
CS61C - Machine Structures

Lecture 4 - C/Assembler Arithmetic and Memory Access

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Overview

- C operators, operands
- Variables in Assembly: Registers
- Comments in Assembly
- Addition and Subtraction in Assembly
- Memory Access in Assembly

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Review C Operators/Operands (1/2)

- **Operators:** +, -, *, /, % (mod);

$\text{¥}7/4=1, 7\%4==3$

- **Operands:**

- **Variables:** lower, upper, fahr, celsius

- **Constants:** 0, 1000, -17, 15.4

- **Assignment Statement:**

Variable = expression

- **Examples:**

```
celsius = 5*(fahr-32)/9;
```

```
a = b+c+d-e;
```

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C Operators/Operands (1/2)

- In C (and most High Level Languages) variables declared first and given a type

- **Example:**

```
int fahr, celsius;  
char a, b, c, d, e;
```

- Each variable can **ONLY** represent a value of the type it was declared as (cannot mix and match int and char variables).

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Assembly Design: Key Concepts

- **Keep it simple!**

- **Limit what can be a variable and what can't**

- **Limit types of operations that can be done to absolute minimum**

- if an operation can be decomposed into a simpler operation, don't include it

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Assembly Variables: Registers (1/4)

- Unlike HLL, assembly cannot use variables

- Why not? Keep Hardware Simple

- **Assembly Operands are registers**

- limited number of special locations built directly into the hardware

- operations can only be performed on these!

- **Benefit:** Since registers are directly in hardware, they are very fast

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Assembly Variables: Registers (2/4)

- Drawback: Since registers are in hardware, there are a predetermined number of them
 - Solution: MIPS code must be very carefully put together to efficiently use registers
- 32 registers in MIPS
 - Why 32? Smaller is faster
- Each MIPS register is 32 bits wide
 - Groups of 32 bits called a **word** in MIPS

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Assembly Variables: Registers (3/4)

- Registers are numbered from 0 to 31
- Each register can be referred to by number or name
- Number references:
 - \$0, \$1, \$2, \$30, \$31

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Assembly Variables: Registers (4/4)

- By convention, each register also has a name to make it easier to code
- For now:
 - \$16 - \$22 **£** \$s0 - \$s7
(correspond to C variables)
 - \$8 - \$15 **£** \$t0 - \$t7
(correspond to temporary variables)
- In general, use names to make your code more readable

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Comments in Assembly

- Another way to make your code more readable: comments!
- Hash (#) is used for MIPS comments
 - anything from hash mark to end of line is a comment and will be ignored
- Note: Different from C.
 - C comments have format /* comment */, so they can span many lines

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Assembly Instructions

- In assembly language, each statement (called an **Instruction**), executes exactly one of a short list of simple commands
- Unlike in C (and most other High Level Languages), each line of assembly code contains at most 1 instruction

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Addition and Subtraction (1/4)

- Syntax of Instructions:
 - 1 2,3,4
where:
 - 1) operation by name
 - 2) operand getting result ("destination")
 - 3) 1st operand for operation ("source1")
 - 4) 2nd operand for operation ("source2")
- Syntax is rigid:
 - 1 operator, 3 operands
 - Why? Keep Hardware simple via regularity

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Addition and Subtraction (2/4)

◦ Addition in Assembly

• **Example:** `add $s0, $s1, $s2` (in MIPS)

Equivalent to: `a = b + c` (in C)

where registers `$s0, $s1, $s2` are associated with variables `a, b, c`

◦ Subtraction in Assembly

• **Example:** `sub $s3, $s4, $s5` (in MIPS)

Equivalent to: `d = e - f` (in C)

where registers `$s3, $s4, $s5` are associated with variables `d, e, f`

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Addition and Subtraction (3/4)

◦ How do the following C statement?

`a = b + c + d - e;`

◦ Break into multiple instructions

`add $s0, $s1, $s2 # a = b + c`

`add $s0, $s0, $s3 # a = a + d`

`sub $s0, $s0, $s4 # a = a - e`

◦ Notice: A single line of C may break up into several lines of MIPS.

◦ Notice: Everything after the hash mark on each line is ignored (comments)

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Addition and Subtraction (4/4)

◦ How do we do this?

`∗f = (g + h) - (i + j);`

◦ Use intermediate temporary register

`add $s0, $s1, $s2 # f = g + h`

`add $t0, $s3, $s4 # t0 = i + j`

need to save `i+j`, but can't use
`f`, so use `t0`

`sub $s0, $s0, $t0 # f = (g+h) - (i+j)`

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Administrivia

◦ Project 1 due Midnight

◦ Lab 3: Your first MIPS program!

◦ HW 2 (due Mon 9/11) and HW3 (9/18) online and available

◦ Reading assignment:

• P&H 3.1-3.3, 3.5, 3.8 (page 145)

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Immediates

◦ Immediates are numerical constants.

◦ They appear often in code, so there are special instructions for them.

◦ Add Immediate:

`addi $s0, $s1, 10` (in MIPS)

`f = g + 10` (in C)

where registers `$s0, $s1` are associated with variables `f, g`

◦ Syntax similar to `add` instruction, except that last argument is a number instead of a register.

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Register Zero

◦ One particular immediate, the number zero (0), appears very often in code.

◦ So we define register zero (`$0` or `$zero`) to always have the value 0.

◦ This is defined in hardware, so an instruction like

`addi $0, $0, 5`

will not do anything.

◦ Use this register, it's very handy!

