>
0	1	x	x	x	x	x	1
1	x	x	0	1	0	1	1
1	x	x	1	0	1	0	1

- ALUctr<0> = !ALUop<2> & ALUop<0>
  + ALUop<2> & !func<3> & func<2> & !func<1> & func<0>
  + ALUop<2> & func<3> & !func<2> & func<1> & !func<0>

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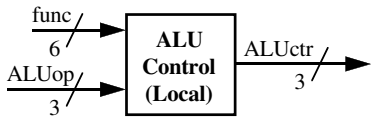
The ALU Control Block



- ALUctr<2> = !ALUop<2> & ALUop<0> +
  ALUop<2> & !func<2> & func<1> & !func<0>
- ALUctr<1> = !ALUop<2> & !ALUop<0> +
  ALUop<2> & !func<2> & !func<0>
- ALUctr<0> = !ALUop<2> & ALUop<0>
  + ALUop<2> & !func<3> & func<2> & !func<1> & func<0>
  + ALUop<2> & func<3> & !func<2> & func<1> & !func<0>

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Step 5: Logic for each control signal

- nPC_sel <= if (OP == BEQ) then EQUAL else 0
- ALUsrc <= if (OP == "Rtype") then "regB" else "immed"
- ALUctr <= if (OP == "Rtype") then funct
  elseif (OP == ORi) then "OR"
  elseif (OP == BEQ) then "sub"
  else "add"
- ExtOp <= _____
- MemWr <= _____
- MemtoReg <= _____
- RegWr: <=_____
- RegDst: <= _____

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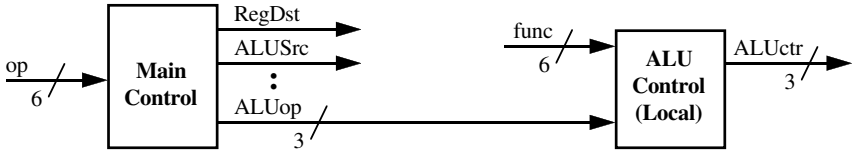
The "Truth Table" for the Main Control



op	00 0000	00 1101	10 0011	10 1011	00 0100	00 0010
	R-type	ori	lw	sw	beq	jump
RegDst	1	0	0	x	x	x
ALUSrc	0	1	1	1	0	x
MemtoReg	0	0	1	x	x	x
RegWrite	1	1	1	0	0	0
MemWrite	0	0	0	1	0	0
Branch	0	0	0	0	1	0
Jump	0	0	0	0	0	1
ExtOp	x	0	1	1	x	x
ALUop (Symbolic)	"R-type"	Or	Add	Add	Subtract	xxx
ALUop <2>	1	0	0	0	0	x
ALUop <1>	0	1	0	0	0	x
ALUop <0>	0	0	0	0	1	x

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The "Truth Table" for RegWrite

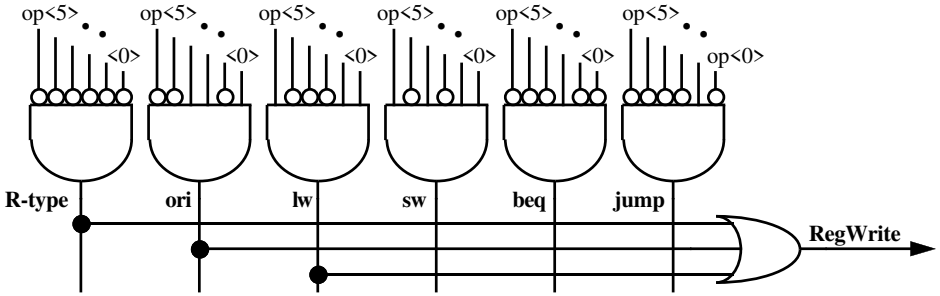
op	00 0000	00 1101	10 0011	10 1011	00 0100	00 0010
	R-type	ori	lw	sw	beq	jump
RegWrite	1	1	1	0	0	0

° RegWrite = R-type + ori + lw

= !op<5> & !op<4> & !op<3> & !op<2> & !op<1> & !op<0> (R-type)

+ !op<5> & !op<4> & op<3> & op<2> & !op<1> & op<0> (ori)

+ op<5> & !op<4> & !op<3> & !op<2> & op<1> & op<0> (lw)