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Assembly Operands: Memory

- C variables map onto registers; what about large data structures like arrays?
- 1 of 5 components of a computer: memory contains such data structures
- But MIPS arithmetic instructions only operate on registers, never directly on memory.
- **Data transfer instructions** transfer data between registers and memory:
 - Memory to register
 - Register to memory

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Data Transfer: Memory to Reg (1/4)

- To transfer a word of data, we need to specify two things:
 - Register: specify this by number (0 - 31)
 - Memory address: more difficult
 - Think of memory as a single one-dimensional array, so we can address it simply by supplying a pointer to a memory address.
 - Other times, we want to be able to offset from this pointer.

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Data Transfer: Memory to Reg (2/4)

- To specify a memory address to copy from, specify two things:
 - A register which contains a pointer to memory
 - A numerical offset (in bytes)
- The desired memory address is the sum of these two values.
- Example: `8($t0)`
 - specifies the memory address pointed to by the value in `$t0`, plus 8 bytes

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Data Transfer: Memory to Reg (3/4)

- Load Instruction Syntax:
 - 1 2,3(4)
 - where
 - 1) operation name
 - 2) register that will receive value
 - 3) numerical offset in bytes
 - 4) register containing pointer to memory
- Instruction Name:
 - `lw` (meaning Load Word, so 32 bits or one word are loaded at a time)

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Data Transfer: Memory to Reg (4/4)

- Example: `lw $t0,12($s0)`

This instruction will take the pointer in `$s0`, add 12 bytes to it, and then load the value from the memory pointed to by this calculated sum into register `$t0`
- Notes:
 - `$s0` is called the **base register**
 - 12 is called the **offset**
 - offset is generally used in accessing elements of array: base reg points to beginning of array

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Data Transfer: Reg to Memory (1/2)

- Also want to store value from a register into memory
- Store instruction syntax is identical to Load instruction syntax
- Instruction Name:
 - `sw` (meaning Store Word, so 32 bits or one word are loaded at a time)

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Data Transfer: Reg to Memory (2/2)

◦ **Example:** `sw $t0,12($s0)`

This instruction will take the pointer in `$s0`, add 12 bytes to it, and then store the value from register `$t0` into the memory address pointed to by the calculated sum

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Pointers v. Values

◦ **Key Concept:** A register can hold any 32-bit value. That value can be a (signed) int, an unsigned int, a pointer (memory address), etc.

◦ If you write `add $t2,$t1,$t0` then `$t0` and `$t1` better contain values

◦ If you write `lw $t2,0($t0)` then `$t0` better contain a pointer

◦ Don't mix these up!

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Addressing: Byte vs. word

◦ Every word in memory has an **address**, similar to an index in an array

◦ Early computers numbered words like C numbers elements of an array:

`Memory[0], Memory[1], Memory[2], ...`

Called the "address" of a word

◦ Computers needed to access 8-bit **bytes** as well as words (4 bytes/word)

◦ Today machines address memory as bytes, hence word addresses differ by 4

`Memory[0], Memory[4], Memory[8], ...`

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Compilation with Memory

◦ What offset in `lw` to select `A[8]` in C?

◦ $4 \times 8 = 32$ to select `A[8]`: byte v. word

◦ Compile by hand using registers:

`g = h + A[8];`

• `g: $s1, h: $s2, $s3: base address of A`

◦ 1st transfer from memory to register:

`lw $t0, 32($s3) # $t0 gets A[8]`

• Add 32 to `$s3` to select `A[8]`, put into `$t0`

◦ Next add it to `h` and place in `g`

`add $s1, $s2, $t0 # $s1 = h + A[8]`

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Notes about Memory

◦ **Pitfall:** Forgetting that sequential word addresses in machines with byte addressing do not differ by 1.

• Many an assembly language programmer has toiled over errors made by assuming that the address of the next word can be found by incrementing the address in a register by 1 instead of by the word size in bytes.

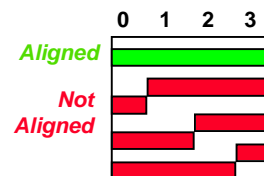
• So remember that for both `lw` and `sw`, the sum of the base address and the offset must be a multiple of 4 (to be **word aligned**)

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More Notes about Memory: Alignment

◦ MIPS requires that all words start at addresses that are multiples of 4 bytes



◦ Called **Alignment**: objects must fall on address that is multiple of their size.

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Role of Registers vs. Memory

- What if more variables than registers?
 - Compiler tries to keep most frequently used variable in registers
 - Writing less common to memory: **spilling**
- Why not keep all variables in memory?
 - Smaller is faster:
registers are faster than memory
 - Registers more versatile:
 - MIPS arithmetic instructions can read 2, operate on them, and write 1 per instruction
 - MIPS data transfer only read or write 1 operand per instruction, and no operation

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“And in Conclusion...” (1/2)

- In MIPS Assembly Language:
 - Registers replace C variables
 - One Instruction (simple operation) per line
 - Simpler is Better
 - Smaller is Faster
- Memory is **byte-addressable**, but `lw` and `sw` access one **word** at a time.
- A pointer (used by `lw` and `sw`) is just a memory address, so we can add to it or subtract from it (using offset).

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“And in Conclusion...” (2/2)

- New Instructions:
add, addi,
sub
lw, sw
- New Registers:
C Variables: `$s0 - $s7`
Temporary Variables: `$t0 - $t9`
Zero: `$zero`

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