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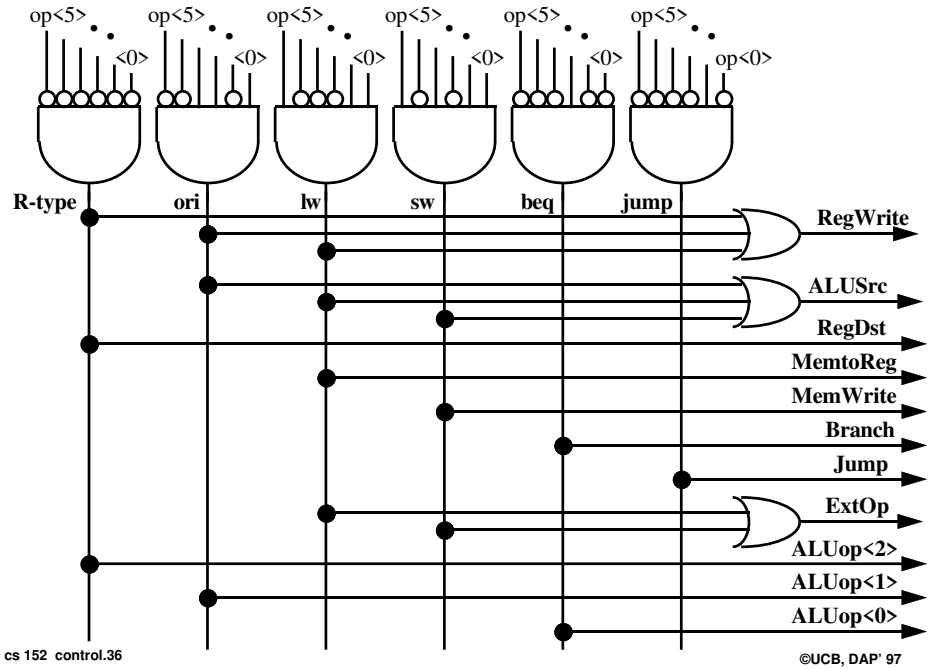
PLA Implementation of the Main Control



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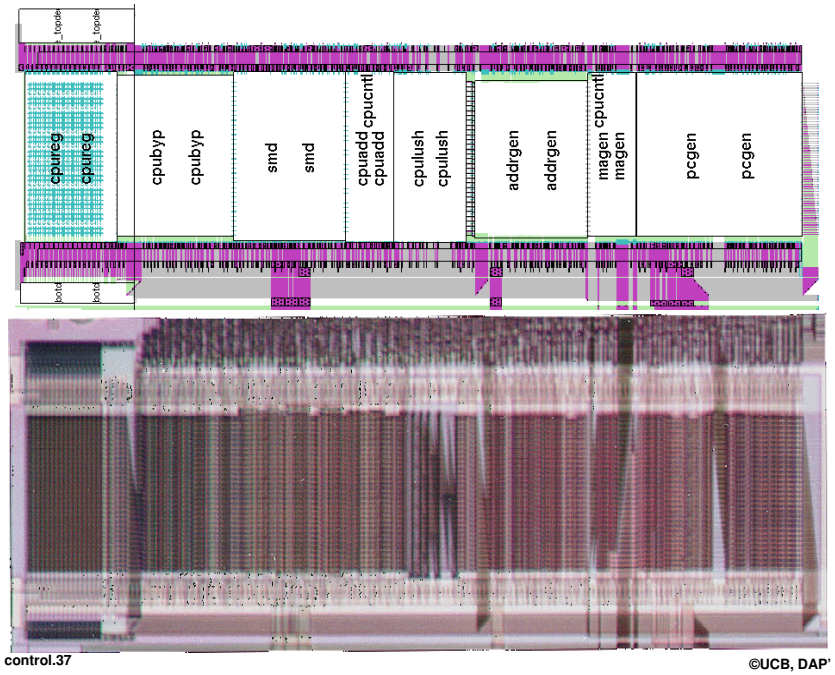
A Real MIPS Datapath (CNS T0)



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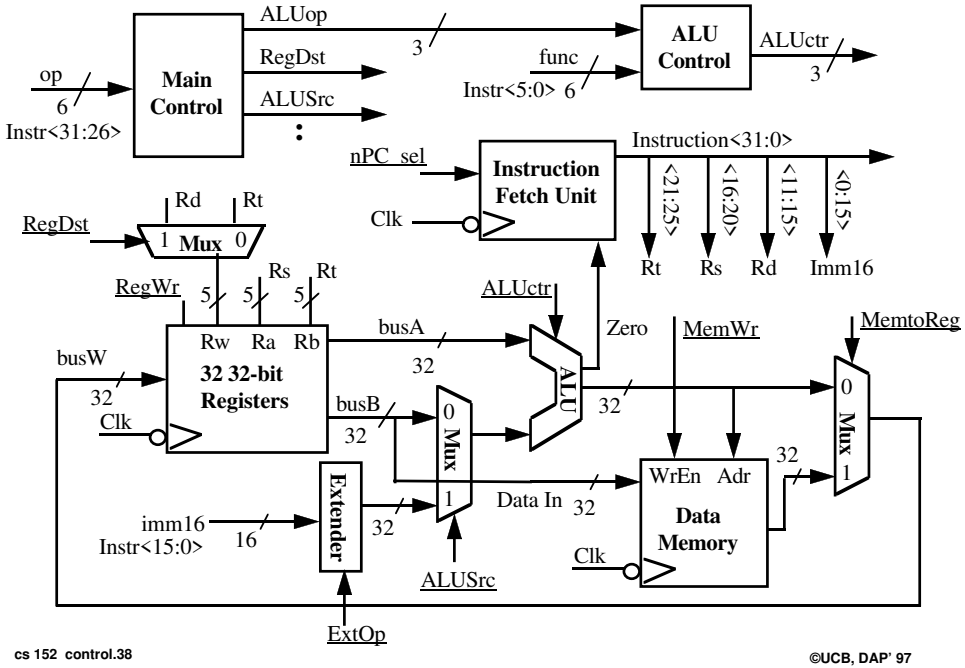
Putting it All Together: A Single Cycle Processor



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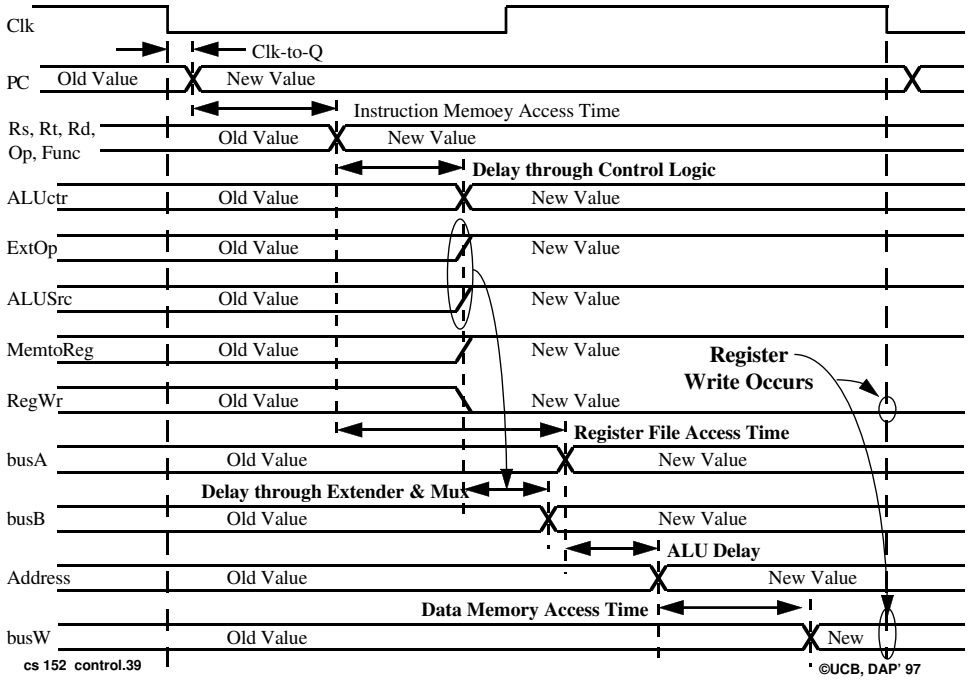
Worst Case Timing (Load)



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Drawback of this Single Cycle Processor

- Long cycle time:
  - Cycle time must be long enough for the load instruction:
    - PC's Clock -to-Q +
    - Instruction Memory Access Time +
    - Register File Access Time +
    - ALU Delay (address calculation) +
    - Data Memory Access Time +
    - Register File Setup Time +
    - Clock Skew
- Cycle time for load is much longer than needed for all other instructions

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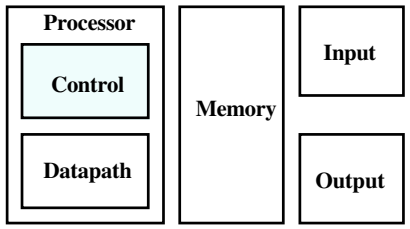
Summary

- Single cycle datapath => CPI=1, CCT => long
- 5 steps to design a processor
  - 1. Analyze instruction set => datapath requirements
  - 2. Select set of datapath components & establish clock methodology
  - 3. Assemble datapath meeting the requirements
  - 4. Analyze implementation of each instruction to determine setting of control points that effects the register transfer.
  - 5. Assemble the control logic
- Control is the hard part
- MIPS makes control easier
  - Instructions same size
  - Source registers always in same place
  - Immediates same size, location
  - Operations always on registers/immediates



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Where to get more information?

- Chapter 5.1 to 5.3 of your text book:
  - Daid Patterson and John Hennessy, "Computer Organization & Design: The Hardware / Software Interface," Second Edition, Morgan Kaufman Publishers, San Mateo, California, 1998.
- One of the best PhD thesis on processor design:
  - Manolis Katevenis, "Reduced Instruction Set Computer Architecture for VLSI," PhD Dissertation, EECS, U C Berkeley, 1982.
- For a reference on the MIPS architecture:
  - Gerry Kane, "MIPS RISC Architecture," Prentice Hall.

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